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# Topflow, a Toolbox for Specialized Power System Analysis 

Master's thesis in Energy and Environmental Engineering<br>Supervisor: Olav Bjarte Fosso<br>July 2020

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## Abstract

Power system analysis is a branch of electrical engineering which is essential in designing electrical power systems. Simulations show if systems operate as expected, can withstand stress, and protect against failures. Many tools for power system analysis exist; one is a toolbox created by NTNU professor Olav Bjarte Fosso in the 90s. The toolbox consists of various FORTRAN-routines that do simulations such as Newton-Rapshon load-flow, continuation load-flow, contingency analysis, and security-constrained DC load-flow. Previous projects have worked towards making the toolbox available in Python, a popular object-orientated programming language. The speed and the functionality of the program is preserved by writing parts of the modernized code in C. The primary goal of this master's thesis is to update the simulationfunction "acsolve", which performs a Newton-Rapshon load-flow, and make sure it matches the performance of existing open-source toolboxes in Python.

The specialization project "Toolbox for Power System Analysis" [1] from 2019 by Åsmund Sælen purposed an initial implementation of acsolve. However, the implementation lacked functions for initializing the input-data; hence, the program could not run simulations on large power systems. This master's thesis presents Python-functions that communicates with Excelfiles, and a more sophisticated way of initializing the parameters of the simulation-functions. An additional algorithm named "decsolve", which performs fast-decoupled load-flows, is implemented. The updated program is tested on various MATPOWER standard-cases to study the code's reliability and performance. Profiling-tools available in Python are used to optimize the initial implementations, and the module "timit" is used to benchmarks the program against other open-source projects.

The result of the master's thesis is the toolbox "Topflow", which consists of translated load-flow functions that originate from the FORTRAN-code developed by Olav Bjarte Fosso. Topflow uses a combination of Python and C to boost the program's speed, and the functions use sparse methods to save memory and optimize the performance. The Newton-Rapshon load-flow function (acsolve) gives reliable results in all the tests, and the speed matches the open-source Python-projects "pypower" and "pandapower". The second simulation-function, decsolve, gave satisfactory results on small and medium-large power systems, but not on large. The conclusion of the master's thesis is that it reached the primary goal of updating acsolve, and gave recommendations on how to resolve the apparent bugs in decsolve. Future contributors can use the documentations in this thesis to complete the translation of the original toolbox.

## Samandrag

Kraftsystemanalyse er ein vesentleg del av å utvikle elektriske kraftsystem. Simuleringar viser om systema opererer som forventa og kan beskytte mot feil. Det finnast mange verktøy for å analysere kraftsystem, eit av dei er ein kjeldekode som NTNU professor Olav Bjarte Fosso utvikla på 90 -talet. Koden, som har fått namn «Topfllow», er bygd opp av ulike FORTRANrutinar som simulerer blant anna Newton-Rhapson last-flyt, CPF-analyser, utfalls-analyser og optimal DC last-flyt. Tidlegare prosjekt har arbeidt for å implementere Topflow i Python, eit mykje brukt objekt-orientert programmeringsspråk. Ytelsen og funksjonaliteten til det originale programmet er teke vare på ved å skrive delar av den moderniserte koden i C. Primærmålet med denne masteroppgåva er å oppdatere simuleringsfunksjonen "acsolve", som utøver NewtonRhapson last-flyt, og sørgjer for at den køyrer like effektivt som liknande modular tilgjengeleg i Python.

Spesialiserings prosjektet "Toolbox for Power System Analysis" [1] frå 2019 av Åsmund Sælen presenterte eit førsteutkast til oppdateringa av acsolve. Denne implementeringa mangla funksjonar for å setje input-data; noko som medførte at programmet ikkje var i stand til å køyre simuleringar på store kraftsystem. Denne masteroppgåva presentera funksjonar som kommuniserer med Excel-filer, og ein meir sofistikert måte å setje parameterane av simuleringsfunksjonane på. Topflow er også oppdatert med ein ny rutine, "decsolve", som utfører last-flyten FDLF. Det oppdaterte programmet er testa på fleire MATPOWER standard system for å undersøke kor påliteleg og rask koden er til å gjennomføre analysane. Tilgjengelege modular i Python er brukt for å optimalisere simularings- funksjonane.

Denne masteroppgåva resulterer i den oppdaterte versjonen av Topflow. Topflow nyttar no ein kombinasjon av Python og C for å auke programmet sin ytelse, saman med metoder som utnytter fordelane til glissne (sparse) matriser. Newton-Rhapson last-flyt funksjonen (acsolve) gjer pålitelege resultat på alle systema den er testa på, og den er minst like rask som dei eksisterande Python-programma "pypower" og "pandapower". Den andre simuleringsfunksjonen, decsolve, gjer tilfredsstillande resultat for små og middels store kraftsystem, men ikkje for store. Denne masteroppgåva konkluderer med at primærmålet, å oppdatere acsolve, er nådd. Oppgåva inneheld også tilrådingar for korleis framtidige prosjekt kan finne feila i rutinen decsolve. Framtidige bidragsytarar kan bruke dokumentasjonen i denne masteroppgåva i arbeidet mot å fullføre omsetjinga av den opphavlege FORTAN-koden, og ferdigstille Topflow.

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## Introduction

### 1.1 Previous work

This master's thesis contributes to the work of updating the toolbox for specialized power system analysis, which prof. Olav Bjarte Fosso developed in the 90s. Previous work includes implementations by Leif Warland at Statnett, the master thesis "Toolbox for Specialized Power System Analysis" by Hege Bruvik Kvandal [2], and the specialization project by Åsmund Sælen [1], which leads to this thesis.

The original toolbox contains multiple tools for analyzing power systems, such as NewtonRapshon load-flow, contingency analysis, continuous power-flow, and DC-optimal load-flow. Prof. Olav Bjarte Fosso wrote the code in FORTRAN, a high-performance programming language, which is suited for scientific computing. The toolbox uses sparse methods to save memory and enhance performance, which, combined with the qualities of FORTRAN, makes the program efficient and able to run on large systems.

Prof. Fosso designed the original program for study purposes only, however, Statnett showed interest in modernizing the code and making it available in Python. Leif Warland began the work of translating the code and laid the foundation for the new toolbox, which he named "Topflow". Hege Bruvik Kvandal continued the work in her specialization project [3] and master's thesis [2].

The focus of the specialization project by Åsmund Sælen [1] was to present a first update of the load-flow routine "acsolve" from the original toolbox. The implementations were tested on a 3-bus system, and compared with solved values form the lecture slides by prof Olav Bjarte Fosso [4], and hand-calculations. The test gave satisfactory results, which was an essential first step in translating the activity.

### 1.2 Scope and problem formulation

The most crucial goal of this thesis is to complete and verify the simulation-function acsolve. This routine performs a Newton-Rapshon load-flow, which is considered the "working horse" of power system analysis [4]. The specialization project by Åsmund Sælen [1] presented an initial update, and gave recommendations on how to make further contributions to the toolbox. These recommendations are listed below.

- Develop functions which initializes the parameters of acsolve
- Test the reliability of the code on larger power-systems
- Test the performance of the program
- Find a different solver for the differential equation-system
- Translate other activities from the original toolbox

This master's thesis builds on the recommendations from [1] to define clear milestones for the work. The milestones are listed below in order of priority.

1. Develop functions which initializes the parameters of acsolve
2. Make acsolve reliable on large systems
3. Find a solver which handles sparse matrices
4. Match the performance of similar open-source projects in Python
5. Translate the fast-decoupled load-flow routine "decsolve" from the original toolbox

Notice that milestone 4 tackles the "performance" of Topflow. The performance of a program may refer to several types of measurements; however, this thesis uses it interchangeably with the code's speed.

The rest of this thesis will discuss the work of completing the milestones above. Chapter 2 explains the basic theory that is needed to understand the work conducted in this master's thesis. Chapter 3 provides a user-guide, which gives the reader an overview of the program and shows the user-interface. Chapter 4 explains the design of the updated toolbox, and introduces the methods for testing the program's reliability and performance. Chapter 5 shows the resulting implementations, and aims to give future contributors a deeper understanding of the current toolbox. Chapter 6 provides the results of the reliability-tests and the speed-tests, and, finally, Chapter 7 concludes the master thesis.

## Theory

The sections 2.1-2.3, 2.5.7 and 2.5.9 are based on the theory-chapter in the specialization project by Åsmund Sælen from 2019[1].

### 2.1 Programming language

This thesis uses three different programming languages. The original code is written in Fortran 95, while the updated version uses C and Python. A vast amount of information is available online, and platforms such as YouTube and Tutorialspoint provide excellent tutorials for free . This section assumes that the reader is familiar with basics of programming, such as statements and loops, and will only give a brief introduction to each programming language. Python will be emphasised because it's the primary language of the updated program, Topflow.

### 2.1.1 Fortran

As one of the oldest programming languages, FORTRAN was created in the 1950s and developed for scientific calculations [5]. FORTRAN is still used, in particular, as a high-performance computing language. It's especially suitable for numerical analysis, which suits the techniques in the original toolbox.

A distinctive feature of FORTRAN is that routines can share data by using so-called "Commonblocks" (separate files that stores input data). If a routine changes a common-block parameter, this change applies to all other routines that share that common-block. Programs written in FORTRAN should minimize the use of this feature, since it may be a source of error that is difficult to resolve. The original toolbox uses common-blocks to initialize the parameters of various functions, an approach that the modernized code must handle differently.[6]

### 2.1.2 C

C is a general-purpose programming language created in 1972, and has since become one of the most used programming languages in the world [7]. It has many similarities to FORTRAN as it's a compiled language that uses many of the same default data types. C has also a thin abstraction layer and a low overhead compared with other high-level programming languages, which makes it useful for computationally intense programs. Programmers can, therefore, rewrite

FORTRAN-codes in C, and simultaneously keep the program close to its original design and performance.

### 2.1.3 Python

Different from Fortran and C, Python is an interpreted language which offers dynamically typing. There is no need for declarations of variables, because the data types are stated at runtime [8]. This feature, among others, make Python easy to learn and gives it intuitive syntax that facilitates productivity and readability. However, the flexibility of Python comes with lower performance compared with FORTRAN and C.

One of the most significant advantages of Python is the comprehensive standard library, and the fact that there are a vast number of third-party libraries and packages available. Among these, some modules make it possible to interface Python with other programming languages. A common way of boosting the performance of Python-programs, is to rewrite slow parts of the code in C and wrap them to pure Python (C_extensions). Topflow uses this technique to exploit the best of both worlds.[9]

## The SciPy stack

The SciPy stack is a collection of libraries that support scientific computing in Python. It consists the 6 open-source packages NumPy, SciPy library, Matplotlib, IPython, SymPy and pandas.[10],[11]

NumPy stands for "Numerical Python" and is a specialized library that enables numerical computing in Python. Its primary usage is to define and do operations on large arrays and matrices. The NumPy array objects benefit from a range of native functions, which rely on well-optimized C-codes. These arrays are, therefore, more efficient than the inbuilt Pythonlists.[12]

The SciPy library is another core package of the SciPy stack. It supports efficient numerical routines for linear algebra, with functions based on implemented standard algorithms [13]. In particular, Topflow relies on the function scipy.sparse.linalg.spsolve() to solve the sparse linear system $A x=b$.

### 2.2 C-extensions

Topflow uses Python as the user-interface, but the program relies on underlying C-codes to boost the performance of the functions. There exist several tools that makes C-codes available in Python (also called wrapping), including Python-C-Api, SWIG, Ctypes, and Cython [14]. The specialization project by Hege Bruvik Kvandal [3] studied these tools, and the following thesis by the same author [2] implemented C-extensions by wrapping the codes with Ctypes, a foreign function library for Python [15]. All the previous work on Topflow uses Ctypes, and the implementations presented in this thesis are no exceptions. The rest of this section explains the process of writing C-extensions with Ctypes.

### 2.2.1 Shared libraries

Ctypes uses shared libraries to export the C code and wrap them to Python. Shared libraries are files containing code that several modules or applications can use simultaneously. The benefits
are that, unlike static linking, the linker does not copy the code into all the modules, which saves memory and disk storage space [16]

On a Windows system, shared libraries are called "dynamic linked libraries" (DLL), hence a ".dll"-file is created when building the library. The documentation "Building C/C++ Extensions on Windows" on python.org [17] states the following on building DLLs: "Windows Python is built in Microsoft Visual C++; using other compilers may or may not work (though Borland seems to)". Based on that information, this project utilizes the Visual Stuido C/C++ compiler for creating DLLs.

Figure 2.1 shows an example of a C-file (operators.c), which contains the functions "multiply" and "add". "multiply" takes in two arguments of the type "double" (decimals) and multiplies the numbers, while "add" takes in two integers and returns their summation. The "__declspec(dllexport)" statement in front of the function declaration is required by Windows to export the code from the DLL.[16]

```
operators.c
#include <stdio.h>
    declspec(dllexport) double multiply(double a, double b)
{
    return a*b;
}
__declspec(dllexport) int add(int a, int b) {
    return a+b;
}
```

Figure 2.1: C-extension exported using __declspec(dllexport)
Topflow uses SCons, a free construction-tool [18], to build DLLs with the function SharedLibrary ('foo', ['fl.c', 'f2.c','f3.c',...]). The first argument of the function sets the name of the DLL, and the second specifies which file(s) to include. The file which contains the "SharedLibrary"function must be called "SConstruct.py" because SCons look for this particular filename when it builds the DLL [18]. Experiences from the specialization project [1] show that locating the SConstruct-file in the same directory as the C source files is the best practice, as it avoids linking problems. Figure 2.2 shows the SConstruct-file that put "add" and "multiply" from operators.c in a DLL named "cfunctions". The action that builds the DLL is to run the command "scons" in the terminal at the directory of the SConstruct-file. Figure 2.3 provides the output of this command, which shows that SCons successfully built the DLL "cfunctions", which now makes "add" and "multiply" accessible in Python. The next step is to help Python use these functions, by wrapping them to pure Python-code.

```
SConstruct.py
SharedLibrary(
    'cfunctions', 'operators.c'
)
```

Figure 2.2: The SConstruct file

```
PS C:\Users\Åsmund\PycharmProjects\acsolve> scons
scons: Reading SConscript files ...
scons: done reading SConscript files.
scons: Building targets ...
cl /Fooperators.obj /c operators.c /nologo
operators.c
link /nologo /dll /out:cfunctions.dll /implib:cfunctions.lib operators.obj
    Creating library cfunctions.lib and object cfunctions.exp
scons: done building targets.
```

Figure 2.3: Command prompt output from a successfully built DLL

### 2.2.2 Wrapper-functions

Ctypes provides C compatible data-types, and the ability to load DLLs and wrap the exported functions to pure Python. The wrapped functions are available from Python like any other callable. [15]

Figure 2.4 shows the Python-file "wrapperfunctions.py", which wraps the exported functions from the DLL "cfunctions" (defined in Figure 2.2). Wrapper-functions must specify the data types of arguments and return values, since Python uses dynamic typing and C does not. The example in Figure 2.4 starts off by importing ctypes and the required C-compatible data types, which in this case are c_double and c_int. The function "ctypes.cdll.LoadLibrary()" loads the DLL "cfunctions" and assigns it to the object "clib". "ctypes.cdll.LoadLibrary()" requires only one input-argument, which is the name or the path of the DLL. The script goes on to specify the data types of the wrapper functions. The function "add" uses integer-types for both its arguments and return values, while "multiply" uses double-types. The corresponding C compatible data types are "c_int" for integers and "c_double" for double. Finally, "wrapperfunctions.py" defines the python-functions "add" and "multiply", which are callable from the Python API. Figure 2.4 show how these functions take input arguments from Python, passes them on to the C-extension that does the work, and returns the result. The script "test.py" in Figure 2.5 tests the functions on a simple example. The output of the test, which is given in Figure 2.6, shows that both "add" and "multiply" are working as expected.

```
wrapperfunction.py
import ctypes
from ctypes import c_double
from ctypes import c_int
clib=ctypes.cdll.LoadLibrary('cfunctions')
clib.multiply.restype = c_double
clib.multiply.argtypes = [c_double, c_double]
clib.add.restype = c_int
clib.add.argtypes = [c_int, c_int]
def multiply(a,b):
    return clib.multiply(a,b)
def add(a,b):
    return clib.add(a,b)
```

Figure 2.4: The wrapper function

```
test.py
import wrapperfunction as wf
c=2
d=3
print('2 + 3 = ',wf.add(c,d))
print('2 * 3 = ',wf.multiply(c,d))
```

Figure 2.5: The Python file test.py testing the c-extensions "add" and "multiply"

```
PS C:\Users\Åsmund\PycharmProjects\acsolve> python test.py
2 + 3 = 5
2*3 = 6.0
```

Figure 2.6: Output of test.py
Ctypes supports NumPy array objects through the library Ctypeslib. In this project, the function numpy.ctypeslib.ndpointer is used to describe the return type and argument types of functions which uses NumPy arrays.[19]

### 2.3 Visual Studio Code

The free source editor Visual Studio Code (VSC) operates as the integrated development environment (IDE) for this project. It was recommended by the previous master's thesis [2] because it provides built-in support to a handful of programming languages, as well as extensions to others. This thesis uses the extensions for Python, C/C++, and Fortran, which provide support such as syntax highlighting and debugging. The IDE makes it possible to compare files side by side, which is very convenient when translating code between multiple languages. [20]

### 2.4 Black

Black is a code formatter for Python which structures the code automatically and makes the layout look the same regardless of the project [21]. The purpose of using Black in this project, is to make work consistent, and the collaboration as smooth as possible. The existing code by Leif Warland and Hege Bruvik Kvandal is formatted with Black, so is the updated program from this thesis. Visual Studio Code supports the formatter, and enables it to format the code every time a file is saved.

### 2.5 Power system analysis

Power system analysis is essential in the work of obtaining optimal operation of the existing power system, as well as planning expansions for the future. Topflow models the grid in a simplified way by line-diagrams and a per unit system, with parameters such as voltages, generated power and consumed power. Some of these parameters are known, others are calculated by
using well-established numerical techniques that gives good approximations. This section explains the theory behind the way Topflow models power systems and uses algorithms to perform simulations. The subsections $2.5 .6,2.5 .7$, and 2.5 .9 are from the specialization project [1] by Åsmund Sælen.

### 2.5.1 Fundamental electrical equations

This subsection contains four fundamental equalities, which the next subsections will use to derive electrical models and algorithms. (2.1) gives the relationship between the admittance $(Y)$, impedance $(Z)$, conductance $(G)$ and suscpetance $(B)$. Ohm's law (2.2) states that the voltage across two points on a conductor is proportional to the current through the conductor. (2.3) is Krichoff's current law, which expresses that the total current that flows into a node is equal to the total currents that flows out. Kirchoff's voltage law (KVL), (2.4), states that the sum of the voltages in a closed loop is zero. Finally, (2.5) specifies that complex power is the product of voltage and the conjugate of the current.

$$
\begin{gather*}
Y=\frac{1}{Z}=G+j B  \tag{2.1}\\
V=Z I  \tag{2.2}\\
\sum I_{i}=0  \tag{2.3}\\
\sum V_{i}=0  \tag{2.4}\\
S=P+j Q=V I^{*} \tag{2.5}
\end{gather*}
$$

### 2.5.2 Buses

The interconnection point between several components of the power system is called a "bus". A bus may be coupled to four different types of components: generators, loads, transmission lines and shunt elements. Figure 2.7 show the configuration of bus "i", which have all the four different component-types connected.


Figure 2.7: Bus "i"

Each bus in the system have four variables which is either known or unknown: voltage magnitude $(|V|)$, voltage angle $(\theta)$, active power injected $(P)$ and reactive power injected $(Q)$. More specifically, $P+j Q$ is the differences between the generated power $\left(P_{G}+j Q_{G}\right)$ and the demanded power $\left(P_{D}+j Q_{D}\right)$, so that $P=P_{G}-P_{D}$ and $Q=Q_{G}-Q_{D}$. The buses are classified into different types based on the combination of the known and unknown variables, as shown in Table 2.1. Which variables that are known depends on whether the bus connects to a load or a generator.

Loads consume a certain amount of active and reactive power, hence load-connected buses are classified as PQ-buses since both $P$ and $Q$ are known. Generators controls the voltage magnitude and the active generated power at the bus, hence generator-connected buses are classified as PV-buses. The reactive power at a generator, $Q_{G}$, is dependent on the topology of the system, since transmission lines, transformers and loads consume reactive power. Furthermore, reactive power is used to control voltages in the system, which is why generators may consume or produce reactive power. Buses such as the one in Figure 2.7, which has both a generator and a load connected, are classified as PV-buses since $Q=Q_{G}-Q_{D}$ is unknown.

The power system model uses an arbitrary PV-bus (often the largest) as a reference to all other buses, hence called "reference bus" or "slack bus". The voltage magnitude of the slack bus is set at 1.0 pu , and the angle at $0^{\circ}$. This bus holds a special role in the load-flow studies, as it balances the total power in the system and provides for the losses.

| Bus type | Known | Unknown |
| :--- | :--- | :--- |
| Generator (PV) | $P,\|V\|$ | $Q, \theta$ |
| Load (PQ) | $P, Q$ | $\|V\|, \theta$ |
| Slack | $\|V\|, \theta$ | $P, Q$ |

Table 2.1: Bus types.

### 2.5.3 Transmission lines

Power is transferred between buses through transmission lines. The representation of these lines depends on the length, broadly categorized into short, medium and long lines. The program presented in this thesis applies the same branch-model for all lines, the nominal- $\pi$ model, which is commonly used to represent medium length transmission-lines ( $80 \mathrm{~km}-240 \mathrm{~km}$ ).

Observing Figure 2.8, the line has an impedance of $Z_{i j}$ and a total line charging admittance, $Y_{S i j}$. By symmetry, the model lumps an equal portion of the total charging admittance on each end of the line. The real part of $Y_{S i j}=G_{S 12}+J B_{S i j}$ is very small, in most cases 0 , which is why $B_{S i j}$ often is the only given value for the shunt element.


Figure 2.8: $\pi$-equivalent transmission line

### 2.5.4 Transformers

Transformers are modeled in the same branch model as section 2.5.3, with some modifications due to the tap ratio and the shift-angle. A transformer with tap ratio $t$ and a phase-shifting angle $\theta_{s}$ is placed at the "from-end" of the line, as shown in Figure 2.9. The following derivation will show how to represent the transformer-branch on the same model as Figure 2.8.


Figure 2.9: Transformer-branch
KCL (2.3) is applied to Figure 2.9 to obtain expressions for $I_{1}$ (2.6) and $I_{2}$ (2.7).

$$
\begin{gather*}
\frac{I_{i}}{t e^{-j \theta_{s}}}=V_{i} t e^{j \theta_{s}}\left(\frac{Y_{S}^{*}}{2}+Y_{i j}^{*}\right)-V_{j} Y_{i j}^{*}  \tag{2.6}\\
I_{j}=-t e^{j \theta_{s}} Y_{i j}^{*} V_{i}+V_{j}\left(\frac{Y_{S}^{*}}{2}+Y_{i j}^{*}\right)-t e^{j \theta_{s}} V_{2} Y_{i j}^{*} \tag{2.7}
\end{gather*}
$$

This system of equation can be written as the augmented matrix shown in (2.8).

$$
\left[\begin{array}{c}
I_{i}  \tag{2.8}\\
I_{j}
\end{array}\right]=\left[\begin{array}{cc}
t^{2}\left(Y_{i j}^{*}+\frac{Y_{s}^{*}}{2}\right) & -t e^{-j \theta_{s}} Y_{i j}^{*} \\
-t e^{j \theta_{s}} Y_{i j}^{*} & Y_{i j}^{*}+\frac{Y_{s}^{*}}{2}
\end{array}\right]\left[\begin{array}{c}
V_{i} \\
V_{j}
\end{array}\right]
$$

(2.8) shows that Figure 2.9 reduces to the $\pi$-equivalent model in Figure 2.8 by substituting the following parameters:

$$
\begin{aligned}
& Z_{i j}=\frac{1}{t e^{-j \theta_{s}} Y_{i j}^{*}} \\
& Y_{S i j}=t^{2}\left(Y_{i j}^{*}+\frac{Y_{s}^{*}}{2}\right)-t e^{-j \theta_{s}} Y_{i j}^{*} \\
& Y_{S j i}=\frac{Y_{s}^{*}}{2}
\end{aligned}
$$

Notice that a transformer with tap ratio $t=1$ and shift angle $\theta_{s}=0$ is mathematically equivalent to a none-transformer branch.

### 2.5.5 Shunt Element

Shunt elements are inductors or capacitors connected to the power system to control the reactive power and thereby control the voltage. Line-losses increases if the loading level of the system is high, which leads to higher voltage drops. To prevent voltages below acceptable levels during such conditions, capacitor banks are switched on at weak busses to increase the voltage. These capacitors provide reactive power, which will increase the voltage.

Shunt elements can be connected either at a bus or on a line, but since reactive compensators have the greatest effect locally, these elements are most often bus-connected. This thesis converts data-sheets from MATPOWER [22] to initialize parameters of power systems. These sheets includes only bus-connected shunt elements, however, Topflow support both types of connections.

### 2.5.6 The problem formulation

Section 2.5.2 explained how buses are categorized by their known and unknown parameters. The goal of a power-flow study is to obtain good approximations for the unknown parameters by solving the so-called "load-flow equations". The derivation of these equations starts by applying Ohm's law (2.2) and Kirchoff's current law (2.3) to each bus in a system with $N$ buses, and thereby obtain the matrix equation-system (2.9).

$$
\left[\begin{array}{c}
I_{1}  \tag{2.9}\\
\vdots \\
I_{N}
\end{array}\right]=Y_{\text {bus }}\left[\begin{array}{c}
V_{1} \\
\vdots \\
V_{N}
\end{array}\right]
$$

$V_{i}$ is the voltage at bus $i, I_{i}$ is the sum of the line-currents flowing in and out of bus $i$, while $Y_{\text {bus }}$ is the $N \times N$ admittance matrix shown in (2.10):

$$
Y_{\text {bus }}=\left[\begin{array}{ccc}
Y_{11} & \ldots & Y_{1 n}  \tag{2.10}\\
\vdots & \ddots & \vdots \\
Y_{n 1} & \ldots & Y_{n n}
\end{array}\right]
$$

(2.11) shows that the diagonal elements of the $Y_{\text {bus }}$ are the sum of the line admittances and shunt admittances that are connected to the respective bus. The off-diagonal element $Y_{i j}$ is the
negative value of the line admittance between bus $i$ and $j$. In a large power system, most buses are not connected directly to another, which means that most element of the $Y_{\text {bus }}$ are zero. This is the definition of a sparse matrix.[23]

$$
\begin{gather*}
Y_{i i}=\sum_{i \neq j}^{n} y_{i j}  \tag{2.11}\\
Y_{i j}=-y_{i j} \tag{2.12}
\end{gather*}
$$

The load-flow equations (2.13)-(2.15) are finally obtained by substituting $I_{i}$ from (2.9) in the equation for complex power (2.5). The equations calculates the net power-injections at bus i. $G_{i j}, B_{i j}$ and $\theta_{i j}$ are the conductance, susceptance and phase shift between bus $i$ and $j$.

$$
\begin{gather*}
S_{i}=P_{i}+j Q_{i}=V_{i} \sum_{j}^{n} Y_{i j} V_{j}  \tag{2.13}\\
P_{i}=\left|V_{i}\right| \sum_{j}^{n}\left|V_{j}\right|\left(G_{i j} \cos \left(\theta_{i j}\right)+B_{i j} \sin \left(\theta_{i j}\right)\right)  \tag{2.14}\\
Q_{i}=\left|V_{i}\right| \sum_{j}^{n}\left|V_{j}\right|\left(G_{i j} \sin \left(\theta_{i j}\right)-B_{i j} \cos \left(\theta_{i j}\right)\right) \tag{2.15}
\end{gather*}
$$

The load-flow equations are non-linear, and therefore solved numerically by approximating the unknown parameters. A common approach is to approximate the voltages by solving the equations that contain known values for $P$ and $Q$. Section 2.5.7 and section 2.5.8 explains the theory of two algorithms that uses this approach to solve (2.14) and (2.15).

### 2.5.7 Newton Rapshon load flow

The first algorithm for solving the load-flow equations (2.14) and (2.15) is based on Newton's method. This method approximates the the solution of problem (2.16) through an iterative process. $f$ is a function defined for the variable $x$, and $c$ is a known constant. (2.17) shows Newton's method for the problem (2.16), and is derived from the first order Taylor-series for the function $f . f^{k}$ and $x^{k}$ are the approximations of $f$ and $x$ in the $k$ th iteration.[23]

$$
\begin{gather*}
f(x)=c  \tag{2.16}\\
x^{k+1}-x^{k}=\frac{c-f\left(x^{k}\right)}{f^{\prime}(x)} \tag{2.17}
\end{gather*}
$$

The method start with the iteration-counter $k=0$, and an initial guess of the value of $x: x^{0}$. (2.17) gives the formula for calculating the next approximation for $x: x^{1}$. The iterative process
continues and gives successively better approximations, until the error $c-f\left(x^{k}\right)$ is of adequate size. The speed of convergence is related to how close the initial guess is to the final solution.

