The Language Parser

BNF-converter

February 8, 2010

This document was automatically generated by the *BNF-Converter*. It was generated together with the lexer, the parser, and the abstract syntax module, which guarantees that the document matches with the implementation of the language (provided no hand-hacking has taken place).

The lexical structure of Parser

Literals

Integer literals ⟨*Int*⟩ are nonempty sequences of digits.

CapIdent literals are recognized by the regular expression ⟨*upper*⟩(⟨*upper*⟩ | ⟨*digit*⟩ | '_') ∗ ⟨*lower*⟩(⟨*letter*⟩ | ⟨*digit*⟩ | '_')∗

LowIdent literals are recognized by the regular expression ⟨*lower*⟩(⟨*letter*⟩ | ⟨*digit*⟩ | '_')∗

Reserved words and symbols

The set of reserved words is the set of terminals appearing in the grammar. Those reserved words that consist of non-letter characters are called symbols, and they are treated in a different way from those that are similar to identifiers. The lexer follows rules familiar from languages like Haskell, C, and Java, including longest match and spacing conventions.

The reserved words used in Parser are the following:

```
FOR  PROC  var
```

The symbols used in Parser are the following:

```
:  =  ;
+  -  (
)
```

Comments

There are no single-line comments in the grammar.
There are no multiple-line comments in the grammar.

The syntactic structure of Parser

Non-terminals are enclosed between ⟨ and ⟩. The symbols ::= (production), | (union) and ϵ (empty rule) belong to the BNF notation. All other symbols are terminals.

⟨*ProgramSpec*⟩ ::= ⟨*ListFuncDecl*⟩

⟨*FuncDecl*⟩ ::= `PROC` ⟨*CapIdent*⟩ `:` ⟨*SingleStmtBlock*⟩

⟨*ListFuncDecl*⟩ ::= ϵ
| ⟨*FuncDecl*⟩ ⟨*ListFuncDecl*⟩

⟨*SingleStmtBlock*⟩ ::= ⟨*ListVar*⟩ ⟨*ListStmt*⟩

⟨*AssignOp*⟩ ::= =

⟨*Stmt*⟩ ::= ⟨*LowIdent*⟩ ⟨*AssignOp*⟩ ⟨*Expr*⟩
| ⟨*CapIdent*⟩
| ⟨*Expr*⟩
| ⟨*LowIdent*⟩ = ⟨*Expr*⟩ `FOR` ⟨*Expr*⟩ `:` ⟨*SingleStmtBlock*⟩

⟨*ListStmt*⟩ ::= ϵ
| ⟨*Stmt*⟩ `;` ⟨*ListStmt*⟩

⟨*Var*⟩ ::= `var` ⟨*LowIdent*⟩

⟨*ListVar*⟩ ::= ϵ
| ⟨*Var*⟩ `;` ⟨*ListVar*⟩

⟨*Expr1*⟩ ::= ⟨*Expr1*⟩ + ⟨*Expr2*⟩
| ⟨*Expr1*⟩ − ⟨*Expr2*⟩
| ⟨*Expr2*⟩

⟨*Expr3*⟩ ::= ⟨*Integer*⟩
| `(` ⟨*Expr*⟩ `)`

⟨*Expr*⟩ ::= ⟨*Expr1*⟩

$\langle Expr2 \rangle \quad ::= \quad \langle Expr3 \rangle$