The Newton-Rapshon load-flow (2.19) is a multi-dimensional version of (2.17), which solves the matrix equation system (2.18). In this version, $f$ is a vector that consists of the known active $(P)$ and reactive $(Q)$ power injections in a power system. $x$ is a vector that contains the angles $(\theta)$ and magnitudes $(\mid V)$ of the unknown voltages: $\left[\begin{array}{c}\theta \\ {[V]}\end{array}\right] . c$ is a vector of the scheduled active $\left(P_{s c h}\right)$ and reactive $\left(Q_{s c h}\right)$ powers and, finally, $J$ is the Jacobian matrix of the vector $f$. (2.20) show the definition of this matrix, which consists of the derivatives of (2.14) and (2.15) with respect to $\theta$ and $|V|$. The sub-matrices $\left[\frac{\partial P}{\partial \theta}\right],\left[\frac{\partial P}{\partial \mid V]}\right],\left[\frac{\partial Q}{\partial \theta}\right]$ and $\left[\frac{\partial Q}{\partial \mid V]}\right]$ are often referred to as $J 1, J 2, J 3$ and $J 4$.

$$
\begin{gather*}
{\left[\begin{array}{c}
P \\
Q
\end{array}\right]=\left[\begin{array}{c}
P_{s c h} \\
Q_{s c h}
\end{array}\right]}  \tag{2.18}\\
{\left[\begin{array}{c}
\theta^{k+1}-\theta^{k} \\
|V|^{k+1}-|V|^{k}
\end{array}\right]=J^{-1}\left[\begin{array}{l}
P_{s c h}-P^{k} \\
Q_{s c h}-Q^{k}
\end{array}\right]}  \tag{2.19}\\
J=\left[\begin{array}{ll}
\frac{\partial P}{\partial \theta} & \frac{\partial P}{\partial|V|} \\
\frac{\partial Q}{\partial \theta} & \frac{\partial Q}{\partial|V|}
\end{array}\right] \tag{2.20}
\end{gather*}
$$

The inverse of the Jacobian matrix can be found in a number of ways, including GaussJordan elimination and Gauss elimination [24]. However, the process of obtaining the inverse of a large matrix is time consuming, which is why programs commonly uses LU-factorization to solve (2.19). [4], [23]

Introducing $T_{i j}=G_{i j} \cos \left(\theta_{i j}\right)+B_{i j} \sin \left(\theta_{i j}\right)$ and $U_{i j}=G_{i j} \sin \left(\theta_{i j}\right)-B_{i j} \cos \left(\theta_{i j}\right)$ is convenient when obtaining the equations for calculating the derivatives in (2.20). The set of equations (2.21) shows how these derivatives are calculated. The variables $G_{i i}=\sum_{j \neq i}\left(G_{S i j}+G_{i j}\right)$ and $B_{i i}=B_{S i}+\sum_{j \neq i}\left(B_{S i j}+B_{i j}\right)$, where $G_{S i j}$ and $B_{S i j}$ are the conductance and susceptance of the line-shunt elements between bus $i$ and bus $j$. $B_{S i}$ is the imaginary part of the bus-connected shunt elements at bus $i$.

$$
\begin{array}{ll}
\frac{\partial P_{i}}{\partial \theta_{i}}=\left|V_{i}\right| \sum_{j \neq i}\left|V_{j}\right| U_{i j} & \frac{\partial P_{i}}{\partial\left|V_{i}\right|}=2\left|V_{i}\right| G_{i i}-\left|V_{i}\right| \sum_{j \neq i}\left|V_{j}\right| T_{i j} \\
\frac{\partial P_{i}}{\partial \theta_{j}}=-\left|V_{i}\right|\left|V_{j}\right| U_{i j} & \frac{\partial P_{i}}{\partial\left|V_{j}\right|}=-\left|V_{i}\right| T_{i j} \\
\frac{\partial Q_{i}}{\partial \theta_{i}}=-\left|V_{i}\right| \sum_{j \neq i}\left|V_{j}\right| T_{i j} & \frac{\partial Q_{i}}{\partial\left|V_{i}\right|}=2\left|V_{i}\right| B_{i i}-\left|V_{i}\right| \sum_{j \neq i}\left|V_{j}\right| U_{i j}  \tag{2.21}\\
\frac{\partial Q_{i}}{\partial \theta_{j}}=\left|V_{i}\right|\left|V_{j}\right| T_{i j} & \frac{\partial Q_{i}}{\partial\left|V_{j}\right|}=-\left|V_{i}\right| U_{i j}
\end{array}
$$

The number of elements in the vectors from (2.18) are dependent on the topology of the power system. In a system with the total number of $n$ buses and $g$ generators (including slack
bus), the number of equations of the type (2.14) is equal to $n-1$, while the number of equations of the type (2.15) is equal to $n-g$. It follows mathematically that $J$ is a $(2 n-g-1) \times(2 n-g-1)$ matrix.

Figure 2.10 shows an example of a three-bus system that consists of one PQ-bus, one PVbus, and one slack bus. The known parameters are given for each bus.


Figure 2.10: Three bus system
(2.22), (2.23) and (2.24) show how to write (2.18), (2.19) and (2.20) for the three-bus system in Figure 2.10 .

$$
\begin{gather*}
{\left[\begin{array}{c}
P_{1} \\
P_{2} \\
Q_{1}
\end{array}\right]=\left[\begin{array}{l}
P_{1(s c h)} \\
P_{2(s c h)} \\
Q_{1(s c h)}
\end{array}\right]}  \tag{2.22}\\
{\left[\begin{array}{c}
\theta_{1}^{k+1}-\theta_{1}^{k} \\
\theta_{2}^{k+1}-\theta_{2}^{k} \\
|V|_{1}^{k+1}-|V|_{1}^{k}
\end{array}\right]=J^{-1}\left[\begin{array}{c}
P_{(1) s c h}-P_{1}^{k} \\
P_{(2) s c h}-P_{2}^{k} \\
Q_{(1) s c h}-Q_{1}^{k}
\end{array}\right]}  \tag{2.23}\\
J=\left[\begin{array}{lll}
\frac{\partial P_{1}}{\partial \theta_{1}} & \frac{\partial P_{1}}{\partial \theta_{2}} & \frac{\partial P_{1}}{\partial V V_{1}} \\
\frac{\partial P_{2}}{\partial \theta_{1}} & \frac{\partial P_{2}}{\partial \theta_{2}} & \frac{\partial P_{2}}{\partial V \|_{1}} \\
\frac{\partial Q_{1}}{\partial \theta_{1}} & \frac{\partial Q_{1}}{\partial \theta_{2}} & \frac{\partial Q_{1}}{\partial|V|_{1}}
\end{array}\right] \tag{2.24}
\end{gather*}
$$

The advantages of the Newton-Rapshon method is that the region of convergence is large, and the number of iterations are few (if it converges). However, the speed of convergence is relatively slow if the initial guess is of poor choice. The Guass-Siedel method is another algorithm with the advantage of getting to the correct region fast, while it converges slow compared to Newton-Rapshon method. Some load-flow solutions are found by combining the two: using Guass-Siedel to get to the right region, followed by Newton-Rapshon to obtain the final solution. [4]

### 2.5.8 Fast-decoupled load-Flow

Another widely-used technique to solve the load-flow equations (2.14) and (2.15) is the Fastdecoupled load-low (FDLF). As the name suggests, it is faster and simpler than the regular Newton-Rapshon load flow, as it uses approximations to:

1. Decouple the system of equations into active $(P \theta)$ and reactive $(Q|V|)$ sub-problems.
2. Construct constant matrices which are independent of the voltage magnitudes and angles.

Despite apparent significant simplifications, the method has shown to perform remarkably well. The paper [25] showed that the success of the FDLF is related to the iteration scheme, which updates either angels or magnitudes after each sub-problem. Furthermore, it showed that some of the approximations commonly used to derive the FDLF are unnecessary due to the successive way of updating the magnitudes and angles. These approximations are still seen in the literature today, even after the publication of [25]. This section will explain three different versions of FDLF and give a brief overview of the findings in [25].

## Derivations

The derivation of the method starts with the first iteration of Newton-Rapshon using flat start, shown in (2.25)

$$
\left[\begin{array}{cc}
H & N  \tag{2.25}\\
M & L
\end{array}\right]\left[\begin{array}{c}
\Delta \theta \\
\Delta|V|
\end{array}\right]=\left[\begin{array}{c}
\Delta P \\
\Delta Q
\end{array}\right]
$$

The matrices $H, N, M, L$ corresponds to the four submatrices of the Jacobian matrix calculated at flat start $(|V|=1.0$ and $\theta=0)$. Note that the flat start results in $H$ and $L$ being the imaginary part of the admittance matrix $Y_{b u s}=G+J B$, in the structures of $\frac{\partial P}{\partial \theta}$ and $\frac{\partial Q}{\partial|V|}$ respectively.

The decoupling of the problem, so that $\Delta \theta$ and $\Delta|V|$ can be calculated separately, can be done in several ways:

1. The reactive power equations of (2.25) are subtracted by the active power equations, which are premultiplied with $M H^{-1}$ :

$$
\left[\begin{array}{cc}
H & N  \tag{2.26}\\
0 & L_{e q}
\end{array}\right]\left[\begin{array}{c}
\Delta \theta \\
\Delta|V|
\end{array}\right]=\left[\begin{array}{c}
\Delta P \\
\Delta Q-M H^{-1} \Delta P
\end{array}\right]
$$

2. The active power equations of (2.25) are subtracted by the reactive power equations, which are premultiplied with $N L^{-1}$ :

$$
\left[\begin{array}{cc}
H_{e q} & 0  \tag{2.27}\\
M & L
\end{array}\right]\left[\begin{array}{c}
\Delta \theta \\
\Delta|V|
\end{array}\right]=\left[\begin{array}{c}
\Delta P-N L^{-1} \Delta Q \\
\Delta Q
\end{array}\right]
$$

3. Both operations (1 and 2) are applied on system (2.25):

$$
\left[\begin{array}{cc}
H_{e q} & 0  \tag{2.28}\\
0 & L_{e q}
\end{array}\right]\left[\begin{array}{c}
\Delta \theta \\
\Delta|V|
\end{array}\right]=\left[\begin{array}{c}
\Delta P-N L^{-1} \Delta Q \\
\Delta Q-M H^{-1} \Delta P
\end{array}\right]
$$

Where the new equivalent matrices are given by:

$$
\begin{align*}
& H_{e q}=H-N L^{-1} M  \tag{2.29}\\
& L_{e q}=L-M H^{-1} N \tag{2.30}
\end{align*}
$$

## Solving the equation-systems

Algorithms can be used to solve the equation-systems (2.26) - (2.28) without introducing any major approximations. For example: the Primal-method solves (2.26) in the following way:

1. Initialize iteration count: $k=0$
2. Calculated the active power mismatches, and compute angle corrections: $\Delta \theta^{(k)}=\frac{\left.H^{-1} \Delta P^{(k)}|V|(k), \theta^{(k)}\right)}{|V|^{(k)}}$
3. Update the angles: $\theta^{(k+1)}=\theta^{(k)}+\Delta \theta^{(k)}$
4. Calculate the reactive power mismatches, and compute magnitude corrections:

$$
\Delta|V|^{(k)}=\frac{L_{e q}^{-1} \Delta Q^{(k)}\left(\mid V V^{(k)}, \theta^{(k+1)}\right)}{|V|^{k+1)}}
$$

5. Update the magnitudes: $|V|^{(k+1)}=|V|^{(k)}+\Delta|V|^{(k)}$
6. $k=k+1$, go to step 2 .

Multiple algorithms can solve the equation-system (2.26) - (2.28) in similar ways as the Primal-method. Appendix(A) provides two other versions, namely the Standard-method and the Dual-method. In contrast to the Primal-method, the Dual starts with a $Q \theta$-iteration and ends with a $P|V|$-iteration. The Dual-method uses the sub-matrices $H_{e q}$ and $L$, while the Standardmethod uses $H_{e q}$ and $L_{e q}$. The methods use approximations to obtain these matrices, which will be discussed in the next section.

## Approximations

It is essential to notice that the derivation of equations (2.26)-(2.28) uses no approximations, merely matrix operations. Each of these equations is, therefore, the same as the first iteration of Newton-Rapshon using flat-start. It's assumed that the voltage magnitudes and angles are close to the flat-start values, so that $|V| \approx 1 p u, \sin \theta \approx 0$ and $\cos \theta \approx 1$. These assumptions keep the sub-matrices constant throughout the load-flow.

The paper [25] showed that the impacts of the matrices M and N are automatically taken into account when solving the Primal-method successively by performing a $P \theta$ iteration and use the updated angles in the following $Q|V|$ sub-problem. The same is true for the Dual-method; therefore, there is no need to neglect these matrices, which was the standard in the traditional derivation of the FDLF methods. [25] [26] [27]

Submatrix $L_{e q}$ is the submatrix $L$ with $b_{i j}$ substituted by $1 / x_{i j}$ (which is why the primal method is often called the BX-version; it consists of $H$ and $L_{e q}$ ). This substitution is exactly true for radial systems and systems with constant $r / x$ ratios, because the operation performed to obtain (2.26) naturally cancel out the resistances in $L_{e q}$. It is a common mistake in the literature, even after the publication of [25], to assume that the resistances are neglected; this is only an apparent approximation. For systems which are not radial nor have constant $r / x$ ratios, on the other hand, the representation of $L_{e q}$ is an approximation, but still an excellent one. [25]

The same holds for $H_{e q}$, its obtained by calculating $H$ when $b_{i j}$ is substituted by $1 / x_{i j}$. Additional approximations need to be done, namely ignoring the effect of PV-buses and shunts. These approximations are why the Primal method performs better than the Dual in most cases.

## Summary

To summarize, the following approximations are made for the FDLF-mehods:

1. $\left|V_{i}\right|=1$
2. $\sin \left(\theta_{i j}\right)=0$
3. $\cos \left(\theta_{i j}\right)=1$
4. The effect of PV-buses and shunts ar neglected when forming $H_{e q}$ (XB- and XX-version only)
5. $L_{e q}$ is the submatrix $L$ calculated with $b_{i j}$ substituted by $1 / x_{i j}$
6. $H_{e q}$ is the submatrix $H$ calculated with $b_{i j}$ substituted by $1 / x_{i j}$
(Note: if the system is radial or have constant r/x ratios, 6. and 7. are not approximations, but exactly true.)

With the above approximations, the general FDLF-equations become (2.31). Table 2.2 show the meaning of $B^{\prime}$ and $B^{\prime \prime}$, along with the sequence of the $P \theta$ - and $Q|V|$-iterations.

$$
\begin{equation*}
\Delta \theta=\frac{B^{\prime-1} \Delta P}{|V|} \quad \Delta|V|=\frac{B^{\prime \prime-1} \Delta Q}{|V|} \tag{2.31}
\end{equation*}
$$

| Algorithm | $B^{\prime}$ | $B^{\prime \prime}$ | Iteration scheme |
| :--- | :--- | :--- | :--- |
| Primal / BX | $H$ | $L_{e q}$ | $P \theta-Q\|V\|$ |
| Dual / XB | $H_{e q}$ | $L$ | $Q\|V\|-P \theta$ |
| Standard / XX | $H_{e q}$ | $L_{e q}$ | $P \theta-Q\|V\|$ |

Table 2.2: $B^{\prime}$ and $B^{\prime \prime}$ for the different algorithms

It can be concluded from the material presented in this section that the resulting method is simpler and much faster than the Newton-Rapshon load flow, and it has a wide range of practical applications [26]. Despite its apparent significant simplifications, the method has shown to perform remarkably well. The paper [25] showed that the success of the decoupled method is very much related to the iteration scheme, where angels/magnitudes are updated after each sub-problem. Values which are apparently neglected in one iteration, are actually taken into account in the next. The various versions have shown to perform differently dependent on the system; on most systems the Primal method is favoured, however the Dual method has shown quick convergence in some cases where the Primal method did not. The Standard version is certainly the easiest of the three, since both sub-matrices are built only by the inverse of the line reactances. Knowing about the different versions of the FDLF gives the method more flexibility and greater performance, which is why all versions are implemented in the toolbox presented in this thesis.

### 2.5.9 Reactive power limitation

The operation of a synchronous generator is limited by three limits: the armature current heating, the field current heating, and the end region heating. Together they form the reactive capability chart of a generator, which show the area the generator can deliver power safely.[28]

Figure 2.11 shows a typical capability chart of a synchronous generator, which describes how the limitations discussed above restricts the reactive generation. The reactive generations of the PV-buses in a power system are unknown, and calculated in each iteration of a load-flow by estimating values for the voltages. These estimations might violate the capability requirements, if the limitations are not taken into account. A common approach to enforce reactive limits, is to consider PV-buses that violate the capability requirements as PQ buses, and set the reactive power at the limit. It is necessary to keep track on which buses are real PQ-buses and which are PV-buses (VAR-limited).pv_var_limit


Figure 2.11: Reactive capability chart of a synchronous generator

### 2.6 Sparse matrices

A matrix which mostly contains elements of zero value is called a sparse matrix. This phenomenon often appears in engineering and scientific computing, especially when solving partial differential equations. Building the full representation of a sparse matrix is inefficient, since the program will spend most of the execution time on storing and processing zeros. Most of the buses in a realistic power system are not connected directly to each other; hence the $Y_{\text {bus }}$ is sparse. Therefore, the advantages of exploiting the sparsity of such a matrix are of great importance for the code presented in this thesis. This section offers three different ways to store sparse matrices, along with their advantages and disadvantages. Figure 2.12 will be used as an example to explain the different methods.

### 2.6.1 Coordinated list (COO-format)

The first method consists of the three arrays Data, Row and Col, which are all of equal size. The names are quite self-explanatory; the Data-array contains the values of the non-zero elements, in any order, the Row- and Col-array specifies which row and column this element holds in the dense matrix. The arrays below shows Figure 2.12 on COO-format.
$\left[\begin{array}{llllll}1 & 0 & 0 & 0 & 0 & 2 \\ 0 & 0 & 3 & 0 & 0 & 0 \\ 0 & 4 & 0 & 0 & 7 & 0 \\ 0 & 0 & 0 & 9 & 8 & 0 \\ 5 & 0 & 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 & 0 & 6\end{array}\right]$

Figure 2.12: Sparse matrix example

$$
\begin{aligned}
\text { Data } & =\left[\begin{array}{lllllllll}
1 & 2 & 3 & 4 & 7 & 9 & 8 & 5 & 6
\end{array}\right] \\
\text { Row } & =\left[\begin{array}{lllllllll}
1 & 1 & 2 & 3 & 3 & 4 & 4 & 5 & 6
\end{array}\right] \\
\text { Col } & =\left[\begin{array}{llllllll}
1 & 6 & 3 & 2 & 5 & 4 & 5 & 1
\end{array}\right]
\end{aligned}
$$

Figure 2.13: Figure 2.12 on COO-format

The Python library scipy.sparse supports the coo-format with the class coo_matrix. The advantages of this format are efficient conversion to and from other sparse formats, especially CSC and CSR. The disadvantages are that the format does not support arithmetic operations or slicing. The intended usage of COO is to construct sparse matrices and facilitate fast conversion among other formats. It's convenient to construct the matrix of the equation $A x=b$ on this format before converting it, since different solvers require different formats.[29]

### 2.6.2 Compressed sparse formats (CSC- and CSR-format)

Formats for efficient access and matrix operation are stored on compressed sparse column (CSC) or compressed sparse row (CSR) formats. The CSC-format consist of three arrays: Data is the (top to bottom, then left to right) non-zero elements of the dense matrix, Row is the row-position of a given element in Data, Colptr holds one element per column, which specifies the position in Data where the given column starts. The arrays below show Figure 2.12 on CSC-format.[30]

$$
\begin{aligned}
\text { Data }= & {\left[\begin{array}{lllllllll}
1 & 2 & 3 & 4 & 7 & 9 & 8 & 5 & 6
\end{array}\right] } \\
\text { Row }= & =\left[\begin{array}{llllllll}
1 & 1 & 2 & 3 & 3 & 4 & 4 & 5
\end{array}\right] \\
& \text { Colptr }=\left[\begin{array}{llllll}
1 & 4 & 3 & 6 & 5 & 2
\end{array}\right]
\end{aligned}
$$

Figure 2.14: Figure 2.12 on CSC-format
The CSR-format is similar to the CSC-format, however this variant compresses the row information instead of the column information. Data is the (left to right, then top to bottom) non-zero elements of the dense matrix, Col is the column-position of a given element in Data, Rowptr holds one element per row, which specifies the position in Data where the given row starts. The arrays below show Figure 2.12 on CSR-format.[31]

$$
\begin{aligned}
\text { Data }= & {\left[\begin{array}{lllllllll}
1 & 2 & 3 & 4 & 7 & 9 & 8 & 5 & 6
\end{array}\right] } \\
\text { Col }= & {\left[\begin{array}{llllllll}
1 & 6 & 3 & 2 & 5 & 4 & 5 & 1
\end{array}\right] } \\
& \text { Rowptr }=\left[\begin{array}{llllll}
1 & 3 & 4 & 6 & 8 & 9
\end{array}\right]
\end{aligned}
$$

Figure 2.15: Figure 2.12 on CSR-format

Both CSC and CSR are supported in Python through the classes csc_matrix and csr_matrix from the library scipy.sparse. The advantages of these formats are efficient arithmetic operations and fast matrix-vector products. The CSC format facilitates efficient column slicing, while the CSR format provides efficient row slicing. The solver used in this toolbox requires the sparse representation of the Jacobian matrix to be on either CSC- or CSR-format. Modifications to matrices on these formats are on the other hand expensive, which is why it's common to use other formats to construct and store sparse matrices before converting them to CSR or CSC for arithmetic operations.[30][31]

## Topflow user guide

Topflow is a toolbox for power system analysis. The current program contains functions for running load-flows on power-systems, which the user can initialize by specifying the input-data in Excel-files. This chapter provides an overview of the toolbox and explains how to setup the program and make use of the implemented load-flow methods.

### 3.1 System requirements

The items needed to use the program are:

1. Python
2. Visual Studio C/C++ compiler

Topflow is only tested with the Visual Studio C/C++ compiler, other compilers may or may not work.

### 3.2 Installation

In addition to the system requirements, it is necessary to have various Python packages installed to run the code. These can be installed with the "pip install"-command in the cmd/terminal on a computer that has successfully installed Python. The full list of required packages, along with their installation-command, are listed below.

1. numpy: pip install numpy
2. scipy: pip install scipy
3. scons: pip install scons
4. openpyxl: pip install openpyxl

The end-user can skip these installations, and automatically manage the dependencies by installing Topflow. Since Topflow is not an open-source project, users must acquire the package from one of the contributors, and install it locally. Locally installations are done by feeding the full path of the package to pip:

```
pip install -e c:\\users\\aasel\\programmering\\topflow
```

Figure 3.1: How to install Topflow locally

### 3.3 Running load-flows

The application of Topflow is to run simulations and solve load-flow problems. The workflow of a load-flow study includes: 1) Preparing the input data in an Excel-file. 2) Initializing the parameters. 3) Solving the case with the proper load flow method(s). 4) Viewing the result of the simulation. The remaining subsections will go through these steps for a 14-bus example case.

### 3.3.1 Input Data

Topflow can only read Excel-files of a specific format. The format contains four separate spreadsheets, which the user must name as follows:

1. Case-identification
2. Bus-data
3. Generator-data
4. Line-data

The sheets must hold these names because this is how the routines that initialize the system identifies the different parameters. Common to all the spreadsheets is that they structure the data on a specific matrix-form, where the first row holds the names of the parameters. Similar to the spreadsheet-names, the user must name these parameters as specified in the next subsections.

## Case-identification

The first data record consists of only two parameters; IC and SBASE. IC $=0$ states that the case is a base case; hence simulations do not permanently change the data. If IC $=1$, the case is not a base case, and simulations may alter the system parameters. SBASE gives the system base in MVA.

Figure 3.2 show how the Case-identification sheet is structured as a $2 \times 2$ matrix. The routine that reads the sheet identifies the correct information by searching for "IC". The matrix may start in other cells than A1, but the matrix structure can not be changed.


Figure 3.2: The Case-identification data sheet for the 14-bus

## Bus-data

The bus-data record contains information about voltages, shunt elements, and loads at the buses. Table 3.1 gives a full list of the parameters. Notice how the consumed power for each bus (PLOAD and QLOAD) is present, while the generated power is not. That information belongs to the "Generator-data" spreadsheet.

| I | External bus numbers |
| :--- | :--- |
| NAME | Name of the buses |
| IDE | List of the buscodes for all buses; <br> 1-PQ, 2-PV, 3-Slack |
| PLOAD | Active load [MW] |
| QLOAD | Reactive load[MVAr] |
| GL | Active component of bus-connected shunt element [MW] |
| BL | Reactive component of bus-connected shunt element [MVAr] |
| AREA | Area postions for each bus |
| VM | Voltage magnetudes [pu] |
| VA | Voltage angles [radians] |
| PLOAD | Active load [MW] |
| QLOAD | Reactive load[MVAr] |
| BASEKV | Base voltage at each bus [kV] |
| ZONE | Zone postition for each bus |
| VMAX | Maximum voltage allowed [pu] |
| VMIN | Minimum voltage allowed [pu] |

Table 3.1: The parameters in the Bus-data record

Similarly to the Case-data, the user must type the bus-data in a matrix where the first row contains the name of the parameters. The routine that reads the bus-data sheet identifies the matrix by searching for "I", and reads the data as long as the column of I is not empty. The matrix may start at an arbitrary cell, but the structure must be of the one in Figure 3.3, and the cell below the last bus-number must be empty.

Topflow reads the information of each row consecutively and initializes the corresponding parameters. The program uses internal bus-numbers when performing operations, which simply are the sequence of which the busses were added to the program, starting at 0 . Conversions between external and internal bus-numbers are done by the functions ext2int and int2ext.

| - | A | A | B | C | D | E | F | G | H | 1 | J | K | L |  | M | N | 0 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | I |  | NAME | IDE | PLOAD | QLOAD | GL | BL | AREA | VM | VA | BASEKV | ZONE |  | Vmax | VMIN |  |
| 2 |  |  | 'Bus 1 | 3 | 0 | 0 | 0 | 0 | 1 | 1,06 | 0 | 0 |  | 1 | 1,06 | 0,94 |  |
| 3 |  |  | 'Bus 2 | 2 | 21,7 | 12,7 | 0 | 0 | 1 | 1,045 | -4,9826 | 0 |  | 1 | 1,06 | 0,94 |  |
| 4 |  |  | 'Bus 3 | 2 | 94,2 | 19 | 0 | 0 | 1 | 1,01 | -12,725 | 0 |  | 1 | 1,06 | 0,94 |  |
| 5 |  |  | 'Bus 4 | 1 | 47,8 | -3,9 | 0 | 0 | 1 | 1,018 | -10,313 | 0 |  | 1 | 1,06 | 0,94 |  |
| 6 |  |  | 'Bus 5 | 1 | 7,6 | 1,6 | 0 | 0 | 1 | 1,02 | -8,7739 | 0 |  | 1 | 1,06 | 0,94 |  |
| 7 |  |  | 'Bus 6 | 2 | 11,2 | 7,5 | 0 | 0 | 1 | 1,07 | $-14,221$ | 0 |  | 1 | 1,06 | 0,94 |  |
| 8 |  |  | 'Bus 7 | 1 | 0 | 0 | 0 | 0 | 1 | 1,062 | -13,36 | 0 |  | 1 | 1,06 | 0,94 |  |
| 9 |  |  | 'Bus 8 | 2 | 0 | 0 | 0 | 0 | 1 | 1,09 | -13,36 | 0 |  | 1 | 1,06 | 0,94 |  |
| 10 |  |  | 'Bus 9 | 1 | 29,5 | 16,6 | 0 | 19 | 1 | 1,056 | -14,939 | 0 |  | 1 | 1,06 | 0,94 |  |
| 11 |  |  | 'Bus 10 | 1 | 9 | 5,8 | 0 | 0 | 1 | 1,051 | -15,097 | 0 |  | 1 | 1,06 | 0,94 |  |
| 12 |  |  | 'Bus 11 | 1 | 3,5 | 1,8 | 0 | 0 | 1 | 1,057 | -14,791 | 0 |  | 1 | 1,06 | 0,94 |  |
| 13 |  | 12 | 'Bus 12 | 1 | 6,1 | 1,6 | 0 | 0 | 1 | 1,055 | -15,076 | 0 |  | 1 | 1,06 | 0,94 |  |
| 14 |  |  | 'Bus 13 | 1 | 13,5 | 5,8 | 0 | 0 | 1 | 1,05 | -15,156 | 0 |  | 1 | 1,06 | 0,94 |  |
| 15 |  | 4 ' | 'Bus 14 | 1 | 14,9 | 5 | 0 | 0 | 1 | 1,036 | -16,034 | 0 |  | 1 | 1,06 | 0,94 |  |
| 16 |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |  |
| , | , |  | Case-identification |  |  | Bus-data |  | Generator-data |  |  | Line-data | ${ }_{+}$ |  |  |  |  |  |

Figure 3.3: The Bus-data sheet for the 14-bus IEEE system

## Generator data

Table 3.2 gives the complete description of the generator-parameters. The user must add the data in a separate sheet on the form shown in Figure 3.4. The same directions for inserting the data applies for the generator-sheet: the matrix must be of correct structure, and the cell below the last bus-number must be empty.

| I | External bus-numbers |
| :--- | :--- |
| PG | Active power for generators [MW] |
| QG | Reactive power for generators [MVAr] |
| QMAX | Maximum reactive generation [MVAr] |
| QMIN | Minimum reactive generation [MVAr] |
| VS | Voltage setpoint for the generators [pu] |
| MBASE | MVA base for the machines (Default is MBASE $=$ SBASE) |
| STAT | Status of the generator (in service: 1, else: 0 ) |
| PMAX | Maximum active generation [MW] |
| PMIN | Minimum active generation [MW] |

Table 3.2: The parameters in the Generator-data record

|  | A |  | B | C | D | E | F | G | H | I | J | K |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | I |  | PG | QG | QMAX | QMIN | VS | MBASE | STAT | PMAX | PMIN |  |
| 2 | 1 |  | 232,4 | -16,9 | 10 | 0 | 1,06 | 100 |  | 332 | 0 |  |
| 3 | 2 |  | 40 | 42,4 | 50 | -40 | 1,045 | 100 |  | 140 | 0 |  |
| 4 | 3 |  | 0 | 23,4 | 40 | 0 | 1,01 | 100 |  | 100 | 0 |  |
| 5 | 6 |  | 0 | 12,2 | 24 | -6 | 1,07 | 100 |  | 100 | 0 |  |
| 6 | 8 |  | 0 | 17,4 | 24 | -6 | 1,09 | 100 |  | 100 | 0 |  |
| 7 |  |  |  |  |  |  |  |  |  |  |  |  |
|  | 1 | * |  | Case-identification |  |  | Bus-data | Generator-data |  |  | Line-data |  |

Figure 3.4: The Generator-data sheet for the 14 -bus IEEE system

## Line data

The last spreadsheet contains the data for the transmission-lines, which also includes information about transformers. Table 3.3 gives the full description of the line-parameters, and Figure 3.5 show how the data must is structured in a Excel-file.

| I | External bus-numbers (from bus) |
| :--- | :--- |
| J | External bus-numbers (to bus) |
| R | Line resistance [pu] |
| X | Line reactances [pu] |
| B | Total line charging susceptance [pu] |
| RATEA | phase A current rating [MVA] |
| RATEB | phase b current rating [MVA] |
| RATEC | phase C current rating [MVA] |
| GI | Conductance of line shunt placed at "I" end [pu] |
| BI | Suceptance of line shunt placed at "I" end [pu] |
| GJ | Conductance of line shunt placed at "J" end [pu] |
| BJ | Suceptance of line shunt placed at "J" end [pu] |
| RATIO | Ratio of the transformer |
| ANGLE | shift angle of transformers |
| STATUS | Indicates status for all lines (In service: 1) |

Table 3.3: The parameters in the Line-data record

| - | B |  | C | D | E | F | G | H | 1 | J | K | L | M | N | 0 | P | Q |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 1 | 1 | 」 | J | R | X | B | RATEA | RATEB | RATEC | GI | BI | GJ | BJ | RATIO | ANGLE | STAT |  |
| 2 |  | 1 | 2 | 0,01938 | 0,05917 | 0,0528 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 3 |  | 1 | 5 | 0,05403 | 0,22304 | 0,0492 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 4 |  | 2 | 3 | 0,04699 | 0,19797 | 0,0438 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 5 |  | 2 | 4 | 0,05811 | 0,17632 | 0,034 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 6 |  | 2 | 5 | 0,05695 | 0,17388 | 0,0346 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 7 |  | 3 | 4 | 0,06701 | 0,17103 | 0,0128 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 8 |  | 4 | 5 | 0,01335 | 0,04211 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 9 |  | 4 | 7 | 0 | 0,20912 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0,978 | 0 | 1 |  |
| 10 |  | 4 | 9 | 0 | 0,55618 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0,969 | 0 | 1 |  |
| 11 |  | 5 | 6 | 0 | 0,25202 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0,932 | 0 | 1 |  |
| 12 |  | 6 | 11 | 0,09498 | 0,1989 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 13 |  | 6 | 12 | 0,12291 | 0,25581 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 14 |  | 6 | 13 | 0,06615 | 0,13027 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 15 |  | 7 | 8 | 0 | 0,17615 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 16 |  | 7 | 9 | 0 | 0,11001 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 17 |  | 9 | 10 | 0,03181 | 0,0845 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 18 |  | 9 | 14 | 0,12711 | 0,27038 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 19 |  | 10 | 11 | 0,08205 | 0,19207 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 20 |  | 12 | 13 | 0,22092 | 0,19988 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 21 |  | 13 | 14 | 0,17093 | 0,34802 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 0 | 1 |  |
| 22 |  |  |  |  |  |  |  |  |  |  | Line-data |  |  |  |  |  |  |
|  | ¢ |  |  | Case-identification |  | Bus-data |  | Generator-data |  |  |  |  |  | $\oplus$ |  |  |  |

Figure 3.5: The Line-data sheet for the 14-bus IEEE system

### 3.3.2 Initializing a Case

At the core of the toolbox is the class "Case". A Case-object encapsulates all the system information needed in load flow studies and contains various class functions for initializing the data and performing arithmetic operations. Simulation-functions will typically take in an initialized Case-object to perform load-flows on the system. Figure 3.6 shows the easiest way of initializing a Case-object: to pass the path of the Excel-file that contains the input-data when creating an instance of the class.
>>> import topflow as tf
>>> case1 = tf.Case('C:/Users/aasel/programmering/topflow/topflow/example_cases/case14.xlsx')
Figure 3.6: How to load the input-data from the Excel-file that contains the 14-bus system

The object "case1" from Figure 3.6 now contains the system parameters for the 14-bus system. The user can verify this with the Case-class function "print_all", which prints all system parameters to the terminal.

```
>>> case1.print_all()
```

Figure 3.7: Print all the parameters of casel to screen

Another way of initializing a Case is through the function topflow.example_case(). This function grants access to several example-cases available in Topflow. The 14-bus system used in this user guide is one of them, and Figure 3.8 shows how to load that particularly system by its case-name: 'case14'. A full list of example cases is available in section 5.4.

```
>>> case1 = tf.example_case('case14')
```

Figure 3.8: How to load case14 with the function "topflow_example_case()"
Notice that only cases stored in the Topflow example-library are available through this method, customized cases must be loaded with the approach from Figure 3.6.

### 3.3.3 Solving the Case

The simulation-functions that perform analysis on the system takes in an initialized Case-object along with the desired settings. Default options apply if personal settings are not specified. The function "topflow.loadflow()" performs a loadflow on the system, and supports regular NewtonRapshon load flow as well as several fast-decoupled load-flow versions. Figure 3.9 shows how to run a Newton-Rapshon load-flow with default settings on case 1 from section 3.3.2.

```
>>> tf.loadflow(case1)
```

Figure 3.9: A regular Newton-Rapshon load-flow
The simulation-functions do not alter the input object in any way, which means that multiple functions can use the same Case-object as a base-case.

### 3.3.4 Accessing the Results

The result of a loadflow simulation is accessible through three different ways:

1. The result is pretty printed to the screen.
2. A Case-object containing the final results is returned from the simulation-functions.
3. The data of a simulation is saved to an Excel-file.

While 2. and 3. are optional, the result is by default printed to the screen, showing the number of iterations for convergence, and the execution time of the simulation. Figure 3.10 show the output of the simulation from subsection 3.3.3.

Convergence.
Number of iterations: 3 in 0.012997627258300781 seconds
Figure 3.10: A typical default terminal-output from running a load-flow

The loadflow-function from Figure 3.9 returns the result as a Case-object, which allows the user to use the data in subsequent analysis. Figure 3.11 show how the solved values of the loadflow-function are assigned to the instance "result1". The Case-class function "get()" makes it easy to access the parameters of that instance. get() takes in three arguments where the first specifies the component name (either 'bus', 'gen' / 'generator' or 'line'), the second is the component number(s), and the third is the parameter-name. It is possible to access multiple parameters with the same function call by defining the second argument as a list of numbers. The function will then return the corresponding parameters in the same order as the input-list. Figure 3.12 shows how to access the solved values in the instance return1 form fig:user ${ }_{g}$ uide ${ }_{a} s s i g n_{r}$ esult.Theparametersthatareaccessedare : voltagemagnitudeatbus 4 , theactivepowe
>>> result1 = tf.loadflow(case1)

Figure 3.11: How to assign the result-object to an instance

```
>>> bus4_vm = result1.get('bus', 4, 'vm')
>>> bus169_pd = result1.get('bus', [1,6,9], 'pd')
>>> gen_pg = result1.get('gen', 'all', 'pg')
```

Figure 3.12: Parameters can be accessed with the Case.get-function
The last way to access the result is by saving the solved data to an Excel-file. Solved cases that are assigned to an instance, such as "result1", can save the data with the Case-class-function "save 2 xl() ". Figure 3.13 shows how to save the instance from Figure 3.11 to an Excel-file.

```
>>> result1.save2xl('result1')
```

Figure 3.13: How to save the data of a result-instance to an Excel-file
Case.save2xl() has only one required argument, "filename", which specifies the name of the Excel-file. When calling the function, the file is by default saved in the current working
directory, and will automatically "pop-up" on the screen. Section 5.2 .3 will discuss available settings for saving the data.

Figure 3.13 shows another way to save the case; by specifying the filename when running a load-flow. Case.save 2 xl is then automatically invoked, and the resulting Excel-file will pop-up in the end of the simulation.

```
>>> tf.loadflow(case1, filename = 'result1')
```

Figure 3.14
This approach can provide information on all the iterations of the load-flow, in contrast to saving the Case-object returned by the loadflow-function, which only contains the final result.

### 3.3.5 Settings

Topflow allows the user to choose between a range of options, including choice of loadflow algorithm, various settings for saving the result to an Excel-file, and how much output the loadflow prints to the screen.

The easiest way to set the desired options is to type them in as keyword arguments in the simulation-function. For example, Figure 3.15 show how to run a fast decoupled loadflow (XX-version) on case1 (from Figure 3.8), with no printed output, saved in a file called "Result_FDXX". Topflow will save the file in the current working directory.

```
>>> result1 = tf.loadflow(case1, version = 'FDXX', print_verbose = 0, filename = 'Result_FDXX')
```

Figure 3.15: How to specify options in the loadflow-function
Alternatively, the settings can be set before the simulation by initializing a topflow.Settings() object. Settings() is a class which contains all the available options for running a loadflow, stored as instance-variables. By calling Settings() without any arguments, the object is initialized with the default settings. This class is convenient to use if the user wants to run multiple simulations with specific settings. Figure 3.16 show how to set personal settings by initializing a Settings() object, and use them in multiple loadflows.

```
>> import topflow as tf
>>>
>>> #Initialize the cases:
>>> case0 = tf.example case('case3')
>>> case1 = tf.example_case('case14')
>>> case2 = tf.example_case('case118')
>>>
>>> #Set personal settings:
>>> mysettings = tf.Settings(version = 'FDXX', print_verbose = 0)
>>>
>>> #Run the loadflows:
>>> result0 = tf.loadflow(case0, mysettings)
>>> result1 = tf.loadflow(case1, mysettings, print_verbose = 1)
Convergence. Number of iterations: 6.5
>>> result2 = tf.loadflow(case2, mysettings)
>>>
```

Figure 3.16: How to use a Settings-instance to specify options
Notice that if a setting is directly specified when running a loadflow, it will override the equivalent setting of the Settings() object. For that reason, the second loadflow in Figure 3.16 is the only one printing output to the screen.

## Method

There are many decisions to make when updating a program, such as the choice of programming language, the architecture of the code, and the user interface. Previous contributors decided to use Python as the administrative language and C-extensions to boost slow functions. The idea is to draw benefits from the flexibility of Python, and the speed of C. Existing implementations uses a well-established technique to interface Python with C, which section 2.2 explained in detail. The updated toolbox in this thesis uses the same technique, since it works well and makes the contribution consistent with the previous work.

The "topflow.loadflow()"-function introduced in the user-guide have two callable sub-activities; "acsolve" and "decsolve". acsolve is based on the Newton-Rapshon load-flow from subsection 2.5.7, and decsolve is based on the fast-decoupled load-flow from subsection 2.5.8. As stated in Chapter 1, the primary goal of this thesis is to ensure that acsolve is reliable, and matches the performance of similar activities in open-source Python-projects. This chapter will, therefore, emphasize on describing how acsolve was implemented and optimized.

### 4.1 The file structure

This project is not aiming to distribute an open-source toolbox, however it is structured as a Python-package ready for distribution. The structure corresponds with the Python Packaging Authority's (PyPA) guidelines on how to package a project [32]. The reason for choosing this structure is that it enables Topflow to be easily installed locally with "pip install", if the contributors share the project. It also makes it easier to distribute the toolbox in the future, if that is desired.


Figure 4.1: The top level structure of the toolbox
The folder "topflow" is the root of the repository, which is folder available from GitHub using the pull- or copy-command. The sub-directory, also called "topflow", is the package containing all the source files. This sub-directory also contains a "__init_.py" file, which is required
to import the directory as a package. This file also makes some functions available from the top level of the package. For example, instead of typing "topflow.loadflow.loadflow()" to access the function "loadflow()" from the file loadflow.py", the user can simply type "topflow.loadflow()".

The second sub-directory of the root, "tests", contains the tests-scripts of the project that are used to study the reliability and performance of the code. Section 4.3 and section 4.4 .1 will introduce these tests.

The "setup.py" file is the scripts that builds the distribution using "setuptools". At a minimum, it specifies the name, version, a short description of the package, and which source files to include. This file can also handle dependencies during installation.

Finally, the README.me is a text file that introduces and explains the package.

### 4.2 User interface

The specialization project by Åsmund Sælen [1] focused on translating the Newton-Rapshon algorithm, acsolve, from the original toolbox to Python and C. The updated code did not include functions for initializing the input data, nor a well-designed user interface. The script that tested acsolve initialized the data manually, and passed the parameters as arguments to the function. The first issue with this method is the way the parameters are initialized. The sub-functions of acsolve requires many parameters, and the total number of arguments would increase even more if the activity should include various options. To initialize all the parameters manually in separate objects is cumbersome, even for small systems. The second issue is related to how Python passes variables: functions with lists as arguments can change these globally. The user of the function can prevent this by passing a copy of the list, but again this is inconvenient to do for many objects.

This thesis solves the issues discussed above by introducing two new classes, namely "Case" and "Settings". A Case-object encapsulates all the information acsolve needs to perform a load-flow, while the Settings-object specifies options such as how much information the loadflow function prints to the terminal or saves to a file. The new classes reduces the number of arguments in acsolve to only two, and introduces new ways to initialize the parameters. The Case-class has functions that can read/write Excel-files on a specific format. This particular format is almost identical to the one used by MATPOWER, because it makes it easy to use the open-source MATPOWER test-cases to verify the implementations of Topflow.

### 4.3 Reliability

The primary goal of this thesis is to verify the Newton-Rapshon load-flow activity, acsolve. This activity creates the foundation of the toolbox, since other techniques relies on the same sub-functions. Other parts of the toolbox are easier to develop and verify, once the NewtonRapshon load-flow is established as a reliable routine.

### 4.3.1 Automated tests

This thesis uses automated tests to verify that the implemented functions are working as expected. The automated tests are functions which execute code, and check if specific assertions are true. These assertions are based on the fact that a known fixed input creates a know fixed output.

There are two types of automated tests: unit tests and integration tests. Unit tests check a single component, or in this case: a function. Integration tests, on the other hand, checks the interaction between multiple functions.[33] This project uses "pytest" as a test-runner. Pytest is a framework that has features such as detailed information about failed assert-statements and auto-discovery of test-scripts. This makes running automated tests efficient and easy.[34]

## Unit tests

Appendix C. 6 and C. 7 contains the unit tests that analyses the sub-functions in Topflow. The outcome of most functions are easy to validate by checking the result of a small example. The examples used to check the functions are imaginary, and customized to check if the function handles all relevant inputs as expected.

For example, Figure 4.2 shows the unit test for the function "maxmism", which finds the positions of the worst power mismatches. The first part of the test sets up the input data, which attempts to check all aspects of the function. The numpy-arrays "pinj" and "qinj" contain the active and reactive power mismatches. "mismloc[0]" and "mismloc[1]" stores the position of the worst mismatch in pinj and qinj respectively. "nbuses" is the number of buses in the system, and "buscod" are the bus-codes. Buscod shows that the first element in pinj and qinj is a PV-bus (2), the second is the slack-bus (3), the third is a PQ-bus (1), and the last is a VAr-limited PV-bus (-2).

```
import numpy as np
def test_maxmism():
    nbuses = 4
    buscod = np.array([2,3,1,-2])
    mismloc = np.array([0,0])
    pinj = np.array([0.0,150.0,-100.0,10.0])
    qinj = np.array([150.0, 0.0, 10.0, -100.0])
    for pqv in [1,2,3]:
        mismloc = np.array([0,0])
        maxmism(pqv, nbuses, buscod, mismloc, pinj, qinj)
        if(pqv == 1):
            assert_array_equal(mismloc, np.array([2,0]))
        elif(pqv == 2):
            assert_array_equal(mismloc, np.array([0,3]))
        else:
            assert_allclose(mismloc, np.array([2,3]))
```

Figure 4.2: The unit test for maxmism

The test in Figure 4.2 places the highest numbers of pinj (150.0) and qinj (also 150.0) at the slack-bus and a generator-bus, respectively. Those placements is to verify that maxmismat ignores these numbers correctly, because of the nature of the bus-types. Negative numbers are used as the worst mismatch in both qinj and pinj to check if maxmism can deal with worst absolute mismatches. The next part of the test calls the function maxmism. The test calls the function inside a for-loop to check all the different variants of the argument "pqv". This approach makes sure that the function also operates as expected for the fast-decoupled loadflow algorithm, "decsolve". pqv == 1: only the worst active mismatch is found, pqv == 2 : only the worst reactive mismatches is found, $\mathrm{pqv}==3$ : both active and reactive mismatches are found. The final stage is to check if the result of the function call is as expected. This is done
by the numpy-function "numpy.testing.assert_array_equal", which compares the known result and the worst mismatch-positions calculated by maxmism. The assertions are that the worst active power-mismatch is -100 at position 2 in pinj, and that the worst reactive power-mismatch is -100 at position 3 in qinj.

The known result of running maxmism on the system in Figure 4.2 can be seen beforehand by evaluating the input data. The same thing cannot be said about the results of more complex functions. Unit tests of such functions bases its assertions on the known result of performing a load-flow on the three-bus system in Figure 4.3. The specialization project [1] studied that system, and verified the result by hand-calculations and the lecture notes form the course ELK14 at NTNU [4]. For example, the unit test for the function which builds the Jacobian matrix (2.20) bases its assertions on the matrix calculated by [1] and [4].


Figure 4.3: Three bus system [1]
The system in Figure 4.3 is set up as a "fixture" in Python. Fixtures are functions that are recognized by pytest and initializes input-data for test functions. A fixture-object can be passed to a test-function as an input argument to provide the required data. Therefore, the data of the system in Figure 4.3 is only initialized once, and passed to multiple test-functions. The consistency of the input data makes the unit tests produce reliable results. [35]

## Integration tests

Integration tests check if the sub-functions of an algorithm interact as expected. They are essential to ensure a robust program that produces consistent, reliable results during the code's development. To run integration tests during development is an efficient way to see if the program still works as expected.

Appendix C. 5 contains the integration test for acsolve. The test takes in fixture-objects, which contains input-data of two power systems; the first is the three-bus system in Figure 4.3, and the second is a 14-bus system that origins from the MATPOWER-repository on GitHub [36]. The integration test bases the first system's assertions on the solution that the specialization project [1] provides. Section 6.2 will show that two different programs, namely pandapower and pypower, gives the same solution when running a Newton-Rapshon load-flow on the 14bus system. Topflow uses that solution to develop assertions for the integration test with the 14-bus system as input data.

### 4.3.2 Comparison tests

The tests in Appendix C. 2 and C. 3 evaluates the reliability of Topflow further by comparing it with open-source programs in Python, namely pandapower [37] and pypower [38]. The programs are tested on six standard cases from the MATPOWER repository on GitHub [22]. The cases origins from two different sources: "case14", "case30" and "case118" are IEEE test-cases from [39], while "case1354pegase", "case2869pegase" and "case9241pegase" stems from the Pan European Grid Advanced Simulation and State Estimation (PEGASE) [40], [41]. The digits in the case-names specify the number of buses in the system, which means that "case14" and "case9241pegase" are the smallest and largest systems respectively.

The test-script in Appendix C. 2 tests the reliability of acsolve by calling the load-flow functions ("loadflow" for Topflow, "runpp" for pandapower and "runpf" for pypower) for each program, and compare the calculated voltages. The functions use flat-start, and a convergence criteria of $10^{-8}$.

### 4.4 The design of the Python-C interface

Table 4.1 show the design of Topflow which was purposed in the specialization project [1]. The project argued to write functions that handle heavy calculations in C and keep the less timeconsuming functions in Python. Which functions that were considered "less time-consuming" was not well justified, since Topflow was not able to study large systems.

| Function | Environment |
| :--- | :--- |
| acsolve | Python |
| Sub-function | Environment |
| giibii | C |
| netinj | C |
| mismat | C |
| zerosp | Python |
| fmaxsp | Python |
| enfqlim | C |
| bujac | C |
| addel | C |
| numpy.linalg.solve | Python |
| t_u | C |
| jacpy | Python |

Table 4.1: The design purposed in [1]
The new Case-class (see section 4.2) makes it possible to initialize data and run load-flows on large power-systems. One of the goals of this thesis is to develop a toolbox that matches the performance of existing open-source programs for Python. This goal introduces a new motivation when deciding the Python-C interface, namely, to optimize the functions. The following sections describe the method used to decide the new design of the toolbox.

### 4.4.1 Optimiziation

Optimization is the process of modifying the code to increase the quality and efficiency of the program. Section 4.3 has already conducted the initial step of optimizing the code by writing automated tests. These tests detect bugs in the current program, but they also ensure that updates do not break the functionality of the code.

This thesis uses tools that profiles code in the process of making Topflow more efficient. Profiling gives information on the time consumed by various parts of the code, and can be used to decide which tasks to optimize. There are many profiling tools, and they give different information on the speed of the program. This project uses a combination of three tools to profile the program, namely "timeit", "cProfile" and "LineProfiler". These tools are all available in Python and can easily be installed as packages with the "pip install"-command.

### 4.4.2 cProfile

The first profiling tool is called cProfile. This module provides statistics of a function, such as how long the execution spends on various parts of the code, and how many times different sub-functions are called. The code snippet in Figure 4.4 is from the user manual of "The Python Profilers" [42], and shows how to profile a function which takes in a single argument.

```
import cProfile
import re
cProfile.run('re.compile("foo|bar")')
```

Figure 4.4: How to profile the function "re.compile()" with cProfile.run() [42]
The example in Figure 4.4 will run the function "re.compile()", and print the result to the screen. Figure 4.4 displays the output.


Figure 4.5: Output from running the code in Figure 4.4
Observing Figure 4.5, the sentence "order by: standard name" specifies that the filenames are ordered alphabetically by name, not by the order in which the functions are called. Table 4.2 gives a short description of the columns in the output from Figure 4.5.
cProfile provides a quick overview of the program's execution time by showing how much time each function spends. This is an excellent first step in locating "bottlenecks" in the program, which makes it easier to optimize the program. The module is designed to provide statistics on a given program, not for comparison with other programs (benchmarking).[42]

| Columns | Description |
| :--- | :--- |
| ncalls | number of calls |
| tottime | Total time spent in the given function <br> (exluding time spent in sub-functions) |
| percall | tottime devided by ncalls |
| cumtime | cumulative time spent in the given function <br> (including time spent in sub-functions) |
| percall | cumtime devided by ncalls |
| filename:lineno(function) | provides the respective data of each function |

Table 4.2: Description of the columns in the output of cProfile.run()

### 4.4.3 LineProfiler

The second profiling tool is "LineProfiler". Time-consuming parts of the code are sometimes not caused by sub-functions, but by actual lines in the code. This is especially true for programs that do scientific computing and uses libraries such as numpy. cProfiler is not able to identify these hotspots in the code, since the module only measures the execution time of functions calls. Line-profiler, on the other hand, measures the execution time of each line inside the given function, and gives therefore additional information about the code. A typical workflow should limit the use of LineProfiler to specific functions, since measuring and printing the execution time of each line in the code would be overwhelming. The approach of this thesis is to us cProfile as the first step in profiling the code, and LineProfiler as the second. [43]

Figure 4.6 shows the file "primes,py", which contains the function "primes()" that calculates the " n " first prime numbers. The function is decorated with @ profile, which tells LinePorifler to profile it. The function is called at the end of the script with $n=1000$ to initialize an instance.

```
@profile
def primes(n):
    A = [True] * (n+1)
    A[0] = False
    A[1] = False
    for i in range(2, int(n**0.5)):
        if A[i]:
            for j in range(i**2, n+1, i):
                A[j] = False
    return [x for x in range(2, n) if A[x]]
primes(1000)
```

Figure 4.6: The file primes.py, which has decorated the function primes()
The easiest way to run LineProfiler is to use the "kernprof"-script. This script comes with the installation of LineProfiler, and is convenient for running the package. Figure 4.7 shows the command that runs the kernprof-script, and profiles the decorated function in primes.py (see Figure 4.6).

The command from Figure 4.7 has to be run in the same directory as the relevant file (primes.py in this case), and the "-v"-option makes LineProfiler immediately print the result

```
kernprof -l -v primes.py
```

Figure 4.7: The command that invokes the kernprof-script
to the terminal. The output of the command is given in Figure 4.8. The results are formatted in a table of 5 columns, which are described in Table 4.3.

| Columns | Description |
| :--- | :--- |
| Line | Line number in the function |
| Hits | The number of times the line was executed |
| Time | Total execution-time of the line (in the timer's unit) |
| Per Hit | Execution-time of the line per call (in the timer's unit) |
| \%Time | Percentage of time spent on the line <br> relevant to the time spent on the whole function |
| Line Content | The actual code written on the line |

Table 4.3: Description of the columns in the output from running LineProfiler
Wrote profile results to line_profiler_ex.py.lprof
Timer unit: 1e-07 s
Total time: 0.0109534 s
File: line_profiler_ex.py
Function: primes at line 3

| Line \# | Hits | Time | Per Hit \% Time |  | Line Contents |
| :---: | :---: | :---: | :---: | :---: | :---: |
|  |  |  |  |  | ========== |
| 3 |  |  |  |  | @profile |
| 4 |  |  |  |  | def primes( $\mathrm{n}=1000$ ): |
| 5 | 1 | 60505.0 | 60505.0 | 55.2 | $\mathrm{A}=$ [True] * ( $\mathrm{n}+1$ ) |
| 6 | 1 | 43.0 | 43.0 | 0.0 | $\mathrm{A}[0]=$ False |
| 7 | 1 | 19.0 | 19.0 | 0.0 | A [1] = False |
| 8 | 30 | 569.0 | 19.0 | 0.5 | for i in range(2, int( $\mathrm{n}^{* *} 0.5$ ) : |
| 9 | 29 | 462.0 | 15.9 | 0.4 | if A[i]: |
| 10 | 1419 | 20163.0 | 14.2 | 18.4 | for $j$ in range( $\left.i^{* *} 2, n+1, i\right)$ : |
| 11 | 1409 | 22311.0 | 15.8 | 20.4 | A[j] = False |
| 12 |  |  |  |  |  |
| 13 | 1 | 5462.0 | 5462.0 | 5.0 | return [ x for x in range( $2, \mathrm{n}$ ) if $\mathrm{A}[\mathrm{x}]$ ] |

Figure 4.8: The output of profiling the prime-function
The first lines of the output gives information about the execution. This includes the "Timer unit", the total execution time of the code, as well as the file and the functions which were profiled. Timer unit specifies the conversion factor to seconds for the information provided in the table containing the results.

Observing the output in Figure 4.8, it can be seen that the most time-consuming lines are lines 7 and 8 . These lines should be the focus when optimizing the code. One way of improving the performance of this particular code is to use the NumPy-library, which provides faster arrayoperations compared with the inbuilt Python-lists.

### 4.4.4 Timeit

The last module for analyzing the performance of the program is "timeit". timeit contains timerelated functions, making it simple to measure the execution time of a specified part of the code.

The module avoid common traps for measuring execution times, and is more reliable than the inbuilt Python timer "time", and the modules cProfile and LineProfile [44]. The purpose of using timeit in this thesis, is to benchmark Topflow against pandapower and pypower.

Figure 4.9 show a simple example of how to use the timeit-module to measure the execution time of a sample code. The execution time of the code in Figure 4.9 is found by using the timeit.timeit() function. This function gives the user the opportunity to specify three input arguments. The keyword argument "stmt" stands for "statement", and is the actual code that the user wants to time. Timit runs the argument "setup" before the actual statement. timeit uses the setup to import the modules that the statement requires. The final argument, "number", refers to the number of times timeit.timeit() runs the code. A high number of measurements gives a more reliable result because it minimizes the influence of background processes. The default value of 1 million measurements can make the tests very slow, so there is a trade-off between the precision and speed of the test itself.

```
import timeit
#Code to executed only once
my_setup = '''
import numpy as np
list = np.array([0,1,2,3,4])
,',
#Code whose execution time is to be measured:
my_code = '''
for i in range(0,5):
    list = np.append(list, i)
print(timeit.timeit(setup = my_setup, stmt = my_code, number = 1000))
```

Figure 4.9: An example of a script which uses the timeit-module

### 4.4.5 Profiling acsolve

The focus in the process of optimizing Topflow is on increasing the performance of the function acsolve. This function performs a Newton-Rapshon load-flow, and its sub-routines creates the foundation for other techniques in the toolbox.

All the profiling in this section is done on the MATPOWER test-case case 1354pegase with default settings. The first stage in locating the most time-consuming part of acsolve, is to profile the code with cProfile.run(), which reveals the most time consuming sub-functions. The result is summarized in Table 4.4, which orders the relevant functions from the most time consumed to the least. Functions with a cumulative time of less than 0.005 s are not included.

Observing Table 4.4, the most time consuming functions are jacpy, which creates the full the Jacobian matrix (2.20), and "numpy.linalg.solve", which solves the differential equation system (2.19). This is no surprise, since the solver numpy.linalg.solve requires a dense matrix as input. A promising alternative to the numpy-solver is another solver available in Python via the Scipy-library, namely "scipy.sparse.linalg.spsolve()". This function solves a sparse linear system, and provides high performance by using the C library "UMFPACK". The input matrix, " $A$ ", should be a sparse matrix on CSC or CSR form to ensure efficiency. Because building the full Jacobian matrix is expensive, a new function called "COO_conv()" is implemented as a replacement for jacpy. This function returns a sparse matrix on COO-format. This format is used because NumPy provides fast conversion from COO format to CSR and CSC through the

| Function | Cumulative time in seconds |
| :--- | :--- |
| acsolve | 1.732 |
| Sub-function | Cumulative time in seconds |
| numpy.linalg.solve | 0.983 |
| jacpy | 0.648 |
| zerosp | 0.021 |
| maxmism | 0.019 |

Table 4.4: The most time consumin functions of acsolve
functions "to.csr()" and "to.csc". Having the Jacobian matrix on COO-format makes it more flexible, since it can efficiently switch the format required by a specific solver.[45],[46]

The other notably slow functions from Table 4.4 are "zerosp" and "maxmism". It was argued in the specialization project, [1], to write these functions in pure Python, because they are simple, and not expected to do complicated computations. The profiling with cProfile shows, on the other hand, that these functions are slower than the more complex routines that are Cextensions. One way of further increasing the code's performance is, therefore, to write all functions as C -extensions.

The impact of the modifications discussed above is analyzed by profiling the new code with cProfile on the same system and settings as the initial profiling-test. The results are summarized in Table 4.5.

| Function | Cumulative time in seconds (s) |
| :--- | :--- |
| acsolve | 0.069 |
| Sub-function | Cumulative time in seconds (s) |
| scipy.sparse.linalg.spsolve | 0.024 |
| to.csr() | 0.003 |
| coo_conv | 0.001 |
| zerosp | 0.000 |
| maxmism | 0.000 |

Table 4.5: The optimized functions

Comparing Table 4.4 and Table 4.5, the modifications have lowered the total execution time of acsolve with 1.663 seconds. This improvement means that the optimized code is 25 times faster than the original design. The increase in performance is mostly due to the sparse linear solver, whose performance is approx. $4000 \%$ higher than the NumPy-solver. The functions coo_conv() and csr_matrix() restructures the Jacobian matrix 162 times faster than jacpy, and the cumulative times of zerosp and maxmism are also improved significantly.

Table 4.5 shows that scipy.sparse.linalg.spsolve is the most time consuming function with a cumulative time of 0.024 s . The cumulative times of all the other functions are each lower than 0.005 s . In fact, the sum of all the sub-functions does not add up to the total execution time of 0.069 s . Figure 4.10 shows a snapshot of the output from cProfile, which reveals that acsolve has a "tottime" of 0.030 s . This means that acsolve is spending 0.030 s outside its sub-functions and 0.39 s inside.

| ncalls | tottime | percall | cumtime | percall filename:lineno(function) |
| :---: | :---: | :---: | :---: | :---: |
| 61 | 0.000 | 0.000 | 0.000 | 0.000 < array function internals>:2(can cast) |
| 15 | 0.000 | 0.000 | 0.000 | 0.000 <__array_function__ internals>:2(copyto) |
| 10 | 0.000 | 0.000 | 0.000 | 0.000 <__array_function__ internals>:2(empty_like) |
| 10 | 0.000 | 0.000 | 0.000 | 0.000 <__array_function__ internals>:2(ndim) |
| 10 | 0.000 | 0.000 | 0.000 | 0.000 <frozen importlib._bootstrap>:389(parent) |
| 1 | 0.000 | 0.000 | 0.069 | 0.069 <string>:1(<module>) |
| 332 | 0.001 | 0.000 | 0.001 | 0.000 _init__.py:506(cast) |
| 115 | 0.000 | 0.000 | 0.000 | 0.000 _asarray.py:16(asarray) |
| 332 | 0.000 | 0.000 | 0.000 | 0.000 _internal.py:251(_init__) |
| 332 | 0.000 | 0.000 | 0.001 | 0.000 -internal.py:268(data_as) |
| 332 | 0.000 | 0.000 | 0.001 | 0.000 _internal.py:346(_as_p-parameter_) |
| 10 | 0.000 | 0.000 | 0.000 | 0.000 _methods.py:28(_amax) |
| 10 | 0.000 | 0.000 | 0.000 | 0.000 _methods.py:32(_amin) |
| 40 | 0.000 | 0.000 | 0.000 | 0.000 util.py:129( prune arrav) |
| 1 | 0.030 | 0.030 | 0.069 | 0.069 acsolve.py:17(acsolve) |
| 6 | 0.000 | 0.000 | 0.001 | 0.000 acsolve_wrapper.py:127(mismat) |

Figure 4.10: a snippet of the output from cProfile

Nearly all the heavy calculations are performed by the sub-routines of acsolve, which makes it unrealistic that almost half of the cumulative time is spent outside these functions. Information about the source of this time-consumption is not accessible from cProfile, since it is not caused by function-calls. LineProfiler, on the other hand, gives information about each line in the function. The code is profiled with LineProfiler by adding the declaration "@ profile" before the definition of acsolve and running the command "kernprof $-1-\mathrm{v}$ acsolve.py" in the terminal.

The output from running LineProfiler showed that acsolve spends much time on imposing flat start, setting up the right-hand side of the equation system (2.19), and updating the voltages. The design from the specialization project [1] handles these operations directly in acsolve as pure Python-code. Figure 4.11 is a snippet of the output from LineProfiler, and shows the part of acsolve which updates the voltages. Column number 5 from the left gives the percentage of the total time spent on a particular line. It can be seen that the activity of updating the voltages in total occupies $35.7 \%$ of the cumulative time spent in acsolve. Together, the three operations discussed above account for the time-consumption outside the sub-functions of acsolve. Each of the operations is written as C -extensions to improve the performance of the code further.

| 250 |  |  |  | \# update voltages and angles: |  |
| :--- | ---: | ---: | ---: | :---: | :---: |
| 251 |  |  |  |  |  |
| 252 | 5 | 116.0 | 23.2 | 0.0 | ip $=0$ |
| 253 | 5 | 65.0 | 13.0 | 0.0 | iq $=0$ |
| 254 | 6775 | 74909.0 | 11.1 | 3.8 | for ib in range(0, obj.nbuses): |
| 255 | 6770 | 115395.0 | 17.0 | 5.9 | if(obj.buscod[ib]!=3): |
| 256 | 6765 | 127939.0 | 18.9 | 6.5 | obj.voang[ib] $+=$ correction[ip] |
| 257 | 6765 | 82635.0 | 12.2 | 4.2 | ip $+=1$ |
| 258 | 6765 | 115785.0 | 17.1 | 5.9 | if(obj.buscod[ib] != 2): |
| 259 | 5470 | 114835.0 | 21.0 | 5.9 | obj.vomag[ib] $+=$ correction[obj.nbuses-1+iq] |
| 260 | 5470 | 69266.0 | 12.7 | 3.5 | iq $+=1$ |

Figure 4.11: A snippet of the output from LineProfiler
The impacts of the newest modifications are analyzed by profiling acsolve with cProfile. Figure 4.12 shows a snippet of the output. A comparison of Figure 4.10 and Figure 4.12 show that the time acsolve spends outside its sub-functions has decreased from 0.030 seconds to 0.001 seconds. The total execution time of acsolve is almost cut in half thanks to the most recent modifications.

The optimizations discussed in this section have greatly improved the performance of the program. The previous design of acsolve presented in [1], measured an execution time of 1.732 seconds on the system case 1354pegase. The new design executes the same operations in 0.038

| ncalls | tottime | percall | cumtime | percall filename:lineno(function) |
| :---: | :---: | :---: | :---: | :---: |
| 61 | 0.000 | 0.000 | 0.000 | 0.000 <__array_function__ internals>:2(can_cast) |
| 15 | 0.000 | 0.000 | 0.000 | 0.000 <__array_function__ internals>:2(copyto) |
| 10 | 0.000 | 0.000 | 0.000 | 0.000 <__array_function__ internals>:2(empty_like) |
| 10 | 0.000 | 0.000 | 0.000 | 0.000 <__array function__ internals>:2(ndim) |
| 10 | 0.000 | 0.000 | 0.000 | 0.000 <frozen importlib._bootstrap>:389(parent) |
| 1 | 0.000 | 0.000 | 0.039 | 0.039 <string>:1(<module>) |
| 380 | 0.001 | 0.000 | 0.001 | 0.000 _init_.py:506(cast) |
| 115 | 0.000 | 0.000 | 0.000 | 0.000 _asarray.py:16(asarray) |
| 380 | 0.000 | 0.000 | 0.000 | 0.000 _internal.py:251(__init__) |
| 380 | 0.000 | 0.000 | 0.001 | 0.000 _internal.py:268(data_as) |
| 380 | 0.000 | 0.000 | 0.001 | 0.000 _internal.py:346(_as_parameter_) |
| 10 | 0.000 | 0.000 | 0.000 | 0.000 _methods.py:28(_amax) |
| 10 | 0.000 | 0.000 | 0.000 | 0.000 _methods.py:32(_amin) |
| 40 | 0.000 | 0.000 | 0.000 | 0.000 util.py:129( prune array) |
| 1 | 0.001 | 0.001 | 0.038 | 0.038 acsolve.py:17(acsolve) |
| 6 | 0.000 | 0.000 | 0.001 | 0.000 acsolve_wrapper.py:127(mismat) |

Figure 4.12: A snippet of the output from cprofile
seconds, which is an improvement in the computational performance of approximately $4500 \%$. In other words, the new design is 45 times faster than the original.

### 4.5 Approach to work

The process of updating Topflow has been time-consuming and demanding, yet exciting and educational. In particular, the field of computer programming required a steep learning curve. With previous experiences limited to introduction-courses in MATLAB and C++, the project introduced many new concepts since it utilizes FORTRAN, C, and Python. The fundamental theoretical knowledge and programming skills were acquired during the specialization project "Toolbox for Specialized Power System Analysis" by Åsmund Sælen [1], which leads to this thesis. This section is for that reason divided into two part: subsection 4.5.1 are paragraphs from the specialization-project, while subsection 4.5.2 describes to working-process of this thesis.

### 4.5.1 Theoretical research and skill development

Theoretical knowledge of power system analysis and the skill of programming with different programming languages are crucial when working on the toolbox. The specialization course "ELK-14: Methods and Algorithms for Power Systems" at NTNU provided useful insight into the algorithms used in the toolbox. This thesis uses some of the lecture slides as references, because they gave a compact, thorough description of the required methods. Besides, the mandatory exercises of the course encouraged the use of Python to create the algorithms. "Learning by doing" is arguably the best way to develop skills in a new programming language, which is why "Elk-14" was also very helpful in learning Python.

Free tutorials from online platforms, in particular YouTube and Tutorialspoint, covers all the relevant programming languages. These sites can be recommended to future contributors as a place to get familiar with coding. Other sites for users to ask and answer questions, such as Stack Overflow, are great for fixing bugs in the code, and has been used extensively through this project.

The original Fortran code was provided by professor Olav Bjarte Fosso, who also provided additional documentation of the code. One resource which was especially helpful was a set of videos where he explained parts of the program in detail. These video-clips were made
available on Google Drive. The work done by Statnett's representative, Leif Warland, as well as the code from the previous master thesis by Hege Kvandal [2], was made available in the private repository "topflow" on Github. This repository consisted of a considerable amount of already translated code, both rewritten C-code and wrapper-functions in Python. Both collaborators shared the same goal of modernizing the toolbox, making Python the administrative language which displays the program to the user. These contributors laid the foundation of the work on the toolbox, and established the technique of using Ctypes and Scons to create C-extensions. The scope and the direction of the work on this thesis came naturally with the established decisions and techniques. [1]

### 4.5.2 Master thesis

The specialization course [1] presented a first update of the Newton-Rapshon load-flow function, acsolve. The goal of the project was to make the code run on a 3-bus system. That system did not consist of components such as transformers and shunt-capacitors. Although the test passed, the simplicity of the system called for additional examination of the implementations.

Designing the user interface, which is essential when initializing the data for large power systems, makes up a considerable proportion of the time spent on this thesis. The structure of the input-data was kept close to the data-formats used by MATPOWER [47], because this software is the primary source of the test-cases used to verify the functions in Topflow. The choice of using Excel as the source of the input data is the result of previously experienced that Excel is a flexible and commonly used software. More research could have been conducted in finding alternative options, which may be more efficient and flexible. The same is true for the current method of reading Excel-files, which uses the Python-package "openpyxl". An alternative package is "pandas", which is part of the SciPy-stack (see section subsection 2.1.3) and potentially provides faster functions for communicating with Excel. The intent of this function was, on the other hand, to develop a functional routine of initializing the parameters of large power-systems, something the current method does.

Another time-consuming part of the work was achieving satisfactory results when running load-flows on large systems. Script for running the comparison-test described in subsection 4.3.2 were utilized after the implementation of the functions that initializes the systemparameters. The tests showed that acsolve did not converge for all the test-cases. The approach of detecting the bugs was to print the results of each sub-function, and manually check the values. This method was inefficient, unreliable, and unsuccessful in locating the errors. The errors were resolved thanks to the discovery of automated tests. These tests made validating code more efficient and reliable, since it's a computer that checks the values, not a human. Section 6.2 will discuss the success of the automated tests further.

The toolbox presented in this thesis uses the same technique for writing C-extensions as the previous contributors, and many of the functions are modifications of code originally translated by Leif Warland at Statnett. However, the lack of documentation made it difficult to understand the overall system, which is why the code of this thesis is not merged back with the original branch on GitHub.

## Implementation

This section aims to give future contributors a deeper understanding of the existing implementations of the toolbox.

### 5.1 Installation test

The installation process and the setup.py file are tested in a virtual environment in a different directory than the Topflow-project. Virtual environments have their own site-directories and Python-packages, isolated from the system site directories [48]. This makes them suitable to test if "pip" installes Topflow correctly. Figure 5.1 shows the command that is typed to create the virtual environment called "topflow-env".

C: \Users \aasel \programmering\environments>python -m venv topflow_env
Figure 5.1: Create the virtual environment topflow_env
To enter the virtual environment, it must be activated by using a script in the virtual environment's binary directory. The script is invoked by running the terminal-command "<venv> Scriptlactivate.bat" in the directory of the virtual environment (<venv> is the name of the virtual environment) [48]. Figure 5.2 shows the command that activates the virtual environment topflow_env (which was created in Figure 5.1).

C: \Users \aasel \programmering\environments>topflow_env\Scripts\activate.bat
Figure 5.2: Active/enter the virtual environment topflow_env
The name of the virtual environment will now appear in brackets in the far left end of the terminal. This means that the user is inside the virtual environment. The installed pythonpackages are seen by running "pip-list", as shown in Figure 5.3.

```
(topflow_env) C:\Users\aasel\programmering\environments>pip list
Package Version
pip 19.2.3
setuptools 41.2.0
```

Figure 5.3: The installed packages in the newly created virtual environment "topflow_env"

There are currently only two packages installed in the virtual environment: pip and setuptools. Topflow can now be installed to check if the setup.py installs the package and all its dependencies correctly. This example uses "pip install -e" to install the package in developingmode (the package appears to be installed, but is still editable) [32]. Since Topflow is not an open source project, pip requires the full path of the local Topflow-project. Figure 5.4 shows the command that installs Topflow.
(topflow_env) C:\Users\aasel\programmering\environments>pip install -e c:\users\aasel\programmering\topflow
Figure 5.4: How to install Topflow locally

To validate the installation, "pip list" is once again used to see the installed packages in the virtual environment. Figure 5.5 provides the output of the command, which shows that Topflow and all its dependencies are installed correctly.


Figure 5.5: The result of installing Topflow

### 5.2 The Case-class

At the core of the toolbox, is the class "Case". A Case-object encapsulates all the information needed in load flow studies, and contain various class functions for initializing the data and performing other operations. An initialized Case-object is typically passed to the functions which performs load flows on the system, also referred to as "simulation-functions". Most of the variables of the Case-class are NumPy-arrays, which provides efficient storage and better ways of handling data compared with the inbuilt Python-lists. Another benefit of using NumPyarrays is the mechanism to specify the data type of the content, which is convenient when working with C-extensions. Appendix B. 5 provides the source code of the Case-class as well as Table 2, which lists the class-variables. This section will discuss the Case-class functions.

### 5.2.1 Printing output to the screen

The first way to access the result of a Case-object is by using the print-functions. These are class-functions which provide a quick display of the data directly to the terminal. They are great for checking if the case has been initialized properly, and can be used to print iterationsummaries to the screen when performing simulations. The latter is done by specifying the "print_verbose" argument when running a loadflow, which range from 0 (no output at all) to 3 (very verbose output). The functions are primarily meant for small cases, since the output will
become overwhelming for large systems. Table Table 5.1 show the different functions with a short description of their operations.

| Function | Description |
| :--- | :--- |
| topflow.Case.print_all() | Prints all the data of the Case-instance to the screen |
| topflow.Case.print_buses() | Prints the bus-data of the Case-instance to the screen |
| topflow.Case.print_gens() | Prints the generator-data of the Case-instance to the screen |
| topflow.Case.print_lines () | Prints the line data of the Case-instance to the screen |

Table 5.1: A description of the print-functions

### 5.2.2 Accessing the parameters with "get"

Another way to access the parameters of the Case-class is function "get(component, number, variable)". This function returns the desired parameter based on the input from the user. The three required arguments must be of the type string, and must be known to function.

The different variants of the arguments are listed in table Table 5.2. The argument "number" can either be a single component-number, or a list of numbers, as described in Section 3.3.4. Notice that the generator- and line-numbers are the order the user has placed the components in the Excel-file. All the parameters which can be accessed by the component "gen" can also be accessed by "bus", the only difference is that the user can access them using generator-numbers instead of bus-numbers. The get-function is convenient if the user only wants to access parts of the result, and it enables the parameters to be used in further analysis.

| Component | Number | Variable | Description |
| :---: | :---: | :---: | :---: |
| bus | Any external bus-number(s) | vm va pd qd pg qg | Voltage magnitude of the bus(es) Voltage angle of the bus(es) <br> Active power demand (load) of the bus(es) <br> Reactive power demand of the bus(es) <br> Active power generation of the bus(es) <br> Reactvie power generation of the bus(es) |
| gen | Any generator -number(s) | $\begin{aligned} & \mathrm{vm} \\ & \mathrm{va} \\ & \mathrm{pg} \\ & \mathrm{qg} \end{aligned}$ | Voltage magnitude Voltage angle Active power generation Reactvie power generation |
| line | Any line -number(s) | $\begin{aligned} & \mathrm{pf} \\ & \mathrm{qf} \\ & \mathrm{pt} \\ & \mathrm{qt} \end{aligned}$ | Active power injected at the "from"-end Reactive power injected at the "from"-end Active power injected at the "to"-end Reactive power injected at the "to"-end |

Table 5.2: The parameters in the Generator-data record
The variables of a Case-instance are also accessible with the default Python-approach: "instance.variable". The purpose of the get-function is merely to give the user an intuitive interface.

### 5.2.3 Loading and saving data

This thesis uses Excel as the source of the input and output data. This is mainly because Excel is a widely used and flexible software. The function for loading data from a Excel-file is "topflow.Case.loadxl(filepath)", which requires that the data is structured on the form described in Section 3.3.1. The only argument, "filepath", is the path of the Excel-file which contains the data. If the file is placed in the current working-directory, the name of the file can also be used. Section 3.3.2 showed how to initialize a Case-object by calling the class with the file-path as an input argument. Case has no required input arguments, but when it's called in that way, the loadxl-function will be called automatically to load the data.

Similarly, "topflow.Case.save2xl(filename, save_path = None, save_verbose = 1)" is the function for writing data from a Case-object to a Excel-file. The only required argument is "filename", which is what the user wants to name the file. The second argument, "save_path" tells the function where to save the file. If this argument is not specified, the file will be saved in the current working directory. Lastly, "save_verbose" tells the function how much information it shall save. If save_verbose $=1$, it only saves the final iteration, while if save_verbose $=2$, it creates two separate spreadsheets: one for the final results, and one for all the iterationsummaries.

To this date, using loadxl and save 2 xl is the most efficient way of working on large systems with the current toolbox. However it's encouraged to implement new functions and ways to initialize the Case-objects in the future, to enhance flexibility.

### 5.2.4 External and internal bus-numbers

The program distinguishes between internal and external bus-numbers to provide flexibility to the code. The internal bus-numbers are a result of the order the buses were added to the Caseobject when initializing the data. This means that the external bus-numbers can be chosen arbitrary, as long as they are unique, and the order of which they are placed doesn't matter. This also applies for the slack-bus. Conversion between internal and external bus data are done with the Case-class functions "int2ext" and "ext2int".

### 5.3 The Settings-class

As described in the example from Section 3.3.5, the class "Settings" can be used to specify options before running a simulation. A initialized Settings-object can be passed to the simulationfunctions, which will customize its operations based on this input. Table 5.3 gives a full list over the available options, and specifies the default settings. If a Settings-object is initialized without any input arguments, these settings will apply. The class only deals with settings which are relevant to the function topflow.loadflow(), since other simulation-functions from the original toolbox are not yet translated. Future contributors should update the Settings-class (source-code given in Appendix B.20) along with new translations.

| Name | Variant | Description | Default | Type |
| :---: | :---: | :---: | :---: | :---: |
| version | $\begin{aligned} & \text { 'NR' } \\ & \text { 'FDXX' } \\ & \text { 'FDBX' } \\ & \text { 'FDXB' } \\ & \text { 'FDBB' } \end{aligned}$ | Load-flow algorithms <br> Newton-Rapshon load-flow Fast-Decoupled XX-version Fast-Decoupled BX-version Fast-Decoupled XB-version Fast-Decoupled BB-version | 'NR' | string |
| flat_start | True False | Start with flat start Start with specified voltages | True | bool |
| enf_qlim | True False | Consider reactive generation limits Ignore reactive generation limits | False | bool |
| max_it |  | Maximum number of iterations allowed | 20 | int |
| conv_tol |  | Convergence tolerance | $10^{-6}$ | float |
| pqv | $2$ | FDLF only <br> Start with a $P\|V\|$-iteration <br> Start with a $Q \theta$-iteration | FDXX: 1 <br> FDBX: 1 <br> FDBB: 1 <br> FDXB: 2 | int |
| filename |  | Name of the file containing the results | None | string |
| save_path |  | Path of the file containing the results | cwd* | string |
| save_verbose | $\begin{aligned} & 1 \\ & 2 \end{aligned}$ | Save the final results Save final result + each iteration-summary | 1 | int |
| print_verbose | $\begin{aligned} & 0 \\ & 1 \\ & 2 \\ & 3 \end{aligned}$ | Don't print anything <br> Print the final result <br> Print the final result more verbose <br> Print each iteration-summary | 1 | int |

Table 5.3: Varaibles of the Settings-class.

* $\mathrm{cwd}=$ current working directory


### 5.4 Example cases

"topflow.example_cases()" is a Case-class function that gives access to several test-cases. Topflow has converted these cases from the MATPOWER repository at GitHub [22]. The cases origins either from IEEE test-cases [39], the Pan European Grid Advanced Simulation and State Estimation (PEGASE) [40],[41] or the book "Power System Analysis" by J.J.Grainger and W.D.Stevenson [49]. Table 5.4 lists all the example cases available in Topflow. The digits in the case-names specify the number of buses in the system; hence case 9241 pegase, which represents the size and complexity of the European high voltage transmission network, is the largest system.

| Example case | Origin |
| :--- | :--- |
| case4gs | J.J.Grainger and W.D.Stevenson |
| case14 | IEEE |
| case30 | IEEE |
| case118 | IEEE |
| case300 | IEEE |
| case1354pegase | PEGASE |
| case2869pegase | PEGASE |
| case9241pegase | PEGASE |

Table 5.4: The design purposed in [1]

### 5.5 The loadflow-function

As the name suggests, "topflow.loadflow(..)" is the function for running load-flows. Based on the input from the user, it calls the correct simulation-function, which performs analyzes on a initialized Case-object. The simulation-functions take in the Case-object as an argument, and makes a deep-copy before performing any other operations. This action ensures that the original Case-object is not modified in any way. The deep-copy is then modified by a iteration-process, and returned from the function. A initialized Case-object can therefore be used as a base-case and passed to multiple simulations without the need of reloading the data from the Excel-file. This approach gives a better performance compared to modifying the original Case-object, since loading a Excel-file can be time consuming and should be done as rarely as possible.

This project presents two different simulation-functions: "acsolve" and "decsolve", which are callable through topflow.loadflow(). They are both based on routines in the original toolbox by prof. Olav Bjarte Fosso.

### 5.5.1 Acsolve

The first function callable from topflow.loadflow(), acsolve, performs a Newton-Rapshon loadflow (see section 2.5.7). A initial proposal of this method was presented in the specialization project [1]. That proposal was only tested on a small 3bus-system, which could not provide any information on the speed of the program. The specialization project [1] argued to keep the less complex sub-functions in Python because they would not affect the performance greatly. Section 4.4.5 in this thesis showed the benefits of having more code in C , and as a result, all the sub-functions of acsolve are now C-extensions. Figure 5.6 show the structure of Topflow, highlighting acsolve and all its sub-routines. A flowchart of the algorithm is provided by Figure 5.7.


Figure 5.6: The structure of Topflow, highlighting the branch of acsolve


Figure 5.7: Flow-chart of the Newton-Rapshon load-flow algorithm, acsolve

### 5.5.2 Decsolve

The second simulation-function, decsolve, is based on the fast decoupled load-flow (FDLF) method from section 2.5.8. Figure 5.9 shows the structure of Topflow, highlighting decsolve, and all its subroutines. The original Fortran-code had separate functions for acsolve and decsolve. A comparison of Figure 5.6 and Figure 5.9 shows that the two activities now shares many of the same sub-functions. This is achieved by introducing the argument "pqv". pqv stands for P-Q-version, and tells the functions whether to calculate active parameters ( $\mathrm{pqv}=1$ ), reactive parameters ( $\mathrm{pqv}=2$ ) or both ( $\mathrm{pqv=3} \mathrm{)} .\mathrm{For} \mathrm{example:} \mathrm{the} \mathrm{sub-function} \mathrm{"netinj"} \mathrm{calculates} \mathrm{either}$ only the active power injections, only the reactive power injections, or both, based on the value of pqv.

The flowchart of decsolve is given in Figure 5.10. It shows that after an iteration with pqv $=1$, the algorithm switches to $\mathrm{pqv}=2$ in the next iteration, and vica versa. The user chooses the initial value of pqv by specifying the FDLF-version when calling the function. The built-in versions are: FDXX, FDXB, FDBX and FDBB. The two first letters stands for "Fast-Decoupled" and the two last refers to the version of $B^{\prime}$ and $B^{\prime \prime}$ (see subsection 2.5.8). For example, the version FDBX specifies that $B^{\prime}$ is built with $b_{i j}$, while $B^{\prime \prime}$ is built by substituting $1 / x_{i j}$ for $b_{i j}$. By default all the versions start with a $P \theta$-iteration ( $\mathrm{pqv}=1$ ), except from FDXB which starts with a $Q|V|$-iteration (pqv = 2). The default versions are based on the theory from the paper [25], however, the user can choose to override the default settings by specifying pqv when calling the "loadflow"-function. Figure 5.8 shows an example where the version "FDXX" is run with pqv $=2$ as the initial value. The order of the $P \theta$ - and $Q|V|$-iterations are not expected to show any significant difference, but it makes the code flexible, as the user can customize the FDLF-algorithm for study purposes.

```
>>> import topflow as tf
>>> case1 = tf.example_case('case14')
>>> result1 = tf.loadflow(case1, version = 'FDXX', pqv=2)
Convergence. Number of iterations: 6.5
```

Figure 5.8: A customized FDLF-version


Figure 5.9: The structure of Topflow, highlighting the branch of decsolve


Figure 5.10: Flow-chart of the fast-decoupled load-flow algorithm, decsolve

## Handling generators on the reactive power limits

A big difference between acsolve and decsolve is that while the Jacobian matrix is built in each iteration in the Newton-Rapshon load-flow technique, constant matrices are built only once in the fast decouple load-flow technique. This also affect the way the program handles generators put on the reactive generation limit (also called Var-limited PV-buses). These buses are treated as PQ-buses, as described in section 2.5.9. If a generator is put on the reactive limit, the dimensions of the Jacobian sub-matrix $J 4$ will change. The only operation required by acsolve to handles this is to change the bus-code of the generator. The Jacobian matrix will then change its dimensions automatically in the next iteration. If the VAR-limitd generator is the slack bus, the next availablle PV-bus is chosen as the new reference.

The same approach could be used in decsolve, however, this would mean that the constant matrices would have to be rebuilt before the next iteration. If generators were often put on and off the reactive limits, the advantage of having "constant matrices" would disappear since they would have to be rebuilt often. To avoid this, the matricies $B^{\prime}$ and $B^{\prime \prime}$ are not built in the dimensions of $J 1$ and $J 4$, as the theory from section 2.5 .8 suggests, but in the dimension $N \times N$
( $N$ being the total number of buses). The rows and columns which must be ignored (slack bus for $B^{\prime}$, slack bus and PV -buses for $\left.B^{\prime \prime}\right)$ are disregarded by adding a large number, $M\left(10^{10}\right.$ is used in Topflow) on the diagonal element of the relevant buses. This will in practice result in the contribution of these rows and columns being zero. When a PV-bus is put on the reactive limit, the diagonal element in $B^{\prime \prime}$ of that bus is subtracted by $M$, and the relevant row and column are active. To put a Var-limited PV-bus off the limit, the opposite operation is performed by adding $M z$ on the diagonal. This approach preserves the advantage of the constant matrices, as well as it makes the program more flexible. Some programs require specific ordering of the buses, such as having PQ-buses numbered before PV-buses and the slack bus being the last bus, but this algorithm allows arbitrary numbering of the buses as well as multiple slack-buses.[50]

## Reliability and performance

### 6.1 Result and discussion of the automated tests

The purposed program from the specialization project, [1] was only tested on a 3-bus system, and showed a number of bugs when it was tested on larger systems. Prior to the automated tests, the approach of detecting bugs was to manually check the calculated values of each function by printing the result to the terminal. This approach is neither efficient nor reliable.

Automated test makes detecting and resolving bugs easier, and the test-runner, pytest, was invoked frequently during the development of the code. Figure 6.1 shows an output from running pytest, which resulted in one failure. Observe that pytest prints the unit-test that failed, highlights the false assertion, and reveals the differences between the assertion and the values calculated by the function. The function that failed is called "maxmism", which is a C-extension and sub-function of acsolve. It calculates the active and the reactive power mismatches ( $\Delta P$ and $\Delta Q$ from (2.19) in subsection 2.5.7). A closer look at the code revealed that the function declared a variable as "int" (integer), instead of "double" (decimal number). The wrong declaration caused the C -function to round off numbers, resulting in inaccurate calculations. This bug was particularly hard to detect, since acsolve converged on all the tests-systems, gave satisfactory results on small system, but slightly inaccurate results on larger systems. The fault could, therefore, been a result of inaccurate or missing information from the large system's input-data. This example shows the utility and efficiency of having automated tests.

Figure 6.2 shows the result of running pytest on the final implementations. Three Pythonscripts were identified as tests: the integration test for acsolve (Appendix C.5), the unit tests for acsolve (Appendix C.6), and the unit test for the Case-class (Appendix C.7). The latter contains the test for the function "loadxl", which loads data from Excel-files. The last line of the output shows that all 12 test passed, which means that they are working as expected. The 3 skipped tests are manually marked with a "skip"-decoration because they are functions with "test" in their names that are not automated tests.

The result from Figure 6.2 gives a strong indication that the functions are implemented correctly. Having automated tests increases the reliability of the current implementations, and it's a quick way to check if the functionality of the program is kept when changes are made to the code in the future. However, the test do not guarantee that there are no bugs in the program. Functions may pass automated tests if the assertions are not good enough to check all the situations a function might encounter. Programmers may also make conceptual mistakes, so that the functions work as expected, but the outcome is incorrect. Conceptual mistakes are

```
platform win32 -- Python 3.8.1, pytest-5.4.2, py-1.8.1, pluggy-0.13.1
rootdir: C:\Users\aasel\programmering\topflow
collected 12 items / 1 skipped / 11 selected
tests\test_acsolve_integration.py
tests\test_acsolve_unit.py ...F.....
tests\test_case.py.
```



Figure 6.1: The result of running pytest, which shows an error in the code.
more likely to be picked up by integrated test than unit tests, and it's often challenging to find the source of the error.

```
platform win32 -- Python 3.8.1, pytest-5.4.2, py-1.8.1, pluggy-0.13.1
rootdir: C:\Users\aasel\programmering\topflow
collected 12 items / 3 skipped / 9 selected
tests\test_acsolve_integration.py .
tests\test_acsolve_unit.py
tests\test_case.py -
```

Figure 6.2: The result of the automated tests in the final implementation

The weakest unit tests implemented in the current program are the ones that make use of the data from the three-bus system in Figure 4.3. This is a small system, which lacks complexity and components such as transformers and bus-connected shunt elements. The system is used due to the lack of a better alternative to obtain verified information, such as a Jacobian matrix for a given input. Future contributors should conduct a literature-study to search for better assertions for these unit tests. On the other hand, the integration test for acsolve asserts the final result of a 14-bus system as well as the three-bus system. The 14-bus system includes several transformers, and bus-connected capacitors. The success of the integration test indicates that the implemented transformer-model work as expected.

### 6.2 Comparison-test of acsolve

Section 4.3.2 explained how the test in Appendix C. 2 analyzes the reliability of acsolve further. Table 6.1 summarizes the results of the tests, by displaying how much the voltages calculated by Topflow deviates from pandapower and pypower respectively. All the simulations are performed by running a Newton-Rapshon load-flow with flat-start and a convergence criteria of $10^{-8}$. A precision of 6 decimals is used for both the voltage magnitudes and angles.

| Case | Software | Worst $\|V\|$ deviation [pu] | Worst $\theta$ deviation[degree] |
| :--- | :--- | :--- | :--- |
| case14 | pypower | 0.0 | 0.0 |
|  | pandapower | 0.0 | 0.0 |
| case30 | pypower | 0.0 | 0.0 |
|  | pandapower | 0.0 | 0.0 |
| case118 | pypower | 0.0 | 0.0 |
|  | pandapower | 0.000041 | 0.002694 |
| case1354pegase | pytpower | 0.0 | 0.0 |
|  | pandapower | 0.0 | 0.0 |
| case2869pegase | pypower | 0.0 | 0.0 |
|  | pandapower | 0.0 | 0.0 |
| case9241pegase | pypower | 0.0 | 0.0 |
|  | pandapower | 0.0 | 0.0 |

Table 6.1: The results calculated by Topflow compared with pandapower and pypower when running a Newton-Rapshon load-flow

It can be observed from Table 6.1 that the IEEE-cases show no difference between Topflow and pypower with a precision of 6 decimals. Pandapower gives similar results, except from case 118. The voltages calculated by pandapower differs from both Topflow and pypower on that system, with a worst deviation of 0.000041 pu for the magnitudes, and $0.002694^{\circ}$ for the angles. It's unlikely that this deviation is caused by implementation-errors in any of the packages, since they both pypower and pandapower include validation-models which compares the packages against commercial software [51]. A more likely explanation is that pypower and pandapower uses different input data for the case118. Pandapower has converted the case from pypower, so the source of error may be this transition. This would explain why Topflow and pypower show the same result, since they uses almost identical date-formats (influenced by MATPOWER).

The tests conducted on the PEGASE-cases show no differences among Topflow, pandapower and pypower. These results are of significant importance for the validation of the implemented transformer-model in Topflow. The IEEE-cases use only ideal transformers, which means that the shift-angles are set to 0 . The PEGASE-cases on the other hand, are more complex, and introduces phase shift transformers. These cases also includes values different from 0 for the active components of bus-connected shunt-elements (GL). The fact that Topflow finds the same voltages with a precision of more than $10^{-6}$ show that the program can handle these special cases.

### 6.3 Comparison of decsolve

The primary focus of this thesis has been on analyzing acsolve. However, the last milestone (see Section 1.2) is to translate the second load-flow function, decsolve, from the original FORTRAN-code. Section 5.5 .2 showed how this function is implemented, and that it uses many of the same sub-functions as acsolve. Therefore, most parts of decsolve are already verified by the tests performed on acsolve.

An additional test (see Appendx C.3), compares the voltages calculated by decsolve and acsolve on the same cases that the test from section 6.2 used. The deviation between the two algorithms should be small, since the use the same function to calculate the power injections, and the same condition for terminating the iterative process. The test uses flat-start and a convergencecriteria of $10^{-8}$ for each simulation. The worst voltage-deviations are calculated with a precision of $10^{-6}$ (for both the magnitudes and the angles).

Table 6.2 summarizes the results of the test, and shows that all the three FDLF-versions (which are given in Appendix A) calculates satisfactory results for the IEEE-cases. However, all the simulations performed on the PEGASE-cases diverge. Pandapower has implemented the Primal- and the Dual-method, and a quick simulation in Python showed that these methods converged on the PEGASE-cases listed in Table 6.2. These results show that there are current bugs in the implementation of decsolve. The error is most certain located in one of the distinct functions of decsolve, since the sub-functions of acsolve show reliable results in both the automated tests and the comparison-test from section 6.2.

| Case | FDLF-version | Succsess- <br> flag | Worst $\|V\|$ <br> deviation [pu] | Worst $\theta$ <br> deviation[deg] |
| :--- | :--- | :--- | :--- | :--- |
| case14 | Standard (XX) | Convergence | 0.0 | 0.0 |
|  | Primal (BX) | Convergence | 0.0 | 0.0 |
|  | Dual (XB) | Convergence | 0.0 | 0.0 |
| case30 | Standard(XX) | Convergence | 0.0 | 0.0 |
|  | Primal(BX) | Convergence | 0.0 | 0.0 |
|  | Dual (XB) | Convergence | 0.0 | 0.0 |
| case118 | Standard (XX) | Convergence | 0.0 | 0.0 |
|  | Primal (BX) | Convergence | 0.0 | 0.0 |
|  | Dual (XB) | Convergence | 0.0 | 0.0 |
| case1354pegase | Standard (XX) | Divergence | - | - |
|  | Primal (BX) | Divergence | - | - |
|  | Dual (XB) | Divergence | - | - |
| case2869pegase | Standard (XX) | Divergence | - | - |
|  | Primal (BX) | Divergence | - | - |
|  | Primal (XB) | Divergence | - | - |
| case9241pegase | Standard (XX) | Divergence | - | - |
|  | Primal (BX) | Divergence | - | - |
|  | Dual (XB) | Divergence | - | - |

Table 6.2: The results calculated by Topflow compared with pandapower and pypower when running a Newton-Rapshon load-flow

Apparent bugs in decsolve are not yet resolved, due to the focus and the limited time-frame
of this thesis. The best approach for future contributors in locating the error(s) is to write unit tests for the distinct sub-functions of decsolve. These tests makes it easier to study the action of each part of the program, and resolve the incorrect implementations.

### 6.4 Conclusion on the reliability of Topflow

Having reliable results of the implemented functions is the most important milestone of this thesis. Writing automated tests is considered a good practice in all stages of a project to ensure that the functions are acting as expected. The fact that all the automated tests passed increases confidence that the implementations are correct. As argued in section 6.1, the automated test do not guarantee that there are no bugs in the program, which is why section 6.2 conducted additional tests. These tests compared simulations performed by Topflow, pypower and pandapower on IEEE- and PEGASE-cases. The results of the simulations compliments the automated tests, and concludes that the Newton-Rapshon load-flow is implemented correctly, and gives reliable results.

Section 6.3 analyzed the reliability of the second load-flow function, decsolve, by running simulations on the cases from section 6.2, and comparing the calculated voltages with acsolve. The results was satisfactory for the IEEE-cases, but not for the PEGASE-cases. The apparent bug is most likely located in one of the distinct sub-functions of decsolve, which is why future contributors are recommended to write unit tests for all the sub-functions of decsolve.

### 6.5 Performance

The computational performance of Topflow is analyzed by the script in Appendix C.4, which compares the convergence times of Topflow, pypower and pandapower when running the NewtonRapshon load-flow. The simulations are performed on the cases from section 6.2 that showed lower voltage deviations than $10^{-6}$ for both the magnitudes and angles. Case118 is the only system that is not suited to compare the performances of the programs. All the load-flows are performed by using flat-start and a convergence criteria of $10^{-8}$. The tests are performed on a computer with AMD Ryzen 5 processor at 210 GHz . The module "Timeit", which was discussed in subsection 4.4.4, is used to measure the execution-times of the programs.

Figure 6.3 compares the average values of 100 measurements for each system using Topflow, pandapower and pypower. The calculated times are given in milliseconds. It's observed that the three programs show similar results on all the systems. Topflow is faster than pypower and pandapower for all cases, except for case9241pegase, where pandapower show the shortest convergence time. The similarities are a result of the fact that the three programs uses the same C-library, UMFPACK, through SciPy to solve the sparse linear equation systems. The speed of Topflow is high due to multiple C-extensions, while pandapower uses the just-in-time (jit) compiler numba [52] with success for larger systems [37]. Pypower provides neither, and is therefore the slowest of the three programs.

The jit-compiler used by pandapower translates pure Python-code to machine-code at runtime to give the program speed similar to C and FORTRAN. This compiler will on the other hand cause a delay the first time a function is called, because it has to compile the code [52]. This delay is not included in the calculated times in Figure 6.3, since it was made sure that timeit.timeit() included an initial Newton-Rapshon load-flow in the setup for pandapower, to invoke the jit-compiler before the actual measurements. The script in Appendix C. 1 measures


Figure 6.3: Calculation times for the Newton-Rapshon load-flow on Standard MATPOWER cases
the delay by subtracting the first runtime of the Newton-Rapshon function of pandapower by the mean of the next 100 runtimes. This operation is done to ensure that the execution-time of the function itself do not add to the delay. The script was invoked 10 times to force the jit-compiler to translate the function each time. The delay measured 2.75 seconds on average, regardless of the size of the system.

The delay is obviously worth it when the user wants to call the same pandapower-function multiple times, especially on large systems. Figure 6.3 shows that the jit-compiler has less effect for smaller systems, also, the delay makes work which involves few function calls more time consuming, not less. Regardless, the test shows that the jit-compiler is a powerful software that shrinks the gap in performance between Python and C.

One of the advantage of Topflow is that it do not need the jit-compiler to enhance its performance, thanks to the implemented C-extensions. This makes the program efficient on small systems and low-demanding work, as well larger systems and more demanding work.

## Conclusion

This master thesis has conducted work on updating Topflow: a toolbox for power system analysis. NTNU professor Olav Bjarte Fosso wrote the original code in FORTRAN, but previous projects have worked towards the end goal of making the toolbox available in Python. These projects have used C-extensions to preserve the speed of the FORTRAN-program. The contribution of this thesis builds on an initial update of the Newton-Rapshon load-flow function from the specialization project "Toolbox for Specialized Power System Analysis" by Åsmund Sælen [1]. Newly added features to the toolbox includes functions which read and initializes data from Excel-files. These functions make Topflow able to run simulations on large power-systems.

Automated tests are now implemented to ensure the reliability of the code. The result of running a Newton-Rapshon load-flow with Topflow has been compared with corresponding simulations by other programs to complement the automated tests. The largest test-system is case 9241 pegase, which represents the size and complexity of the European high voltage transmission network [53]. The tests confirm that the implementations are working as expected, and show that Topflow gives reliable results on large systems.

Various profiling-tools have been used to optimize the Newton-Rapshon load-flow function, and justify the design of the Python-interface with C. The most significant optimization is the use of the new solver, scipy.sparse.linalg.spsolve(), which handles sparse matrices. The updated version of Topflow is 45 times faster than the initial version from the specialization project [1], and matches the speed of the open-source projects pandapower and pypower.

The fast-decoupled load-flow routine "decsolve" is also implemented in the updated toolbox presented in this thesis. This implementation is not as well-documented as the Newton-Rapshon load-flow function, but it shares many of the same sub-routines, which have been verified by automated tests.

This thesis's conclusions complete the milestones, and the primary goal set at the start of the working process. Future work should complete the documentation of decsolve, by creating automated tests for the remaining sub-functions, and comparing the result of the load-flow with the Newton-Rapshon activity. The natural next step in translating the remaining parts of the original toolbox is to update the routine for continuation power-flow and contingency analysis. These routines are extensions of the load-flow activities implemented in this thesis, and will introduce new functionality to the program. Hopefully, future contributors can use the documentation presented by this thesis to complete the toolbox, and add additional features to handle the challenges in the future's power systems.

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## Appendix

Appendix A. contains additional information on the Fast-decoupled load-flow form section subsection 2.5.8. Appendix B. provides the source-code of Topflow. Existing parts of the package that are not utilized in this master's thesis (such as the part Hege Bruvik Kvandal developed concerning DC optimal load-flow [2]) are not included. Appendix C. consists of the tests discussed in chapter 6 , which analyses the reliability and performance of the program. The tests uses Excel-files that are too large to include in the Appendix, however, the data is accessible in the GitHub repository [22] where the cases were converted from. Finally, Appendix D. contains the setup.py-file, which is not a part of the Topflow source-code, but located at the root of the repository to support local installations.

## A Fast-decoupled load-flow versions

## Approximations

1. $\left|V_{i}\right|=1$
2. $\sin \left(\theta_{i j}\right)=0$
3. $\cos \left(\theta_{i j}\right)=1$
4. The effect of PV-buses and shunts ar neglected when forming $H_{e q}$ (XB- and XX-version only)
5. $L_{e q}$ is the submatrix $L$ calculated with $b_{i j}$ substituted by $1 / x_{i j}$
6. $H_{e q}$ is the submatrix $H$ calculated with $b_{i j}$ substituted by $1 / x_{i j}$
(Note: if the system is radial or have constant $r / x$ ratios, 6. and 7. are not approximations, but exactly true.)

## Standard Algorithm (XX-version)

1. Initialize iteration count: $k=0$
2. Calculated the active power mismatches, and compute angle corrections:
$\Delta \theta^{(k)}=\frac{H_{e q}^{-1} \Delta P^{(k)}\left(|V|^{(k)}, \theta^{(k)}\right)}{|V|^{(k)}}$
3. Update the angles: $\theta^{(k+1)}=\theta^{(k)}+\Delta \theta^{(k)}$
4. Calculate the reactive power mismatches, and compute magnitude corrections:
$\Delta|V|^{(k)}=\frac{L_{e q}^{-1} \Delta Q^{(k)}\left(|V|^{(k)}, \theta^{(k+1)}\right)}{|V|^{(k+1)}}$
5. Update the magnitudes: $|V|^{(k+1)}=|V|^{(k)}+\Delta|V|^{(k)}$
6. $k=k+1$, go to step 2 .

## Primal Algorithm (BX-version)

1. Initialize iteration count: $k=0$
2. Calculated the active power mismatches, and compute angle corrections:

$$
\Delta \theta^{(k)}=\frac{\left.H^{-1} \Delta P^{(k)}|V|^{(k)}, \theta^{(k)}\right)}{|V|^{(k)}}
$$

3. Update the angles: $\theta^{(k+1)}=\theta^{(k)}+\Delta \theta^{(k)}$
4. Calculate the reactive power mismatches, and compute magnitude corrections: $\Delta|V|^{(k)}=\frac{L_{e q}^{-1} \Delta Q^{(k)}\left(|V|^{(k)}, \theta^{(k+1)}\right)}{|V|^{(k+1)}}$
5. Update the magnitudes: $|V|^{(k+1)}=|V|^{(k)}+\Delta|V|^{(k)}$
6. $k=k+1$, go to step 2 .

## Dual Algorithm (XB-version)

1. Initialize iteration count: $k=0$
2. Calculate the reactive power mismatches, and compute magnitude corrections:
$\Delta|V|^{(k)}=\frac{L^{-1} \Delta Q^{(k)}\left(|V|^{(k)}, \theta^{(k)}\right)}{|V|^{(k)}}$
3. Update the magnitudes: $|V|^{(k+1)}=|V|^{(k)}+\Delta|V|^{(k)}$
4. Calculated the active power mismatches, and compute angle corrections: $\Delta \theta^{(k)}=\frac{H_{e q}^{-1} \Delta P^{(k)}\left(|V|^{(k+1)}, \theta^{(k)}\right)}{|V|^{(k+1)}}$
5. Update the angles: $\theta^{(k+1)}=\theta^{(k)}+\Delta \theta^{(k)}$
6. $k=k+1$, go to step 2 .

## B Topflow

This section contains the source-code of Topflow. Appendix sorts the files alphabetically, and Table 1 gives an overview with a short description of the function(s)/class(es) within.

| File | Description of the function(s)/class(es) |
| :--- | :--- |
| __init__.py | Defines that Topflow is a package |
| acsolve_wrapper.py | Wrapper-functions for the C-extensions of the function "acsolve" |
| acsolve.py | Performs the Newton-Rapshon load-flow |
| bmatrix.py | Builds the constant sub-matrices B’ and B" for the FDLF-method |
| case.py | Stores the system-parameters, communicates with Excel |
| coo_conv.c | Restructures the sparse matrix built by "bujac" to coo-format. |
| decsolve_wrapper.py | Wrapper-functions for the C-extensions of the function "decsolve" |
| decsolve.py | Performs the fast-decoupled load-flow (FDLF) |
| enforce_qlim.py | Enforces reactive power limits on generators |
| flatstart.py | Imposes flat start |
| jacobi.c | Constructs the Jacobian matrix |
| jacobi.h | Headers for jacobi.c |
| loadflow.py | Administers the loadflow-functions "acsolve" and "decsolve". |
| maxmism.py | Finds the power-mismatches ( $\Delta P$ and $\Delta Q)$ |
| netinj.c | Calculates the net power-injections ( $P$ and $Q$ |
| sConstruct | Recuired by SCons to build the DLL that exports the C-extensions |
| select_ver.c | Selects the versions of B’ and B" in the FDLF |
| set_rhs.c | Sets the righ-hand-side of the equation system (2.19) |
| Settings | Specifies settings (load-flow-version, print- and save-options, etc.) |
| topflow.h | Defines some basic identities |
| update_voltages.py | Updates voltage magnitudes, $(\|V\|)$ and voltage angles $(\theta)$ |
| zerosp.c | Handles zero impedance lines |

Table 1: The files in Toplow's source-code

## B. 1 __init_.py

```
1 #This file defines Topflow as a python-package
2 #It also makes the functions in case.py and run.py avail-
# a able for import directly by "from topflow import somefunction"
4 from.case import Case
5 from .loadflow import *
6 from .example_cases.example_cases import example_case
from .example_cases.example_cases import example_case_list
8 from .settings import *
```


## B. 2 acsolve.py

This file contains the function acsolve, which performs the Newton-Rapshon load-flow.

```
import sys
```

2
3 sys.path.append(".")
4 import numpy as np
5 import copy
6 import time
7 import scipy
8 import topflow.acsolve_wrapper as ac
9 from scipy.sparse import csr_matrix
10 from scipy.sparse import isspmatrix_csr
11 from topflow.jacpy import *
12
13
14 \#acsolve performs a Newton-Rapshon load-flow
15
16 def acsolve(
17 case,
18
opt
19 ) :
20 \#Start the timer
21
22

```
    start_clock = time.time()
```

    \#Make a deep copy to prevent changes of the base-case
    obj = copy.deepcopy(case)
    obj.convergence = False
    \#set some required values
    obj.icount \(=0\)
    melem \(=5\) * (obj.nbuses + 2 * obj.nlines)
    \(\mathrm{pqv}=3\)
    \#The following values are set to 0 , because they are used in CPF,
    not NR:
    pdelta \(=0.0\)
    qdelta \(=0.0\)
    alfa \(=n p . z e r o s(o b j . n b u s e s, ~ d t y p e=n p . d o u b l e)\)
    beta \(=n p . z e r o s(o b j . n b u s e s, d t y p e=n p . d o u b l e)\)
    ac.flatstart(
        obj.nbuses,
        obj.buscod,
        obj.vomag,
        obj.voang,
        opt.flat_start
        )
    while True:
    ```
    # Caclulate net injection at all buses
```

    # Caclulate net injection at all buses
    ac.netinj(
    ac.netinj(
        pqv,
        pqv,
        obj.nbuses,
        obj.nbuses,
        obj.nlines,
        obj.nlines,
        obj.pinj,
        obj.pinj,
        obj.qinj,
        obj.qinj,
        obj.vomag,
        obj.vomag,
        obj.voang,
        obj.voang,
        obj.gii,
        obj.gii,
        obj.bii,
        obj.bii,
        obj.gij,
        obj.gij,
        obj.bij,
        obj.bij,
        obj.ratio,
        obj.ratio,
        obj.shift_angle,
        obj.shift_angle,
        obj.ifrom,
        obj.ifrom,
        obj.ito,
        obj.ito,
        obj.ibstat
        obj.ibstat
    )
    )
    # Update reactive gneration and calculate the missmatch
    # Update reactive gneration and calculate the missmatch
    vector
vector
\# pinj and qinj are updated; they are now pdelta and qdelta (
\# pinj and qinj are updated; they are now pdelta and qdelta (
rhs)
rhs)
\# mism[0] contains the bus with worst Pmism, m[1] contains
\# mism[0] contains the bus with worst Pmism, m[1] contains
the bus with worst Qmism
the bus with worst Qmism
ac.mismat(
ac.mismat(
pqv,
pqv,
obj.nbuses,
obj.nbuses,
obj.genbus,
obj.genbus,
obj.buscod,
obj.buscod,
obj.pgen,
obj.pgen,
obj.qgen,
obj.qgen,
obj.pload,
obj.pload,
obj.qload,
obj.qload,
obj.pinj,
obj.pinj,
obj.qinj,
obj.qinj,
alfa,
alfa,
beta,
beta,
pdelta,
pdelta,
qdelta,
qdelta,
)
)
\# Modify for zero impedance lines
\# Modify for zero impedance lines
ac.zerosp(obj.nlines, obj.ifrom, obj.ito, obj.ibstat, obj.
ac.zerosp(obj.nlines, obj.ifrom, obj.ito, obj.ibstat, obj.
xinv, obj.pinj)
xinv, obj.pinj)
ac.zerosp(obj.nlines, obj.ifrom, obj.ito, obj.ibstat, obj.
ac.zerosp(obj.nlines, obj.ifrom, obj.ito, obj.ibstat, obj.
xinv, obj.qinj)

```
    xinv, obj.qinj)
```

```
        # Enforce reactive limitations for generators. If they are on
        the limit, the buscode is changed to -2
        if (obj.icount > 1 and opt.enf_qlim == True):
        #Enforce reactive limits
        obj.mismloc[1] = ac.enforce_qlim(True, obj.nbuses, obj.
mismloc[1], obj.buscod, obj.genbus, obj.numbus, obj.pinj, obj.qinj
, obj.qgen, obj.qmin, obj.qmax, obj.vomag, obj.vgbus)
        #update the number of generators on the reactive limit
        obj.set_nlimgens()
        #Update isa
        #obj.set_isa()
        # Find worst location of mismatches after zero-imp
modification
        ac.maxmism(pqv, obj.nbuses, obj.buscod, obj.mismloc, obj.pinj
, obj.qinj)
    # iteration summary
    if(opt.print_verbose == 3):
        print('\niteration:',obj.icount)
        print('Voltage magnetudes:', obj.vomag)
        print('Voltage angles:', obj.voang)
        print("Worst active power mismatch:", obj.pinj[obj.
mismloc[0]], "at bus:", obj.mismloc[0])
        print("Worst reactive power mismatch:", obj.qinj[obj.
mismloc[1]], "at bus:", obj.mismloc[1])
    #Write to Excel file
    if(opt.filename is not None and opt.save_verbose == 2):
        #Saves each iteration only if save_verbose is True (its
False by default)
        obj.save2xl(opt.filename, opt.save_path, save_verbose =
True)
    #Terminate?
    if abs(obj.pinj[obj.mismloc[0]]) < opt.conv_tol and abs(obj.
qinj[obj.mismloc[1]]) < opt.conv_tol:
    #Convergence
        obj.convergence = True
        stop_clock = time.time()
        runtime = stop_clock - start_clock
        if(opt.print_verbose != 0):
        print("\nConvergence. \nNumber of iterations:", obj.
icount, 'in', runtime, 'seconds')
        if(opt.print_verbose > 1):
                                    print("Final voltage angle:", obj.voang)
```

```
        print("Final voltag magnitude:", obj.vomag)
        print('pgen', obj.pgen)
        if(opt.filename is not None):
            obj.save2xl(opt.filename, opt.save_path, save_verbose
    = opt.save_verbose)
        return obj
    elif obj.icount > opt.max_it:
        print("\nDivergence")
        break
    # Initializes the jacobian lists with -1
    jacbi = np.zeros((melem), dtype=np.double)
    jcol = -np.ones((melem), dtype=np.int32)
    ipv = -np.ones((melem), dtype=np.int32)
    isa = -np.ones((2*obj.nbuses), dtype = np.int32)
    # Building the sparse representation of the jacobian matrix
    ac.bujac(
        obj.nbuses,
        obj.nlines,
        jacbi,
        obj.vomag,
        obj.voang,
        obj.gii,
        obj.bii,
        obj.gij,
        obj.bij,
        obj.ratio,
        jcol,
        ipv,
        obj.ifrom,
        obj.ito,
        obj.ibstat,
    )
    #Restructure the matrix form bujac to coo-format:
    #Find the sixe of the matrix, to initialize the lists for the
coo-matrix
    jacsize = ac.jacsize(jcol)
    #Initialize the list for the coo-matrix
    row = np.zeros((jacsize), dtype = np.int 32)
    col = np.zeros((jacsize), dtype = np.int32)
    data = np.zeros((jacsize), dtype = np.double)
    #Build the coo-matrix
    ac.coo_conv(
        obj.nbuses,
        obj.ngens,
```

```
        obj.buscod,
        isa,
        ipv,
        row,
        jcol,
        col,
        jacbi,
        data
    )
        #Build the sparse matrix on CSR-format
        sparse = csr_matrix((data, (row,col)), [2*obj.nbuses-obj.
ngens + obj.nlimgens-1, 2*obj.nbuses-obj.ngens + obj.nlimgens-1])
    # Set right hand side in the equations:
    rhs = np.zeros(2*obj.nbuses-obj.ngens + obj.nlimgens - 1)
    ac.set_rhs(
        3,
        obj.nbuses,
        obj.buscod,
        obj.vomag,
        obj.pinj,
        obj.qinj,
        rhs
    )
    # Solve the eqution:
    correction = scipy.sparse.linalg.spsolve(sparse,rhs)
    # update voltages and angles:
    ac.update_voltages(
        3,
        obj.nbuses,
        obj.buscod,
        obj.vomag,
        obj.voang,
        correction
    )
    #obj.savejacobi("LoadFlowResult.xlsx",icount, jacobi)
    obj.icount += 1
```


## B. 3 acsolve_wrapper.py

```
import numpy as np
import numpy.ctypeslib as npct
from pathlib import Path
import ctypes
5 from ctypes.util import find_library
6 \text { from ctypes import c_int}
7rom ctypes import c_double
from ctypes import c_bool
9 from ctypes import POINTER
10 from ctypes import byref
import copy
from scipy.sparse import coo_matrix
import os
```

14
15 libpath $=\operatorname{str}\left(P a t h\left(\ldots f i l e \_\right.\right.$) .parent.absolute()) + '<br>clibrary.dll'
16 \# input type for the samplelib function
17 \# must be a double array, with single dimension that is contiguous
18 ar_1d_double = npct.ndpointer(dtype=np.double, ndim=1, flags="
CONTIGUOUS")
9 ar_1d_int $=$ npct.ndpointer(dtype=np.int32, ndim=1, flags="CONTIGUOUS"
)
20 ar_1d_bool = npct.ndpointer(dtype=np.bool, ndim=1, flags="CONTIGUOUS"
)
21
22
23 clib $=$ ctypes.cdll.LoadLibrary(libpath)
24
25 \#aclib.giibii.restype = None
26 \#aclib.giibii.argtypes = (
27 \# [c_int] * 3 + [ar_1d_double] * 13 + [ar_1d_int] * 4 + [
ar_1d_bool] * 2
28 \#)
29
30 clib.set_isa.restype $=$ None
31 clib.set_isa.argtypes = [c_int] + [ar_1d_int]*2
32 clib.flatstart.restype $=$ None
33 clib.flatstart.argtypes $=\left[c \_i n t\right]+\left[a r \_1 d \_i n t\right]+\left[a r \_1 d \_d o u b l e\right] * 2+$
[c_bool]
34
35 clib.netinj.restype $=$ None
36 clib.netinj.argtypes $=$ (
37 [c_int] * 3 + [ar_1d_double] * 10 + [ar_1d_int] * 2 + [ar_1d_bool
]
38 )
39
40 clib.mismat.restype $=$ None
41 clib.mismat.argtypes $=$ [c_int]*2 + [ar_1d_int] * 2 + [ar_1d_double] *
8 + [c_double] * 2

```
4 2
clib.maxmism.restype = None
clib.maxmism.argtypes = [c_int]*2 + [ar_1d_int] *2 + [ar_1d_double]*2
4 5
46 clib.enforce_qlim.restype = c_int
47 clib.enforce_qlim.argtypes = (
48 [c_bool] + [c_int] * 2 + [ar_1d_int] * 3 + [ar_1d_double] * 7
49 )
50
51 clib.zerosp.restype = None
52 clib.zerosp.argtypes = [c_int] + [ar_1d_int]*2 + [ar_1d_bool] + [
        ar_1d_double]*2
53
54 clib.set_rhs.restype = None
clib.set_rhs.argtypes = [c_int]*2 + [ar_1d_int] + [ar_1d_double]*4
56
57 clib.t_u.restype = None
58 clib.t_u.argtypes = [ar_1d_double] * 2 + [c_double] * 4
59
60 clib.addel.restype = c_int
61 clib.addel.argtypes = (
62 [c_int] * 2 + [c_double] + [c_int] + [ar_1d_double] + [ar_1d_int]
    * 2 + [c_int]
63 )
64
65 clib.bujac.restype = None
66 clib.bujac.argtypes = [c_int] * 2 + [ar_1d_double] * 8 + [ar_1d_int]
        * 4 + [ar_1d_bool]
6 7
68 clib.jacsize.restype = c_int
69 clib.jacsize.argtypes = [ar_1d_int]
70
71 clib.coo_conv.restype = None
72 clib.coo_conv.argtypes = [c_int]*2 + [ar_1d_int]*6 + [ar_1d_double]*2
73
74 clib.update_voltages.resytpe = None
5 clib.update_voltages.argtypes = [c_int]*2 + [ar_1d_int] + [
        ar_1d_double]*3
76
7 7 \text { def set_isa(}
78 nbuses,
79 buscod,
80 isa
81 ) :
82 clib.set_isa(
83 nbuses,
        buscod,
        isa
    )
87
```

```
def flatstart(
    nbuses,
    buscod,
    vomag,
    voang,
    flat start
):
    clib.flatstart(
    nbuses,
    buscod,
    vomag,
    voang,
    flat_start
    )
    def netinj(
    pqv,
    nbuses,
    nlines,
    pinj,
    qinj,
    vomag,
    voang,
    gii,
    bii,
    gij,
    bij,
    ratio,
    shift_angle,
    ifrom,
    ito,
    ibstat
    ):
    clib.netinj(
            pqv,
            nbuses,
            nlines,
            pinj,
            qinj,
            vomag,
            voang,
            gii,
            bii,
                    gij,
                    bij,
                    ratio,
                    shift_angle,
                    ifrom,
                    ito,
                    ibstat
```

```
138
1 3 9
140
141 def mismat(
    pqv,
    nbuses,
    genbus,
    buscod,
    pgen,
    qgen,
    pload,
    qload,
    pinj,
    qinj,
    alfa,
    beta,
    pdelta,
    qdelta,
):
    clib.mismat(
            pqv,
            nbuses,
            genbus,
            buscod,
            pgen,
            qgen,
            pload,
                qload,
                pinj,
                qinj,
                alfa,
                beta,
                pdelta,
                qdelta,
    )
def maxmism(
    pqv,
        nbuses,
        buscod,
        mismloc,
        pinj,
        qinj
    ):
        clib.maxmism(
            pqv,
            nbuses,
            buscod,
            mismloc,
            pinj,
```

```
188
189
1 9 0
191
193 verbose,
194 nbuses,
195 qmism,
196 buscod,
197 genbus,
198 numbus,
199 pinj,
200 qinj,
201
202
203
204
205
206 ):
207
208
209
210
211
212
213
214
215
216
217
218
219
220
221
222
223 def zerosp(
224 nlines,
225 ifrom,
226
227
228
229
230 ) :
232
233
234
235
236
237
```

```
        qinj
```

        qinj
    192 def enforce_qlim(
192 def enforce_qlim(
231 clib.zerosp(
231 clib.zerosp(

```
    )
```

    )
    ,
,
qgen,
qgen,
qmin,
qmin,
qmax,
qmax,
vomag,
vomag,
vgbus,
vgbus,
return clib.enforce_qlim(
return clib.enforce_qlim(
verbose,
verbose,
nbuses,
nbuses,
qmism,
qmism,
buscod,
buscod,
genbus,
genbus,
numbus,
numbus,
pinj,
pinj,
qinj,
qinj,
qgen,
qgen,
qmin,
qmin,
qmax,
qmax,
vomag,
vomag,
vgbus,
vgbus,
)
)
ito,
ito,
ibstat,
ibstat,
xinv,
xinv,
xinj
xinj
nlines,
nlines,
ifrom,
ifrom,
ito,
ito,
ibstat,
ibstat,
xinv,
xinv,
xinj

```
                xinj
```

```
238 )
239
240 def set_rhs(
241 pqv,
242 nbuses,
243 buscod,
244 vomag,
245 pinj,
246 qinj,
247 rhs
248 ):
249 clib.set_rhs(
250 pqv,
251
252
253
254 pinj,
255 qinj,
256 rhs
257
258
259 def t_u(u, t, delta_angle, gij, bij, ratio):
260 clib.t_u(u, t, delta_angle, gij, bij, ratio)
261
262
263
def addel(i, j, elem, idim, aa, jcol, ipv, ip):
    return clib.addel(i, j, elem, idim, aa, jcol, ipv, ip)
265
266
267 def bujac(
268 nbuses,
269 nlines,
270 jacbi,
271 vomag,
272 voang,
273 gii,
274 bii,
275 gij,
276 bij,
277 ratio,
278 jcol,
279 ipv,
280 ifrom,
281 ito,
282 ibstat,
283 ):
284 clib.bujac(
285 nbuses,
286
287
nlines,
jacbi,
```

```
        vomag,
        voang,
        gii,
        bii,
        gij,
        bij,
        ratio,
        jcol,
        ipv,
        ifrom,
        ito,
        ibstat,
    )
def jacsize(jcol):
    return clib.jacsize(jcol)
303
304
305
306 def coo_conv(
307 nbuses,
308 ngens,
309 buscod,
310 isa,
311 ipv,
312 row,
313 jcol,
314 col,
315 jacbi,
316 data
317 ) :
318 clib.coo_conv(
        nbuses,
        ngens,
        buscod,
        isa,
            ipv,
            row,
            jcol,
            col,
            jacbi,
            data
    )
def update_voltages(
    pqv,
    nbuses,
    buscod,
    vomag,
    voang,
    correction
```

338 ):
339 clib.update_voltages(
340 pqv,
341
342 nbuses,

343
344
345 buscod,
vomag,
voang,
correction
346 )

## B. 4 bmatrix.c

```
#include<stdbool.h>
```

\#include<math.h>
_declspec(dllexport) void bmatrix(double *xx, double *value, int *
row, int *col, int *isa, int *ifrom, int *ito, bool *ibstat, int
nbuses, int nlines, int pqv) \{
8 //bver str for B-matrix version: $1=B^{\prime}, 2=B^{\prime \prime}$
9 //Kva diagonaler som skal vere store blir modifisert seinare nr
du veit kva type det skal vere

```
int i,k;
```

int $j=0$;
int $1=n b u s e s ;$
int offset $=0$;
if(pqv == 2) \{
//the bus-positions in $B^{\prime \prime}$ lies in the scond half of isa
offset += nbuses;
\}
for (k=0; k<nlines; k++)
\{
if(ibstat[k]) \{
//Add the values of the diagonal elements
value[ifrom[k]] += xx[k];
value[ito[k]] += xx[k];
//Add the position of the diagonal elements
row[ifrom[k]] = ifrom[k];
col[ifrom[k]] = ifrom[k];
row[ito[k]] = ito[k];
col[ito[k]] = ito[k];
//Set a big number at a diagonal element of a bus which
is not included in the jacobien
//That is slack for J1 or generator buses for J4
if(isa[ifrom[k] + offset] == -1 \&\& value[ifrom[k]] < pow
$(10,10)$ ) \{
value[ifrom[k]] += pow(10,10);
\}
if(isa[ito[k] + offset] == -1 \&\& value[ito[k]] < pow
$(10,10)$ ) \{

61

```
            value[ito[k]] += pow(10,10);
            }
        //Build the lower triangle of the Bmatrix (its
        symmetrical)
            //if(ifrom[k] > ito[k]){
                value[l] = -xx[k];
                row[l] = ifrom[k];
                col[l] = ito[k];
                l++;
            / / }
            //else{
                value[l] = -xx[k];
                row[l] = ito[k];
                col[l] = ifrom[k];
                l++;
            // }
                }
    }
```

2 \}

## B. 5 case.py

| Class variable | Description | Datatype |
| :---: | :---: | :---: |
| sbase | Apparent power base of the system | double |
| Bus data |  |  |
| nbuses | Number of buses in the system | int |
| busname[nbuses] | Name of the buses | numpy.str |
| area | Area postions for each bus | numpy.str |
| zone | Zone postition for each bus | numpy.str |
| numbus[nbuses] | External bus numbers | numpy.int32 |
| buscod[nbuses] | List of the buscodes for all buses; 1-PQ, 2-PV, 3-Slack, -2-PV (VAR limited) | numpy.int32 |
| basekv[nbuses] | Base voltage at each bus | numpy.double |
| gl[nbuses] | Active component of bus-connected shunt element | numpy.double |
| bl[nbuses] | Reactive component of bus-connected shunt element | numpy.double |
| vomag[nbuses] | Voltage magnetudes in pu | numpy.double |
| voang[nbuses] | Voltage angles in radians | numpy.double |
| pload[nbuses] | Active load vector for the base case | numpy.double |
| qload[nbuses] | Reactive load vector for the base case | numpy.double |
| Generator data |  |  |
| nlimgens | Number of Generators on the reactive limit | int |
| ngens | Number of generators in the system | numpy.int32 |
| genbus[nbuses] | Pointers to internal generators | numpy.int32 |
| vgbus[ngens] | Voltage setpoint for the generators | numpy.double |
| pgen[ngens] | Active power for generators | numpy.double |
| qgen[ngens] | Reactive power for generators | numpy.double |
| pmax | Maximum active generation | numpy.double |
| qmax | Maximum reactive generation | numpy.double |
| qmax | Minimum reactive generation | numpy.double |
| Line data |  |  |
| nlines | Number of lines in the system | int |
| ifrom[nlines] | From bus number for all lines | numpy.int32 |
| ito[nlines] | To bus number for all lines | numpy.int32 |
| ibstat[nlines] | Indicates status for all lines (In service: 1) | numpy.bool |
| xinv[nlines] | Inverse of the line reactances | numpy.double |
| gij[nlines] | Real part of line suceptance | numpy.double |
| bij[nlines] | Imaginary part of line suceptance | numpy.double |
| gii[nlines] | Real part of the diagonal element in the Ybus | numpy.double |
| bii[nlines] | Imaginary part the diagonal element in the Ybus | numpy.double |
| tapno[nlines] | Pointers to transformer description (0-not a transformer) | numpy.int32 |

Table 2: Class variables of the Case-class. [] denotes the size of the array

1 import os
2 import openpyxl as opyxl
3 import numpy as np
4 import cmath
5 import math
6 import copy
7
8 \#This file contains the case-class, which is used to store the 9 \#data of power-systems
10
1 class Case:
def init $\qquad$ (
self, filepath = None ) :
\#Case data
self.icount = 0
self.convergence = False
\#Bus data
self.sbase $=500$ \#Default MVA value
self.numbus $=$ np.empty(0, dtype $=n p . i n t 32)$
self.slackbusnr = []
self.busname $=$ np.empty(0, dtype $=$ np.str)
self.buscod = np.empty(0, dtype = np.int32)
self.basekv = np.empty(0, dtype = np.double)
self.gs $=$ np.empty (0, dtype $=n p . d o u b l e)$
self.bs = np.empty(0, dtype = np.double)
self.area $=$ np.empty(0, dtype $=n p . i n t 32)$
self.zone = np.empty(0, dtype = np.int32)
self.nbuses = 0
self.vomag = np.empty(0, dtype = np.double)
self.voang = np.empty(0, dtype = np.double)
self.mismloc $=$ np.zeros(2, dtype $=$ np.int32)
\#Generator data
self.genbus $=$ np.empty(0, dtype $=$ np.int 32 )
self.ngens $=0$
self.nlimgens $=0$
self.genstat $=$ np.empty(0, dtype $=n p . b o o l)$
self.genbase $=$ np.empty(0, dtype $=n p . d o u b l e)$
self.vgbus = np.empty(0, dtype = np.double)
self.pgen $=$ np.empty(0, dtype = np.double)
self.qgen $=n p . e m p t y(0, d t y p e=n p . d o u b l e)$
self.qmax $=$ np.empty(0, dtype = np.double)
self.qmin $=$ np.empty (0, dtype = np.double)
self.pmax $=$ np.empty (0, dtype $=n p . d o u b l e)$
self.pmin $=$ np.empty(0, dtype = np.double)
self.qlim_out $=0.01$
self.qlim_in = 0.005
self.dvlim $=0.002$
\#Branch data
self.nlines $=0$
self.ito $=n p . e m p t y(0, d t y p e=n p . i n t 32)$
self.ifrom $=n p . e m p t y(0, d t y p e=n p . i n t 32)$
self.xinv $=n p . e m p t y(0, d t y p e=n p . d o u b l e)$
self.b $=n p . e m p t y(0, d t y p e=n p . d o u b l e)$
self.gi $=n p . e m p t y(0, d t y p e=n p . d o u b l e)$
self.bi $=n p . e m p t y(0, d t y p e=n p . d o u b l e)$
self.gj $=n p . e m p t y(0, d t y p e=n p . d o u b l e)$
self.bj $=n p . e m p t y(0, d t y p e=n p . d o u b l e)$
self.gii $=$ np.empty (0, dtype $=n p . d o u b l e)$
self.bii $=n p . e m p t y(0, d t y p e=n p . d o u b l e)$
self.gij = np.empty(0, dtype = np.double)
self.bij = np.empty(0, dtype = np.double)
self.ratio = np.empty(0, dtype = np.double)
self.shift_angle = np.empty(0, dtype = np.double)
self.ibstat $=$ np.empty(0, dtype $=n p . b o o l)$
self.tapno $=$ np.empty $(0, d t y p e=n p . i n t 32)$
\#Load data
self.pload $=$ np.empty (0, dtype $=n p . d o u b l e)$
self.qload $=n p . e m p t y(0, d t y p e=n p . d o u b l e)$
self.pdelta $=n p . \operatorname{empty}(0$, dtype $=n p . d o u b l e)$
self.qdelta $=n p . \operatorname{empty}(0, d t y p e=n p . d o u b l e)$
self.alfa $=n p . e m p t y(0, d t y p e=n p . d o u b l e)$
self.beta $=n p . e m p t y(0$, dtype $=n p . d o u b l e)$
if(filepath is not None):
self.loadxl(filepath)
def loadxl(self,filepath):
\#This functions reads data from an Excel-file which has a
specific format,
\#and initilizes the inctance variables.
\#"filepath" can only be the name of the file if it's saved in
the same folder as the toolbox
\#it's therefore recomended to use the full path of the file
when calling loadxl()
\#Open the Excel-file, and get the sheets
wb = opyxl.load_workbook(filepath)
case = wb['Case-identification']
bus = wb['Bus-data']
gen $=$ wb['Generator-data']
line =wb['Line-data']

```
    #Case identification
    istart = 1
    jstart = 1
    #In case the user did not begin the data in cell (1,1);
    stop = False
    if(str(case.cell(row = istart, column = jstart).value) != 'IC
') :
            for r in range(1, 40):
                        for c in range(1,29):
                        if(str(case.cell(row = r, column = c).value) == '
IC'):
                    istart = r
                    jstart = c
                    stop = True
                    break
            if(stop):
                    break
            #Read the data from the case-sheet
    self.sbase = float(case.cell(row = istart+1, column = jstart
+ 1).value)
    #Bus data
            self.nbuses = 0
    istart = 1
    jstart = 1
    #In case the user did not begin the data in cell (1,1);
    stop = False
    if(str(bus.cell(row = istart, column = jstart).value) != 'I')
:
```

```
        for r in range(1, 40):
```

        for r in range(1, 40):
            for c in range(1,29):
            for c in range(1,29):
                if(str(bus.cell(row = r, column = c).value) == 'I
                if(str(bus.cell(row = r, column = c).value) == 'I
            '):
                    istart = r
                    jstart = c
                    stop = True
                    break
            if(stop):
                break
    buscod_col = jstart + 2
    pload_col = jstart + 3
    qload_col = jstart + 4
    gs_col = jstart + 5
    bs_col = jstart + 6
    ```
```

        area_col = jstart + 7
        vomag_col = jstart + 8
        voang_col = jstart + 9
        basekv_col = jstart + 10
        zone_col = jstart + 11
        #Read the data from the bus-sheet
        i = istart + 1
        while (bus.cell(row = i,column = jstart).value is not None):
            self.numbus = np.append(self.numbus,int(bus.cell(row = i,
    column = jstart).value))
if(int(bus.cell(row = i, column = buscod_col).value) ==
3):
self.slackbusnr = self.nbuses
self.buscod = np.append(self.buscod, int(bus.cell(row = i
, column = buscod_col).value))
self.pload = np.append(self.pload, float(bus.cell(row = i
, column = pload_col ).value)/self.sbase)
self.qload = np.append( self.qload, float(bus.cell(row =
i, column = qload_col).value)/self.sbase)
self.gs = np.append(self.gs, float(bus.cell(row = i,
column = gs_col).value)/self.sbase)
self.bs = np.append(self.bs, float(bus.cell(row = i,
column = bs_col).value)/self.sbase)
self.area = np.append(self.area,int(bus.cell(row = i,
column = area_col).value))
self.vomag = np.append(self.vomag, float(bus.cell(row = i
, column = vomag_col).value))
self.voang = np.append(self.voang, float(bus.cell(row = i
, column = voang_col).value)/180*math.pi)
self.basekv = np.append(self.basekv,float(bus.cell(row =
i, column = basekv_col).value))
self.zone = np.append(self.zone,int(bus.cell(row = i,
column = zone_col).value))
self.nbuses += 1
i+=1
self.pinj = np.zeros(self.nbuses, dtype=np.double)
self.qinj = np.zeros(self.nbuses, dtype=np.double)
self.set_isa()
\#Generator data
istart = 1
jstart = 1
self.genbus = -np.ones(self.nbuses,dtype = np.int32)
\#In case the user did not begin the data in cell (1,1);

```
```

    stop = False
    if(str(gen.cell(row = istart + 1, column = jstart).value) !=
    'I'):
for r in range(1, 40):
for c in range(1,29):
if(str(gen.cell(row = r, column = c).value) == 'I
'):
istart = r
jstart = c
stop = True
break
if(stop):
break
pgen_col = jstart + 1
qgen_col = jstart + 2
qmax_col = jstart + 3
qmin_col = jstart + 4
vgbus_col = jstart + 5
genbase_col = jstart + 6
genstat_col = jstart + 7
pmax_col = jstart + 8
pmin_col = jstart + 9
\#Read the data from the generato-sheet
i = istart + 1
while (gen.cell(row = i, column = jstart).value is not None):
mbase = float(gen.cell(i,genbase_col).value)
self.bs[self.ext2int(int(gen.cell(row = i, column =
jstart).value))] = self.bs[self.ext2int(int(gen.cell(row = i,
column = jstart).value))]*self.sbase/mbase
self.genbus[self.ext2int(int(gen.cell(row = i, column =
jstart).value))] = i-istart-1
self.pgen = np.append(self.pgen, float(gen.cell(row = i,
column = pgen_col).value)/mbase)
self.qgen = np.append(self.qgen, float(gen.cell(row = i,
column = qgen_col).value)/mbase)
self.qmax = np.append(self.qmax, float(gen.cell( row = i,
column = qmax_col).value)/mbase)
self.qmin = np.append(self.qmin, float(gen.cell( row = i,
column = qmin_col).value)/mbase)
self.vgbus = np.append(self.vgbus, float(gen.cell( row =
i, column = vgbus_col).value))
self.vomag[self.ext2int(int(gen.cell(row = i, column =
jstart).value))] = float(gen.cell( row = i, column = vgbus_col).
value)
self.genbase = np.append(self.genbase, mbase)
self.genstat = np.append(self.genstat, bool(gen.cell(i,
genstat_col).value))
self.pmax = np.append(self.pmax, float(gen.cell(i,

```
```

pmax_col).value)/mbase)
self.pmin = np.append(self.pmin, float(gen.cell(i,
pmin_col).value)/mbase)
i += 1
self.ngens = self.pgen.size
\#Line data
istart = 1
jstart = 1
\#In case the user did not begin the data in cell (1,1);
stop = False
if(str(line.cell(row = istart + 1, column = jstart).value) !=
'I'):
for r in range(1, 40):
for c in range(1,29):
if(str(line.cell(row = r, column = c).value) == '
I'):
istart = r
jstart = c
stop = True
break
if(stop):
break
ito_col = jstart + 1
r_col = jstart + 2
x_col = jstart + 3
b_col = jstart + 4
gi_col = jstart + 8
bi_col = jstart + 9
gj_col = jstart + 10
bj_col = jstart + 11
ratio_col = jstart + 12
shift_col = jstart + 13
ibstat_col = jstart + 14
\#Read the data from the line-sheet
i = istart + 1
while (line.cell(row = i,column = jstart).value is not None):
busi = int(line.cell(row = i, column = jstart).value)
busj = int(line.cell(row = i, column = ito_col).value)
r = float(line.cell(row = i, column = r_col).value)
x = float(line.cell(row = i, column = x_col).value)
ratio = float(line.cell(row = i, column = ratio_col).
value)
self.ifrom = np.append(self.ifrom, self.ext2int(busi))
self.ito = np.append(self.ito, self.ext2int(busj))
self.xinv = np.append(self.xinv, 1/x)

```
```

        self.b = np.append(self.b, float(line.cell(row = i,
    column = b_col).value))
self.gi = np.append(self.gi,float(line.cell(row = i,
column = gi_col).value))
self.bi = np.append(self.bi,float(line.cell(row = i,
column = bi_col).value))
self.gj = np.append(self.gj,float(line.cell(row = i,
column = gj_col).value))
self.bj = np.append(self.bj,float(line.cell(row = i,
column = bj_col).value))
self.ratio = np.append(self.ratio, ratio) if ratio != 0.0
else np.append(self.ratio, 1.0)
self.shift_angle = np.append(self.shift_angle, float(line
.cell(row = i, column = shift_col).value)/180*math.pi)
self.ibstat = np.append(self.ibstat, bool(line.cell(row =
i, column = ibstat_col).value))
self.gij = np.append(self.gij,(1/complex(r,x)).real)
self.bij = np.append(self.bij,(1/complex(r,x)).imag)
self.nlines += 1
i += 1
self.gii = np.zeros(self.nbuses, dtype = np.double)
self.bii = np.zeros(self.nbuses, dtype = np.double)
self.tapno = -np.ones(self.nlines)
self.giibii()
def save2xl(self, filename, save_path = None, save_verbose = 1):
\#This function saves the Case-object to a Excel file (which
it creates)
\#In case the user did not end the file-name with ".xlsx":
save_as = filename if filename[-5:] == '.xlsx' else filename

+ '.xlsx'
if(save_path is not None):
\#Specify the location of the result-file. (The default
location will be the current working dir)
save_as = save_path + '/' + save_as
\#The Final-result-sheet
if(self.convergence):
\#Final-results sheet
try:
\#If the workbook already exist: overwrite it
wb = opyxl.load_workbook(save_as)
try:
\#If a Final-result sheet already exist, reset it

```
```

    ws = wb.get_sheet_by_name('Final-results')
    wb.remove_sheet(ws)
    ws = wb.create_sheet("Final-results")
        except:
            #If it doesent exist, create it
            ws = wb.create_sheet("Final-results")
        if(save_verbose == 1):
            try:
                #Try to delete the iteration sheet from
    previous loadflows
")
wb.remove_sheet(ws2)
except:
pass
except:
\#If the workbook doesen't exist: create a new one
wb = opyxl.Workbook()
ws = wb.active
ws.title = 'Final-results'
ws.append([])
ws.append(["Iterations",self.icount])
ws.append(['Worst mismatch', 'mism[pu]', 'Bus'])
ws.append(['P',self.pinj[self.mismloc[0]],self.mismloc
[0]])
[1]])
ws.append([])
ws.append(['Bus', 'type', 'Vomag[V]', 'Voang[deg]','Pload
', 'Qload', 'Pgen', 'Qgen','Qmax', 'Qmin', 'Violation'])
for i in range(0, self.nbuses):
\#Check if there are any reactive power violations
ig = self.genbus[i]
qmax = None if ig == -1 else self.qmax[ig]
qmin = None if ig == -1 else self.qmin[ig]
violation = None
if(ig>=0):
violation = 1 if self.qgen[ig]>self.qmax[ig] or
self.qgen[ig]<self.qmin[ig] else 0
\#Write generated power to the sheet only if it is a
generator
pg = None
qg = None
if(self.genbus[i] != -1):
pg = self.pgen[self.genbus[i]]
qg = self.qgen[self.genbus[i]]
ws.append([self.numbus[i], self.buscod[i], self.vomag
[i], self.voang[i]*l80/math.pi, self.pload[i], self.qload[i], pg,

```
```

qg, qmax, qmin, violation])
elif(save_verbose == 2):
\#Iteration-summary sheet
try:
\#Check if the workbook already exist
wb = opyxl.load_workbook(save_as)
if(self.icount == 0):
\#Initialize the Iteration sheet
try:
\#If a Iteration sheet already exist,
overwrite it
ws2 = wb.get_sheet_by_name("Iteration-summary
")
wb.remove_sheet(ws2)
ws2 = wb.create_sheet("Iteration-summary")
except:
\#If it doesent exist, create it
ws2 = wb.create_sheet("Iteration-summary")
else:
\#Add data to the sheet in each iteration
ws2 = wb.get_sheet_by_name("Iteration-summary")
except:
\#The workbook doesen't exist: create a new one
wb = opyxl.Workbook()
ws2 = wb.active
ws2.title = "Iteration-summary"
\#Append this to the sheet each iteration:
ws2.append([])
ws2.append(["iteration",self.icount])
ws2.append(['Worst mismatch', 'mism[pu]', 'Bus'])
ws2.append(['P',self.pinj[self.mismloc[0]],self.mismloc
[0]])
ws2.append(['Q',self.qinj[self.mismloc[1]], self.mismloc
[1]])
ws2.append([])
ws2.append(['Bus','type', 'Vomag[V]', 'Voang[deg]', '
Pload', 'Qload', 'Pgen', 'Qgen','Qmax', 'Qmin', 'Violation'])
for i in range(0, self.nbuses):
\#Check if there are any reactive power violations:
ig = self.genbus[i]
qmax = None if ig == -1 else self.qmax[ig]
qmin = None if ig == -1 else self.qmin[ig]
violation = None
if(ig>=0):
violation = 1 if self.qgen[ig]>self.qmax[ig] or
self.qgen[ig]<self.qmin[ig] else 0
\#Write generated power to the sheet only if it is a
generator

```
pg = None
qg = None
if(self.genbus[i] != -1):
pg = self.pgen[self.genbus[i]]
\(q g=\) self.qgen[self.genbus[i]]
ws2.append([self.numbus[i], self.buscod[i], self.
vomag[i], self.voang[i]*180/math.pi, self.pload[i], self.qload[i], pg, qg, qmax, qmin, violation])
wb. save (save_as)
if(self.convergence):
\#pop up the excel file
os.startfile(save_as)
\#isa[2*nbuses] is a list which contains the buses position in the jacobian matrix.
\#If an element \(=-1\), the row is not a part of the jacobian matrix
def set_isa(self):
\(i p=0\)
iq = self.nbuses-1
self.isa \(=-n p\).ones(self.nbuses*2, dtype \(=n p . i n t 32\) )
for ib in range(0,self.nbuses):
if(self.buscod[ib] < 3):
self.isa[ib] = ip
ip += 1
if(self.buscod[ib]<2):
self.isa[ib+self.nbuses] = iq iq += 1
def giibii(self):
\#This function is used by loadxl() when it reads the linedata.
\#It calculates system parameters that are needed to
\# impose the impacts of transfromers and shunt-elements.
\#Consider writing it as a C-extension
self.gii \(=\) copy.deepcopy (self.gs)
self.bii \(=\) copy.deepcopy(self.bs)
for il in range(0,self.nlines):
ifr \(=\) self.ifrom[il]
itr \(=\) self.ito[il]
if(self.ibstat[il] != 0):
self.gii[ifr] += self.gi[il]
self.bii[ifr] += self.bi[il]
```

            self.gii[itr] += self.gj[il]
            self.bii[itr] += self.bj[il]
                tap = 1/self.ratio[il]
                    self.gii[ifr] += self.gij[il]*tap*tap
                            self.bii[ifr] += (self.bij[il] + self.b[il]/2)*tap*
    tap
self.gii[itr] += self.gij[il]
self.bii[itr] += (self.bij[il] + self.b[il]/2)
\#self.xii[ifr] += tap*tap*(self.xinv[il] + self.b[il
]/2)
\#self.xii[itr] += self.xinv[il]*tap + self.b[il]/2
def maxmism(self, pqv):
\#Only used for comparison of speed between python and C.
\#It's replaced by the C-extesnion "maxmism"
\#Finds the largest missmatch and return the positionvector
mismloc
\#mismloc[0]: position of worst P mismatch in pinj
\#mismloc[1]: position of worst Q mismatch in qinj
ptemp = 0
qtemp = 0
for ib in range(0, self.nbuses):
if(self.buscod[ib] != 3 and abs(self.pinj[ib]) > ptemp
and pqv != 2):
self.mismloc[0] = ib
ptemp = abs(self.pinj[self.mismloc[0]])
if(self.buscod[ib] < 2 and abs(self.qinj[ib]) > qtemp and
pqv != 1):
self.mismloc[1] = ib
qtemp = abs(self.qinj[self.mismloc[1]])
def set_nlimgens(self):
self.nlimgens = 0
for k in self.buscod:
if k == -2: self.nlimgens +=1
def get(self, component, number, variable):
\#the get function gives the user access the the parameters of
a Case-instance
\#To handle both a singe component, and a list of componenets
number = np.array(number)
output = np.zeros(0)
if(str.lower(component) == 'bus'):

```
\#Gets access to data at the buses (loads and generators included)
\#Check if the busnumber is correct
if(number.size == 1):
return self.get_bus(number, variable)
else:
for i in range (0, number.size) :
if(self.get_bus(number[i], variable) == None): return None
output = np.append(output, self.get_bus(number[i
```

], variable))

```
            return output
        elif(str.lower(component) in ['gen', 'generator']):
            \#Get access to generator-parameters
            if(number.size == 1):
                    return self.get_gen(number, variable)
            else:
                for i in range (0, number.size) :
                    if(self.get_gen(number[i], variable) \(==\) None):
                    return None
                    output \(=\) np.append(output, self.get_gen(number[i
], variable))
            return output
    elif(str.lower(component) == 'line'):
            \#Get access to line-parameters
            if(number.size == 1):
            return self.get_line(number, variable)
            else:
            for i in range(0, number.size):
                        if(self.get_line(number[i], variable) == None):
                    return None
                    output \(=\) np.append(output, self.get_line(number[i
], variable))
            return output
    def get_bus(self,busnumber, variable):
    if(busnumber not in ['all', 'ALL']):
            try:
            busnumber \(=\) self.ext2int(busnumber)
            except:
            print('Error:',busnumber, ' is not a external bus-
number' )
            return None
    if(str.lower(variable) == 'vm'):
    output = self.vomag if busnumber in ['all', 'ALL'] else
```

self.vomag[busnumber]
elif(str.lower(variable) == 'va'):
output = self.voang/math.pi*180 if busnumber in ['all','
ALL'] else self.voang[busnumber]/math.pi*180
elif(str.lower(variable) == 'pd'):
output = self.pload if busnumber in ['all', 'ALL'] else
self.pload[busnumber]
elif(str.lower(variable) == 'qd'):
output = self.qload if busnumber in ['all', 'ALL'] else
self.qload[busnumber]
elif(str.lower(variable) == 'pq'):
output = self.pload if busnumber in ['all', 'ALL'] else
self.pload[self.genbus[busnumber]]
elif(str.lower(variable) == 'qg'):
output = self.qload if busnumber in ['all', 'ALL'] else
self.qload[self.genbus[busnumber]]
else:
print('Error: ',busnumber, ' has no variable ', variable)
return None
return output
def get_gen(self, genumber, variable):
if(genumber not in ['all', 'ALL'] and (self.ngens<genumber or
genumber<0) ):
print('Error: ', genumber, 'is not a valid generator-
number')
return None
if(str.lower(variable) == 'vm'):
output = self.vomag if genumber in ['all', 'ALL'] else
self.vomag[genumber]
elif(str.lower(variable) == 'va'):
output = self.voang/math.pi*180 if genumber in ['all', '
ALL'] else self.voang[genumber]/math.pi*180
elif(str.lower(variable) == 'pg'):
output = self.pgen if genumber in ['all', 'ALL'] else
self.pgen[genumber]
elif(str.lower(variable) == 'qg'):
output = self.qgen if genumber in ['all', 'ALL'] else
self.qgen[genumber]
else:
print('Error: ',genumber, ' has no variable ', variable)
return None
return output
def get_line(self, genumber, variable):
\#Gets access to line-flow data
if(self.ngens<genumber or genumber<0):

```
```

        print('Error: ', genumber, 'is not a valid generator-
    number')
return None
if(str.lower(variable) == 'pf'):
return None
elif(str.lower(variable) == 'qf'):
return None
elif( str.lower(variable) == 'pt'):
return None
elif( str.lower(variable) == 'qt'):
return None
else:
print('Error: ', variable, ' is not a valid line-variable
')
return None
def print_buses(self):
print('Bus data:')
print('Busnr:',self.numbus)
print('Busnames:',self.busname)
print('BaseKV:',self.basekv)
print('Buscode:', self.buscod)
print('slack-bus:',self.int2ext(self.slackbusnr))
print('Bus-connected shunt(Active comp), gl:',self.gs)
print('Bus-connected shunt(Reactive comp), bl',self.bs)
print('Area:',self.area)
print('Zones:',self.zone)
print('Voltage magnitudes:',self.vomag)
print('Voltage angles',self.voang)
print('Pinj:',self.pinj)
print('Qinj:',self.qinj)
print('Active load:',self.pload)
print('Reactive load:',self.qload)
def print_gens(self):
print('Generator data:')
print('Genbus:',self.genbus)
print('Ngen:', self.ngens)
print('pgen:',self.pgen)
print('qgen:',self.qgen)
print('qmax:',self.qmax)
print('qmin:',self.qmin)
print('vgbus:',self.vgbus)
def print_lines(self):

```
```

    print('Linedata:')
    print('Ifrom:', self.ifrom)
    print('Ito:', self.ito)
    print('GII:', self.gii)
    print('BII:', self.bii)
    print('GIJ:', self.gij)
    print('BIJ:', self.bij)
    print('Ratio:', self.ratio)
    print('Shift_angle:',self.shift_angle)
    def print_all(self):
\#print all the data
self.print_buses()
print('\n')
self.print_gens()
print('\n')
self.print_lines()
def ext2int(self,busnr):
\#convert busnumber (from external to internal)
return int(np.where(self.numbus == busnr)[0][0])
def int2ext(self,busnr):
\#convert busnumber (from internal to external)
return self.numbus[busnr]

```

\section*{B. 6 coo_conv.c}
```

\#include <stdbool.h>
\#include <math.h>

```
\(\qquad\)
```

declspec(dllexport) void coo_conv(int nbuses, int ngens, int *
buscod, int *isa, int *ipv, int *row, int *jcol, int *col, double
*jacbi, double *data) {
int ib, ip, iq, nz, ir, ic, inext;
ip = 0;
iq = nbuses - 1;
//set isa
for(ib = 0; ib <n.buses; ib++) {
isa[ib] = -1;
isa[ib+nbuses] = -1;
if(abs(buscod[ib] != 3)) {
isa[ib] = ip;
ip++;
}
if(abs(buscod[ib])<2) {
isa[ib+nbuses] = iq;
iq ++;
}
}
//copy the information of jacbi, ipv, and jcol
//over to data, row and col, which will be on coo-format
nz = 0;
for(ir = 0; ir<2*nbuses; ir++) {
if(isa[ir] != -1){
inext = ir;
while(inext != -1){
ic = jcol[inext];
if(isa[ic] != -1 \&\& jacbi[inext] != 0.0){
row[nz] = isa[ir];
col[nz] = isa[ic];
data[nz] = jacbi[inext];
nz ++;
}
inext = ipv[inext];
}
}

```
3
4
5
46
\(47 \quad\}\)
48
49
50 \}

\section*{B. 7 decsolve_wrapper.py}
```

1 import numpy as np
2 from pathlib import Path
3 import numpy.ctypeslib as npct
import ctypes
5 from ctypes.util import find_library
6 from ctypes import c_int
from ctypes import c_double
from ctypes import c_bool
9 from ctypes import c_char_p
10 from ctypes import POINTER
11 from ctypes import byref
12 import copy
13 from scipy.sparse import coo_matrix
14 import os
15
16 libpath = str(Path(__file__).parent.absolute()) + '<br>clibrary.dll'
17
18 \# input type for the samplelib function
19 \# must be a double array, with single dimension that is contiguous
20 ar_1d_double = npct.ndpointer(dtype=np.double, ndim=1, flags="
CONTIGUOUS")
21 ar_1d_int = npct.ndpointer(dtype=np.int 32, ndim=1, flags="CONTIGUOUS"
)
22 ar_1d_bool = npct.ndpointer(dtype=np.bool, ndim=1, flags="CONTIGUOUS"
)
23
24 clib = ctypes.cdll.LoadLibrary(libpath)
25
26 clib.h_matrix.restype = None
27 clib.h_matrix.argtypes = [ar_1d_double] * 3 + [ar_1d_int] * 5 + [
c_int] * 2
28
29 clib.l_matrix.restype = None
30 clib.l_matrix.argtypes = [ar_1d_double] * 3 + [ar_1d_int] * 5 + [
c_int] * 4
31
32 clib.heq_matrix.restype = None
33 clib.heq_matrix.argtypes = [ar_1d_double] * 2 + [ar_1d_int] * 5 + [
c_int] * 3
34
35 clib.leq_matrix.restype = None
36 clib.leq_matrix.argtypes = [ar_1d_double] * 2 + [ar_1d_int] * 5 + [
c_int] * 4
37
38 clib.bmatrix.restype = None
39 clib.bmatrix.argtypes = [ar_1d_double] * 2 + [ar_1d_int] * 6 + [
ar_1d_bool] + [c_int] * 3

```
40
```

clib.selectver.restype = None
clib.selectver.argtypes = [ar_1d_double]*4 + [c_int] + [c_char_p]
'''def netpinj(
nbuses,
nlines,
pinj,
vomag,
voang,
gii,
bii,
gij,
bij,
ratio,
ifrom,
ito,
ibstat
):
clib.netpinj(
nbuses,
nlines,
pinj,
vomag,
voang,
gii,
bii,
gij,
bij,
ratio,
ifrom,
ito,
ibstat
)
def netqinj(
nbuses,
nlines,
qinj,
vomag,
voang,
gii,
bii,
gij,
bij,
ratio,
ifrom,
ito,
ibstat
89 ) :
90 clib.netqinj(

```
43
    74
\begin{tabular}{|c|c|}
\hline 91 & nbuses, \\
\hline 92 & nlines, \\
\hline 93 & qinj, \\
\hline 94 & vomag, \\
\hline 95 & voang, \\
\hline 96 & gii, \\
\hline 97 & bii, \\
\hline 98 & gij, \\
\hline 99 & bij, \\
\hline 100 & ratio, \\
\hline 101 & ifrom, \\
\hline 102 & ito, \\
\hline 103 & ibstat \\
\hline 104 & ) \\
\hline 105 & \\
\hline 106 & \\
\hline 107 def & pmismat ( \\
\hline 108 & nbuses, \\
\hline 109 & genbus, \\
\hline 110 & buscod, \\
\hline 111 & pgen, \\
\hline 112 & pload, \\
\hline 113 & pinj, \\
\hline 114 & beta, \\
\hline 115 & pdelta, \\
\hline 116 ) : & \\
\hline 117 & clib.pmismat( \\
\hline 118 & nbuses, \\
\hline 119 & genbus, \\
\hline 120 & buscod, \\
\hline 121 & pgen, \\
\hline 122 & pload, \\
\hline 123 & pinj, \\
\hline 124 & beta, \\
\hline 125 & pdelta \\
\hline 126 & ) \\
\hline 127 & \\
\hline 128 def & qmismat ( \\
\hline 129 & nbuses, \\
\hline 130 & genbus, \\
\hline 131 & buscod, \\
\hline 132 & qgen, \\
\hline 133 & qload, \\
\hline 134 & qinj, \\
\hline 135 & alfa, \\
\hline 136 & qdelta, \\
\hline 137 ) : & \\
\hline 138 & clib.qmismat( \\
\hline 139 & nbuses, \\
\hline 140 & genbus, \\
\hline
\end{tabular}

\section*{buscod,}
qgen, qload, qinj, alfa, qdelta,
    )
' ' '
def h_matrix(bii, bij, buscod, ifrom, ito, nbuses, nlines):
    melem \(=5 *\) (nbuses \(+2 *\) nlines)
    value \(=n p . z e r o s(m e l e m, d t y p e=n p . d o u b l e)\)
    row \(=n p . z e r o s(m e l e m, d t y p e=n p . i n t 32)\)
    col \(=\) np.zeros(melem, dtype \(=\) np.int 32 )
    clib.h_matrix(bii, bij, value, row, col, buscod, ifrom, ito,
    nbuses, nlines)
    erase \(=\) np.where(value \(==0.0\) )
    value = np.delete(value, erase)
    row \(=\) np. delete (row, erase)
    col \(=n p\).delete (col, erase)
    return coo_matrix((value, (row, col)), [nbuses-1, nbuses-1])
def l_matrix(bii, bij, buscod, ifrom, ito, nbuses, nlines, ngens,
    nglim):
        melem \(=2 *\) (nbuses \(+2 \star\) nlines)
        value \(=n p . z e r o s(m e l e m, ~ d t y p e=n p . d o u b l e)\)
        row \(=\) np.zeros (melem, dtype \(=n p\).int 32 )
        col \(=\) np.zeros(melem, dtype \(=n p . i n t 32\) )
        clib.l_matrix(bii, bij, value, row, col, buscod, ifrom, ito,
        nbuses, nlines, ngens, nglim)
        erase \(=\) np.where (value \(==0.0\) )
        value = np.delete(value, erase)
        row = np.delete(row, erase)
        col \(=n p\).delete (col, erase)
        return coo_matrix((value, (row, col)), [nbuses-ngens + nglim,
        nbuses-ngens + nglim])
    def heq_matrix(xinv, buscod, ifrom, ito, nbuses, nlines, slackbusnr):
    melem \(=2 *\) (nbuses \(+2 *\) nlines)
        value \(=n p . z e r o s(m e l e m, ~ d t y p e=n p . d o u b l e)\)
        row \(=n p . z e r o s(m e l e m, ~ d t y p e=n p . i n t 32)\)
        col \(=n p . z e r o s(m e l e m, d t y p e=n p . i n t 32)\)
        clib.heq_matrix(xinv, value, row, col, buscod, ifrom, ito, nbuses
        , nlines, slackbusnr)
        erase \(=\) np. where (value \(==0.0\) )
        value \(=n p . d e l e t e(v a l u e, ~ e r a s e)\)
    row \(=\) np.delete (row, erase)
    col \(=\) np.delete (col, erase)
    return coo_matrix((value, (row,col)), [nbuses-1, nbuses-1])
def leq_matrix(xinv, buscod, ifrom, ito, nbuses, nlines, ngens, nglim
    ) :
    melem \(=2 *\) (nbuses \(+2 \star\) nlines)
    value \(=n p . z e r o s(m e l e m, ~ d t y p e=n p . d o u b l e)\)
    row \(=n p . z e r o s(m e l e m, ~ d t y p e=n p . i n t 32)\)
    col \(=\) np.zeros (melem, dtype \(=n p . i n t 32\) )
    clib.leq_matrix(xinv, value, row, col, buscod, ifrom, ito, nbuses
    , nlines, ngens, nglim)
    erase \(=\) np.where(value \(==0.0\) )
    value = np.delete(value, erase)
    row \(=\) np.delete (row, erase)
    col \(=\) np.delete (col, erase)
    return coo_matrix((value, (row,col)), [nbuses-ngens + nglim,
    nbuses-ngens + nglim])
def bmatrix(xx, buscod, ifrom, ito, ibstat, nbuses, nlines, pqv):
    melem \(=2 *\) (nbuses \(+2 \star\) nlines)
    value \(=n p . z e r o s(m e l e m, ~ d t y p e=n p . d o u b l e)\)
    row \(=n p . z e r o s(m e l e m, ~ d t y p e=n p . i n t 32)\)
    col \(=\) np.zeros (melem, dtype \(=n p\).int32)
    isa \(=-n p\).ones( \((2 \star\) nbuses \()\), dtype \(=n p\). int 32\()\)
    clib.bmatrix(xx, value, row, col, buscod, isa, ifrom, ito, ibstat
    , nbuses, nlines, pqv)
    '''erase = np.where(value ==0.0)
    value = np.delete(value, erase)
    row \(=\) np.delete (row, erase)
    col \(=n p . d e l e t e(c o l, ~ e r a s e) ' ' '\)
    return coo_matrix((value, (row, col)), [nbuses, nbuses])
def selectver(xp, xq, xinv, bij, nlines, bver):
    clib.selectver(xp, xq, xinv, bij, nlines, bver.encode('utf-8'))

\section*{B. 8 decsolve.py}
```

import sys
2
3 sys.path.append(".")
import numpy as np
import copy
6 from scipy.sparse.linalg import spsolve
7 from scipy.sparse import csr_matrix
8 import topflow.acsolve_wrapper as ac
9 import topflow.decsolve_wrapper as dec
10
11
12 \#decsolve performs a fast-decoupled load-flow
13 \#Versions:
14 \#FDXX (Standard-method): pqv = 1, bver = 'XX'
15 \#FDBX (Primal-method): pqv = 1, bver = 'BX'
16 \#FDXB (Dual-method): pqv = 2, bver = 'XB'
17 \#Custom: the user can choose iteration-sequence(pqv) and the versions
of B' and B''(bver)
18
19 def decsolve(case, opt):
20
21
22
23

```
#make a deep copy to prevent change of the base-case
```

\#make a deep copy to prevent change of the base-case
obj = copy.deepcopy(case)
obj = copy.deepcopy(case)
obj.convergence = False
obj.convergence = False
pdelta = 0.0
pdelta = 0.0
qdelta = 0.0
qdelta = 0.0
xp = np.zeros(obj.nlines)
xp = np.zeros(obj.nlines)
xq = np.zeros(obj.nlines)
xq = np.zeros(obj.nlines)
alfa = np.zeros(obj.nbuses, dtype=np.double)
alfa = np.zeros(obj.nbuses, dtype=np.double)
beta = np.zeros(obj.nbuses, dtype=np.double)
beta = np.zeros(obj.nbuses, dtype=np.double)
genlim = np.zeros(obj.ngens, dtype = np.bool)
genlim = np.zeros(obj.ngens, dtype = np.bool)
if(opt.pqv == 1):
if(opt.pqv == 1):
\#The standard starts with P-calculations
\#The standard starts with P-calculations
pqv = 1
pqv = 1
elif(opt.pqv == 2):
elif(opt.pqv == 2):
\#Dual algorithm starts with Q-calculations
\#Dual algorithm starts with Q-calculations
pqv = 2
pqv = 2
else:
else:
print('Error: pqv must have the value 1 or 2, not', opt.pqv)
print('Error: pqv must have the value 1 or 2, not', opt.pqv)
\#Set flat start (or not)
\#Set flat start (or not)
ac.flatstart(
ac.flatstart(
obj.nbuses,
obj.nbuses,
obj.buscod,
obj.buscod,
obj.vomag,

```
    obj.vomag,
```

```
        obj.voang,
        opt.flat_start
        )
    #Select bm and bmm version (XX, XB, BX, BB)
    dec.selectver(xp, xq, obj.xinv, obj.bij, obj.nlines, opt.bver)
    #Build the constant matrices
    bm = csr_matrix(dec.bmatrix(xp,obj.buscod, obj.ifrom, obj.ito,
obj.ibstat, obj.nbuses, obj.nlines, 1))
    bmm = csr_matrix(dec.bmatrix(xq,obj.buscod, obj.ifrom, obj.ito,
obj.ibstat, obj.nbuses, obj.nlines, 2))
    while True:
    # Caclulate net injection (pinj or qinj) at all buses
    ac.netinj(
            pqv,
            obj.nbuses,
            obj.nlines,
            obj.pinj,
            obj.qinj,
            obj.vomag,
            obj.voang,
            obj.gii,
            obj.bii,
            obj.gij,
            obj.bij,
            obj.ratio,
            obj.shift_angle,
            obj.ifrom,
            obj.ito,
            obj.ibstat
    )
    # Update reactive gneration and calculate the missmatch
vector
    # pinj and qinj are updated; they are now pdelta and qdelta (
rhs)
    # mism[0] contains the bus with worst Pmism, m[1] contains
the bus with worst Qmism
    ac.mismat(
        pqv,
            obj.nbuses,
            obj.genbus,
            obj.buscod,
            obj.pgen,
            obj.qgen,
            obj.pload,
```

```
        obj.qload,
        obj.pinj,
        obj.qinj,
        alfa,
        beta,
        pdelta,
        qdelta
        )
        # Modify for zero impedance lines
        if(pqv == 1):
        ac.zerosp(obj.nlines, obj.ifrom, obj.ito, obj.ibstat, obj
.xinv, obj.pinj)
        elif(pqv == 2):
            ac.zerosp(obj.nlines, obj.ifrom, obj.ito, obj.ibstat, obj
.xinv, obj.qinj)
    # Find worst location of mismatches after zero-imp
modification
        ac.maxmism(pqv, obj.nbuses, obj.buscod, obj.mismloc, obj.pinj
, obj.qinj)
    #Enfroce reactive limitations
        if (obj.icount > 1.5 and opt.enf_qlim == True and pqv == 2):
        obj.mismloc[1] = ac.enforce_qlim(True, obj.nbuses, obj.
mismloc[1], obj.buscod, obj.genbus, obj.numbus, obj.pinj, obj.
qinj, obj.qgen, obj.qmin, obj.qmax, obj.vomag, obj.vgbus)
            #Update bmm
            for ib in range(0, obj.nbuses):
                if(obj.buscod[ib] == -2 and bmm[ib,ib] >= pow(10,10))
:
                    #Set PV bus as Var-limited
                    bmm[ib,ib] -= pow(10,10)
                elif(obj.buscod[ib] == 2 and bmm[ib,ib] < pow(10,10))
:
                    #Set Var-limited PV bus back as PV-bus
                    bmm[i.b,i.b] += pow(10,10)
        # iteration summary
        if(opt.print_verbose == 3):
        print('\niteration:',obj.icount)
        print('Voltage magnetudes:', obj.vomag)
        print('Voltage angles:', obj.voang)
        print("Worst active power mismatch:", obj.pinj[obj.
mismloc[0]], "at bus:", obj.mismloc[0])
        print("Worst reactive power mismatch:", obj.qinj[obj.
mismloc[1]], "at bus:", obj.mismloc[1])
        #Write to Excel file
        if(opt.filename is not None and opt.save_verbose == 2):
```

\#Saves each iteration only if save_verbose is True (its False by default)
obj.save $2 x l\left(f i l e n a m e=o p t . f i l e n a m e, ~ s a v e \_p a t h=o p t\right.$.

```
save_path, save_verbose = 2)
```

    \#Terminate?
    if abs(obj.pinj[obj.mismloc[0]]) < opt.conv_tol and abs(obj.
    qinj[obj.mismloc[1]]) < opt.conv_tol:
\#Convergence
obj.convergence $=$ True
if (opt.print_verbose ! = 0):
print("\nConvergence. Number of iterations:", obj.
icount)
if(opt.print_verbose > 1):
print("Final voltage angle:", obj.voang)
print("Final voltag magnitude:", obj.vomag)
if(opt.filename is not None):
obj.save2xl(opt.filename, opt.save_path, save_verbose
= opt.save_verbose)
return obj
elif obj.icount > opt.max_it*2:
print("\nDivergence")
return 0
\#Right hand side in the equations:
rhs = np.zeros(obj.nbuses)
ac.set_rhs(
pqv,
obj.nbuses,
obj.buscod,
obj.vomag,
obj.pinj,
obj.qinj,
rhs
)
\#Solve the equation
if(pqv == 1):
correction $=$ spsolve(bm, rhs)
elif(pqv == 2):
correction $=$ spsolve(bmm, rhs)
\#Update voltages (magnitude or angles):

```
ac.update_voltages(
    pqv,
    obj.nbuses,
    obj.buscod,
    obj.vomag,
    obj.voang,
    correction
)
if(pqv == 1):
    #The standard ends with Q-calculations
    pqv = 2
elif(pqv == 2):
    #Dual algorithm ends with P-calculations
    pqv = 1
obj.icount += 0.5
```


## B. 9 enforce_qlim.c

\#include <stdio.h>
\#include <stdbool.h>
3
4 \#define $\operatorname{ABS}(x)(((x)<0) \quad ?(-(x)):((x))) / *$ if loop: ABS (X) gives
the absolute value of $x * /$
5
6 //verbose = True; specificaions are printed to the screen.
7
8
__declspec(dllexport) int enforce_qlim(bool verbose, int nbuses, int
qmism,
12 double qmarg;
13 bool violation;
14 int ig;
15
int $s=1$;
for (int $k=0$; $k$ <nbuses; $k++$ )
\{
ig = genbus[k];
if (ig >= 0)
\{
if (buscod[k] >= 2)
\{
/* Enforce violated reactive units */

/* Exceeds maximum ?*/
if (violation)
\{
if (verbose)
\{
printf("Ulim gen @ bus \%5d ",
numbus[k]);
printf("qmax $=\% 5.4 f$, qgen $=\%$
6.4f $\backslash n ", ~ q m a x[i g], ~ q g e n[i g]) ; ~$
//qmax is now the setpoint, find the
missmatch (qinj[k])
int *buscod, int *genbus, int *numbus, double *pinj,
double *qinj, double *qgen, double *qmin, double *
qmax, double *vomag, double *vgbus)
\}
34
35
fixed at maximum
qinj[k] = qmax[ig] - qgen[ig];
//the reactive generation is now
qgen[ig] = qmax[ig];

```
    //choose another slack-bus if
    //the current slack-bus violates the
    limits
    if(buscod[k] == 3){
    pinj[k] = 0.0;
    //chose the next PV-bus as the new
    if(s == nbuses){
        //go to the start of buscod if
            // is the last generator
            s -= nbuses;
        }
        while(buscod[k+s] != 2){
        s++;
    }
    buscod[k+s] = 3;
    s = 1;
}
buscod[k] = -2;
}
else
{
//Below minimum?
violation = (qmin[ig] > qgen[ig]);
if (violation){
    if(verbose)
        {
                printf("Llim gen @ bus %5d ",
    numbus[k]);
                                    printf("qmin = % 5.4f, qgen =
    % 6.4f \n", qmin[ig], qgen[ig]);
        //qmin is now the setpoint, find the
        missmatch(qinj[k])
        qinj[k] = qmin[ig] - qgen[ig];
        //the reactive generation is now
    fixed at minumum
        qgen[ig] = qmin[ig];
        //choose another slack-bus if it
    violates the limit
        if(buscod[k] == 3) {
        pinj[k] = 0.0; //0 mismatch
        //if the slack bus is on the
    limit;
        //choose the next PV-bus as the
```

```
new slack-bus
```

```
                                    if(s >= nbuses){
```

                                    if(s >= nbuses){
                                    s = 1;
                                    s = 1;
                                    }
                                    }
                                    while(buscod[k+s] != 2){
                                    while(buscod[k+s] != 2){
                                    s++;
                                    s++;
                                    }
                                    }
                                    buscod[k+s] = 3;
                                    buscod[k+s] = 3;
                                    s = 1;
                                    s = 1;
                                    }
                                    }
                                    buscod[k] = -2;
                                    buscod[k] = -2;
                                    }
                                    }
                                }
                                }
                                if (violation &&ABS(qinj[k]) > ABS(qinj[qmism])){
                                if (violation &&ABS(qinj[k]) > ABS(qinj[qmism])){
                                    //update the worst mismatch
                                    //update the worst mismatch
                                    qmism = k;
                                    qmism = k;
                    }
                    }
                    }
                    }
                    //Relax nonbinding units
                    //Relax nonbinding units
            else if (buscod[k] == -2 && qmin[ig] != qmax[ig] )
            else if (buscod[k] == -2 && qmin[ig] != qmax[ig] )
                    if (vgbus[ig] == 0.0)
                    if (vgbus[ig] == 0.0)
                            {
                            {
                        vgbus[ig] = vomag[k];
                        vgbus[ig] = vomag[k];
            }
            }
                            if ((qgen[ig] == qmax[ig] && vomag[k]>vgbus[ig])
                            if ((qgen[ig] == qmax[ig] && vomag[k]>vgbus[ig])
                            || (qgen[ig] == qmin[ig] && vomag[k]<vgbus[ig]))
                            || (qgen[ig] == qmin[ig] && vomag[k]<vgbus[ig]))
                        {
                        {
                                    if(verbose)
                                    if(verbose)
                            {
                            {
                                    printf("Relaxed generator constraints
                                    printf("Relaxed generator constraints
    on %d \n", numbus[k]);
on %d \n", numbus[k]);
}
}
vomag[k] = vgbus[ig];
vomag[k] = vgbus[ig];
buscod[k] = 2;
buscod[k] = 2;
qinj[k] = 0;
qinj[k] = 0;
}
}
}
}
}
}
}
}
return qmism;

```
    return qmism;
```

126 \}

## B. 10 flatstart.c

```
#include <stdbool.h>
2 #include <math.h>
3
4 __declspec(dllexport) void flatstart(int nbuses, int *buscod, double
        *vomag, double *voang, bool flat_start) {
5
int ib;
7
if(flat_start == true){
9 for(ib= 0; ib<nbuses; ib++) {
10 if(buscod[ib] != 3){
11 voang[ib] = 0.0;
12 if(buscod[i.b] != 2 && vomag[i.b] != 1.0){
13 vomag[ib] = 1.0;
14 }
15 }
16 }
17 }
19 }
```

18

## B. 11 jacobi.c

```
#include <stdbool.h>
#include <math.h>
5 #include "jacobi.h"
7 void t_u(double *u, double *t, double delta_angle, double gij, double
    bij, double ratio)
{
9 double cosda, sinda;
15 else{
16 cosda = cos(delta_angle);
17 sinda = sin(delta_angle);
18 }
25 int addel(int i, int j, double elem, int idim, double *aa, int *jcol,
    int *ipv, int ip)
```

```
    if (i == j)
```

    if (i == j)
        {
            aa[i] += elem;
                jcol[i] = i;
            }
    else
        while (true)
            if (ipv[i] == -1) //0 used in fortran. -1 here because
        index begins with 0 in C/python
            {
                                ipv[i] = ip;
                                jcol[ip] = j;
                                aa[ip] = elem;
                                ip++;
                                break;
            }
            else
            {
                i = ipv[i];
    ```
3
4
6
10
11
12
13
14
19
20
21
22
23 \}
24
26 \{
27
28
29
30
31
32
33
34
35
36
37
38
39
56 double *jacbi, double *vomag, double *voang, double *gii,
    double *bii, double *gij,
                double *bij, double *ratio,
                        int *jcol, int *ipv, int *ifrom, int *ito, bool *ibstat
                        )
59
60 \{
61
```

if (jcol[i] == j)
{
aa[i] += elem;
break;

```
                            \}
                \}
    return ip;
\}
void bujac(int nbuses, int nlines,
                        int k, i,j, i2, j2, idim, i_p, i_q, ip;
                double volt, elem, delta_angle;
                double uij[1], uji[1], tij[1], tji[1];
    idim \(=2\) *nbuses;
    ip \(=2\) *nbuses;
    for (i=0; i<nbuses; i++)
        \{
            volt \(=2 *\) vomag[i];
            elem = volt*gii[i];
            \(j=i+n b u s e s ;\)
            ip = addel(i, j, elem, idim, jacbi, jcol, ipv, ip);
            elem = -volt*bii[i];
            ip = addel(j, j, elem, idim, jacbi, jcol, ipv, ip);
        \}
    for ( \(k=0\); \(k<n l i n e s ; ~ k++\) )
        \{
            if (ibstat[k])
            \{
                    i \(=\) ifrom[k];
                    \(j=i t o[k] ;\)
            i2 \(=\) i + nbuses;
            j2 \(=\) j + nbuses;
            delta_angle = voang[i] - voang[j];
            volt = vomag[i]*vomag[j];
            t_u(uij, tij, delta_angle, gij[k], bij[k], ratio[k]);
            t_u(uji, tji,-delta_angle, gij[k], bij[k], ratio[k]);
            // DP/DA
        elem = -volt**uij;
        elem = -volt**uji;
        // DP/DV
        // DQ/DA
        elem = volt**tij;
        elem = volt**tji;
        // DQ/DV
    //which are differnet form 0 .
    int size = 0;
    while(true) \{
        if(jcol[size] == -1) \{
            return size;
        ip = addel(i, j, elem, idim, jacbi, jcol, ipv, ip);
        ip \(=\) addel(i, i,-elem, idim, jacbi, jcol, ipv, ip);
        ip = addel(j, i, elem, idim, jacbi, jcol, ipv, ip);
        ip \(=\) addel(j, j,-elem, idim, jacbi, jcol, ipv, ip);
        elem \(=\)-vomag[i]**tij;
        ip = addel(i, j2, elem, idim, jacbi, jcol, ipv, ip);
        elem = -vomag[j]**tji;
        ip = addel(j, i2, elem, idim, jacbi, jcol, ipv, ip);
        elem = -vomag[j]**tij;
        ip = addel(i, i2, elem, idim, jacbi, jcol, ipv, ip);
        elem \(=\)-vomag[i]**tji;
        ip \(=\) addel(j, j2, elem, idim, jacbi, jcol, ipv, ip);
        ip = addel(i2, j, elem, idim, jacbi, jcol, ipv, ip);
        \(i p=a d d e l(i 2, i,-e l e m, ~ i d i m, ~ j a c b i, ~ j c o l, ~ i p v, ~ i p) ; ~\)
        ip = addel(j2, i, elem, idim, jacbi, jcol, ipv, ip);
        ip = addel(j2, j,-elem, idim, jacbi, jcol, ipv, ip);
        elem = -vomag[i]**uij;
        ip = addel(i2, j2, elem, idim, jacbi, jcol, ipv, ip);
        elem = -vomag[j]**uij;
        ip = addel(i2, i2, elem, idim, jacbi, jcol, ipv, ip);
        elem = -vomag[j]**uji;
        ip = addel(j2, i2, elem, idim, jacbi, jcol, ipv, ip);
        elem = -vomag[i]**uji;
        ip \(=\) addel (j2, j2, elem, idim, jacbi, jcol, ipv, ip);
    //This routine finds the number of element in the
    //sparse jacobian matrix bult by bujac,
```

145
}
size++;
147
}
return 0;
149 }

```

\section*{B. 12 jacobi.h}
```

1 \#ifndef HEADER_jacobi
2 \#define HEADER_jacobi
3
4 __declspec(dllexport) void t_u(double *u, double *t, double
delta_angle, double gij, double bij, double ratio);
5
declspec
double *aa, int *jcol, int *ipv, int ip);
6 __declspec(dllexport) void bujac(int nbuses, int nlines,
7 double *jacbi, double *vomag, double *voang, double *gii,
double *bii, double *gij,
8 double *bij, double *ratio,
9 int *jcol, int *ipv, int *ifrom, int *ito, bool *ibstat
10 );
1 1
12 __declspec(dllexport) int jacsize(int *jcol);
13 \#endif /* HEADER_jacobi */

```

\section*{B. 13 loadlflow.py}

1
2 from topflow.case import Case
3 from topflow.acsolve import acsolve
4 from topflow.settings import Settings
5 from topflow.decsolve import decsolve
6 import copy
7
8 \#This is the administrative function of the load-flow algorithms acsolve and decsolve

9
10 def loadflow(Case, mysettings = Settings(), version = None, flat_start = None, enf_qlim = None, max_it = None, conv_tol = None , pqv = None, filename = None, save_path = None, save_verbose = None, print_verbose \(=\) None):
11
12
13
14
15
```

        args = locals()
    ```
        myset_copy \(=\) copy.deepcopy (mysettings)
        for i in args:
            if( i not in ['Case', 'mysettings', 'args'] and args[i] is
        not None):
            myset_copy.set_var(i, args[i])
17
18
        if(myset_copy.version \(==\) 'NR'):
            result \(=\) acsolve (Case, myset_copy)
        elif(myset_copy.version in ['FDXB', 'FDBX', 'FDXX', 'FDBB']):
            result \(=\) decsolve(Case, myset_copy)
        else:
            print (myset_copy.version,' is not a valid load flow version')
            return None
        return result

\section*{B. 14 maxism.c}
```

    #include <stdbool.h>
    \#include <math.h>

# \#include "topflow.h"

4
5
_declspec(dllexport) void maxmism(int pqv, int nbuses, int *buscod,
int *mismloc, double *pinj, double *qinj){
6
7 int ib;
8
9
10
11
12
13
14
15
16
17
18
19
20
21
22
23
}
24
25 }

```

\section*{B. 15 mismat.c}

1
2 \#include "math.h"
3
4
5 void pmismat (int ib, int *genbus, int *buscod, double *pgen, double * pload, double *pinj,
double *beta, double pdelta)
7 \{
8 //This subroutine calculates the active power mismatch and
updates the active generation
//for the bus ib
int ig = genbus[ib];
if(abs(buscod[ib]) == 2) \{
//pinj is now the mismatch of the PV-bus (including Varlimited PV-bus)
pinj[ib] = pgen[ig]-pload[ib]-pinj[ib]-beta[ib]*pdelta;
\}
else if (buscod[ib] == 3) \{
//Update active generation for the slack bus
pgen[ig]=pinj[ib]+pload[ib] + beta[ib]*pdelta;
//pinj[ib] = 0 ;
\}
else\{
// Active mismatches of PQ -buses
pinj[ib]=-pload[ib]-pinj[ib]-beta[ib]*pdelta;
\}
5 \}
26
27 void qmismat(int ib, int *genbus, int *buscod, double *qgen, double * qload, double *qinj, double *alfa, double qdelta)
28
29 \{
30

32
33
34
35
        //This subroutine calculates the reactive power missmatch and
        updates the reactive generation
        //for the bus ib
        int ig = genbus[ib];
        if(buscod[ib] >= 2) \{
            //Update reactive genertation at PV- and slack-buses
            qgen[ig]= qinj[ib]+qload[ib] + alfa[ib]*qdelta;
            //qinj[ib] \(=0\);
        \}
        else if(buscod[ib] == -2) \{
            //Reactive mismatch of the Var-limited generator buses
            qinj[ib] = qgen[ig]-qload[ib]-qinj[ib]-alfa[ib]*qdelta;
        \}
```

44 else{
//Reactive mismatches of PQ-buses
qinj[ib]=-qload[ib]-qinj[ib]-alfa[ib]*qdelta;
}
48 }
4 9
50
declspec(dllexport) void mismat(int pqv, int nbuses, int *genbus,
int *buscod,
double *pgen, double *qgen, double *pload, double *
qload, double *pinj, double *qinj,
double *alfa, double *beta, double pdelta, double
qdelta)
53 {
55 //This subroutine calculates the relevant mismatches, and updates
the active and/or reactive generations
//pqv decides what to calculate
//pqv = 1: Active power, 2: Reactive power, 3: Both active and
reactive power
for (ib=0;ib<nbuses;ib++) {
if (buscod[ib] != 4)
{
if(pqv != 1){
qmismat(ib, genbus, buscod, qgen, qload, qinj, alfa,
qdelta);
}
if(pqv !=2){
pmismat(ib, genbus, buscod, pgen, pload, pinj, beta,
pdelta);
}
6 7
}
68
69 }
70 }

```

\section*{B. 16 netinj.c}
```

    #include <stdbool.h>
    #include <math.h>
    #include "jacobi.h"
    4
5
_declspec(dllexport) void netinj(int pqv, int nbuses, int nlines,
double *pinj, double *qinj, double *vomag, double *voang,
double *gii, double *bii, double *gij, double *bij,
double *ratio, double *shift_angle, int *ifrom, int *ito, bool *
ibstat)
{
8 int i, j, k;
9 double volt, volt2, delta_angle, sina, cosa;
1 0
11
12
13
14
15
16
];
$\qquad$

``` double *pinj, double *qinj, double *vomag, double *voang, double *gii, double *bii, double *gij, double *bij, double *ratio, double *shift_angle, int *ifrom, int *ito, bool * ibstat)
7 \{
    double uij[1],tij[1], uji[1],tji[1];
    /*Calculate the injection on all buses*/
    //pqv decides what netinj calculates:
    //1: pinj, 2: qinj, 3: both pinj and qinj
    for (i=0; i<nbuses; i++)
            {
            volt2 = vomag[i]*vomag[i];
            if(pqv!= 1) {
                qinj[i] = -volt2*bii[i];
            }
            if(pqv != 2)
                pinj[i] = volt2*gii[i];
            }
        for (k=0; k<nlines; k++)
            {
                if (ibstat[k])
                                    i = ifrom[k];
                                    j = ito[k];
                                    volt = vomag[i]*vomag[j];
                            delta_angle = voang[i] - voang[j] - shift_angle[k
                    t_u(uij,tij,delta_angle,gij[k],bij[k],ratio[k]);
                    t_u(uji,tji,-delta_angle,gij[k],bij[k],ratio[k]);
                    if(pqv != 1){
                            qinj[i] -= volt**uij;
                            qinj[j] -= volt**uji;
}
if(pqv != 2){
                                    pinj[i] -= volt**tij;
                                    pinj[j] -= volt**tji;
```

```
4 5
46
47
48 }
4 9
50
```

$\qquad$

```
declspec(dllexport) void netpinj(int nbuses, int nlines, double *
    pinj, double *vomag, double *voang,
    double *ratio, int *ifrom, int *ito, bool *ibstat)
52 {
53 int i, j, k;
54 double volt, volt2, delta_angle, sina, cosa;
55 double uij[1],tij[1], uji[1],tji[1];
56
5 7
58
59 for (i=0; i<nbuses; i++)
60
6 1
62
63
6 4
6 5
\(\qquad\)
``` declspec (dllexport) void netqinj(int nbuses, int nlines, double * qinj, double *vomag, double *voang,
                        double *gii, double *bii, double *gij, double *bij,
    double *ratio, int *ifrom, int *ito, bool *ibstat)
85 {
86 int i, j, k;
87 double volt, volt2, delta_angle, sina, cosa;
88 double uij[1],tij[1], uji[1],tji[1];
90 /*Calculate the injection on all buses*/
```

91
92
93
94
95
96
97
98

```
for (i=0; i<nbuses; i++)
```

for (i=0; i<nbuses; i++)
{
{
volt2 = vomag[i]*vomag[i];
volt2 = vomag[i]*vomag[i];
qinj[i] = -volt2*bii[i];
qinj[i] = -volt2*bii[i];
}
}
for (k=0; k<nlines; k++)
for (k=0; k<nlines; k++)
{
{
if (ibstat[k])
if (ibstat[k])
{
{
i = ifrom[k];
i = ifrom[k];
j = ito[k];
j = ito[k];
volt = vomag[i]*vomag[j];
volt = vomag[i]*vomag[j];
delta_angle = voang[i] - voang[j];
delta_angle = voang[i] - voang[j];
t_u(uij,tij,delta_angle,gij[k],bij[k],ratio[k]);
t_u(uij,tij,delta_angle,gij[k],bij[k],ratio[k]);
t_u(uji,tji,-delta_angle,gij[k],bij[k],ratio[k]);
t_u(uji,tji,-delta_angle,gij[k],bij[k],ratio[k]);
qinj[i] -= volt**uij;
qinj[i] -= volt**uij;
qinj[j] -= volt**uji;
qinj[j] -= volt**uji;
}
}
}

```
        }
```


## B. 17 sConstruct.py

```
#This file is used bu SCons to build the shared library "clibrary"
#Clibrary contains all the C-extesnions of Topflow
4 \text { SharedLibrary(}
5 "clibrary",
6 [
7 "set_isa.c",
8 "flatstart.c",
9 "netinj.c",
10 "mismat.c",
11 "maxmism.c",
12 "zerosp.c",
13 "enforce_qlim.c"
14 "set_rhs.c",
15 "jacobi.c",
16 "update_voltages.c",
17 "coo_conv.c",
18 "bmatrix.c",
19 "hlmatrix.c",
20 "selectver.c"
21 ]
```

3
22 )

## B. 18 select_ver.c

```
#include<stdbool.h>
2 #include<math.h>
3 #include<string.h>
4
5 #define ABS(x) (((x) < 0) ? (-(x)) : ((x))) /* if loop: ABS(X) gives
    the absolute value of x*/
6
7
_declspec(dllexport) void selectver(double *xp,double *xq,double *
    xinv, double *bij, int nlines, char *bver){
        int k;
        //strcmp compare two strings
        if(strcmp(bver, "XX")) {
            //STD version
            for(k = 0; k<nlines;k++) {
                xp[k] = xinv[k];
                xq[k] = xinv[k];
                }
        }
        else if(strcmp(bver,"XB")) {
            //XB-version (dual)
            for(k = 0; k<nlines;k++) {
                xp[k] = xinv[k];
                xq[k] = ABS(bij[k]);
```

```
22
23
    }
    else if(strcmp(bver, "BX")){
        //BX-version (primal)
        for(k = 0; k<nlines;k++) {
            xp[k] = ABS(bij[k]);
            xq[k] = xinv[k];
        }
    }
    else if(strcmp(bver, "BB")){
        //BB version
        for(k = 0; k<nlines;k++) {
        xp[k] = ABS(bij[k]);
        xq[k] = ABS(bij[k]);
    }
37
38 }
```


## B. 19 set_rhs.c

```
#include <stdbool.h>
2
3
4
_declspec(dllexport) void set_rhs(int pqv, int nbuses, int *buscod,
    double *vomag, double *pinj, double *qinj, double *rhs){
5
6 int iq, ip, ib;
7 iq = ip = 0;
8
9
10
11
12
13
14
15
16
17
18
1 9
20
21
22
23
30 }
            for( ib = 0; ib <nbuses; ib++) {
            if(pqv == 3) {
                //Right hand side of the Newton Rapshon equations
                if(buscod[ib] != 3){
                    rhs[ip] = pinj[ib];
                if(buscod[ib] != 2) {
                    rhs[iq + nbuses-1] = qinj[ib];
                            iq += 1;
                }
                ip ++;
                }
            }
            else if(pqv == 1){
                //Right hand side of the active power equations in
    FDLF
        rhs[ib] = pinj[ib]/vomag[ib];
            }
            else if(pqv == 2){
                        //Right hand side of the reactive power equations in
    FDLF
                        rhs[ib] = qinj[ib]/vomag[ib];
            }
```


## B. 20 settings.py

\#This file contains the Setting-class, which is used to specify
2 \#the settings of the loadflow-functions (currently acsolve and decsolve)
3 class Settings:

4
def __init__(self, version = 'NR', flat_start = True, enf_qlim =
False, max_it $=20$, conv_tol $=0.000001$, pqv $=$ None, filename $=$
None, save_path = None, save_verbose = 1, print_verbose = 1):
self.version = version
self.flat_start = flat_start
self.enf_qlim = enf_qlim
self.max_it = max_it
self.conv_tol = conv_tol
self.pqv = pqv
self.bver = 'XX' \#Default
self.set_fdlf_ver(version, pqv) \#sets self.bver and pqv
self.filename = filename
self.save_path = save_path
self.save_verbose = save_verbose
\# 1: Save only the final result (Default).
\# 2: Save the final reslt in one sheet, and the iteration-
summaries in a second sheet.
self.print_verbose = print_verbose
\# 0: Do NOT print anything to screen.
\# 1: Print the final result to screen: Convergence/Divergence
- (Default)
\# 2: Print a more verbose final result-summary to the screen.
\# 3: Same as 2: + print a summery of each iteration
def print_settings(self):
print('\nVersion:',self.version,
'\nflat_start:',self.flat_start,
'\nEnforce reactive limits:', self.enf_qlim,
'\nmaximum iterations:', self.max_it,
'\nconvergence tolerance',self.conv_tol)
def set_fdlf_ver(self, version, pqv):
if(version in ['FDXB', 'FDBX', 'FDXX', 'FDBB']):
self.bver = version[2:]
if(pqv is None):
\#if the iteration-version is not specified, the
default values for the variuos FDLF-versions are used
self.pqv $=2$ if self.version $==$ 'FDXB' else 1
elif(pqv in $[1,2]):$
self.pqv = pqv
else:
print('unvalid value for pqv')

```
def set_var(self, name, value):
    #set the instance variables
    setattr(self,name, value)
    if(name in ['version', 'pqv']):
                #update bver and pqv
                self.set_fdlf_ver(self.version, self.pqv)
```


## B. 21 topflow.h

1 \#ifndef HEADER_nettanalyse /*to prevent double declaration of identifiers*/
2 \#define HEADER_nettanalyse
3
4 \#define PQ_BUS 1
5 \#define PV_BUS 2
6 \#define SWING_BUS 3
7 \#define DISCONNECTED_BUS 4
8
9
10 \#define $\operatorname{ABS}(x)(((x)<0)$ ? (-(x)) : ((x))) /* if loop: ABS (X) gives the absolute value of $x * /$
11 \#define CONNECTED(i, con) (con[i] < 0 ? i : con[i])
12
13 \#endif

## B. 22 update_voltages.py

```
#include <stdbool.h>
3
_declspec(dllexport) void update_voltages(int pqv, int nbuses, int *
    buscod, double *vomag, double *voang, double *correction){
        int ip, iq, ib;
        ip = iq = 0;
        for( ib = 0; ib <nbuses; ib++) {
            if( pqv == 3){
                //Update for Newton Rapshon algorithm
                if(buscod[ib] != 3) {
                voang[ib] += correction[ip];
                ip ++;
                if(buscod[ib] != 2) {
                                    vomag[ib] += correction[iq + nbuses-1];
                                    iq ++;
                                    }
                        }
                }
                else if(pqv == 1){
                        //Update voltage anges (Active power eq fro FDLD)
                        if(buscod[ib] != 3){
                voang[ib] += correction[ib];
                        }
                }
                else if(pqv == 2){
                        //Update voltage magnitudes (Reactive power eq for FDLF)
                        if(buscod[ib] != 3 && buscod[ib] != 2){
                                    vomag[ib] += correction[ib];
                            }
                        }
        }
}
```

2

## B. 23 zerosp.c

```
    #include <stdbool.h>
2 #include <math.h>
3
4 __declspec(dllexport) void zerosp(int nlines, int *ifrom, int *ito,
    bool *ibstat, double *xinv, double *xinj){
        int il, ifr, itr;
        for(il = 0; il < nlines; il++) {
            if(fabs(xinv[il] > 9000 && ibstat[il] == true)){
                    ifr = ifrom[il];
                        itr = ito[il];
```

$\}$

```
        if(fabs(xinj[ifr]) > fabs(xinj[itr])) {
        xinj[ifr]=xinj[ifr]+xinj[itr];
        xinj[itr]=0.0;
            }
            else{
                xinj[itr]=xinj[itr]+xinj[ifr];
        xinj[ifr]=0.0;
            }
```

\}
$\}$

## C Tests

## C. 1 pandapower_delay.py

```
1 import timeit
2
3 #Measure the delay of the numba-compiler in pandapower
4
5 mysetup = (
6
7 '''import pandapower as pp
import pandapower.networks as pn
9 net = pn.case14()
10 '''
11)
12
13 numba_delay = timeit.timeit( setup = mysetup, stmt = 'pp.runpp(net)',
            number = 1)
14 runtime_case14 = timeit.timeit( setup = mysetup, stmt = 'pp.runpp(net
    )', number = 100)/100
15
16 numba_delay = numba_delay - runtime_case14
17
18 print('Numba delay:', numba_delay)
```


## C. 2 reliability_acsolve.py

```
1 import sys
2 sys.path.append("..")
3 import topflow as tf
import numpy as np
import math
6 \text { import pandapower as pp}
7mport pandapower.networks as pn
8
# #Reliability test
10 #Compare topflow results to pandapower and pypower
11
12 #Initialize topflow cases
13 case4gs = tf.example_case('case4gs')
14 case14 = tf.example_case('case14')
15 case30 = tf.example_case('case30')
16 case118 = tf.example_case('case118')
17 case300 = tf.example_case('case300')
18 case1354 = tf.example_case('case1354pegase')
19 case2869 = tf.example_case('case2869pegase')
20 case9241 = tf.example_case('case9241pegase')
21
22 #Solve topflow cases
23 result4gs = tf.loadflow(case4gs, conv_tol = 0.00000001, print_verbose
    = 0)
2 4 ~ r e s u l t 1 4 ~ = ~ t f . l o a d f l o w ( c a s e 1 4 , ~ c o n v \_ t o l ~ = ~ 0 . 0 0 0 0 0 0 0 1 , ~ p r i n t \_ v e r b o s e
        = 0)
25 result30 = tf.loadflow(case30, conv_tol = 0.00000001, print_verbose =
            0)
26 result118 = tf.loadflow(case118, conv_tol = 0.00000001,
        print_verbose = 0)
27 result300 = tf.loadflow(case300, conv_tol = 0.00000001, print_verbose
            = 0)
28 result1354 = tf.loadflow(case1354 , conv_tol = 0.00000001,
    print_verbose = 0)
29 result2869 = tf.loadflow(case2869 , conv_tol = 0.00000001,
    print_verbose = 0)
30 result9241 = tf.loadflow(case9241 , conv_tol = 0.00000001,
    print_verbose = 0)
31
32 #initialize pandapower cases
33 net4gs = pn.case4gs()
34 net14 = pn.case14()
35 net30 = pn.case30()
36 net118 = pn.case118()
37 net300 = pn.case300()
38 net1354 = pn.case1354pegase()
39 net2869 = pn.case2869pegase()
40 net9241 = pn.case9241pegase()
```

```
4 1
42 #solve pandapower cases
43 pp.runpp (net4gs)
44 pp.runpp (net14)
45 pp.runpp (net30)
46 pp.runpp (net118)
47 pp.runpp (net300)
48 pp.runpp(net1354)
49 pp.runpp(net2869)
50 pp.runpp(net9241)
51
5 2
53 #compare case4gs
54 panda_vm_dev = 0
55 panda_va_dev = 0
56 worst_panda_vm_bus = 0
57 worst_panda_va_bus = 0
58 pyp_vm_dev = 0
59 pyp_va_dev = 0
60 worst__pyp_vm_bus = 0
61 worst_pyp_va_bus = 0
6 2
63 for i in range(0, result4gs.vomag.size):
64 ext = case4gs.int2ext(i)
65 if abs(result4gs.get('bus', ext, 'vm')-net4gs.res_bus.vm_pu[i])>
    abs(panda_vm_dev):
66 panda_vm_dev = result4gs.get('bus', ext, 'vm')-net4gs.res_bus
    .vm_pu[i]
                worst_panda_vm_bus =ext
        if abs(result4gs.get('bus', ext, 'va')-net4gs.res_bus.va_degree[i
        ])> abs(panda_va_dev):
                panda_va_dev = result4gs.get('bus', ext, 'va')-net4gs.res_bus
        .va_degree[i]
                worst_panda_va__bus = ext
        if abs(result4gs.vomag[i]-case4gs.vomag[i])> abs(pyp_vm_dev):
                pyp_vm_dev = result4gs.vomag[i]-case4gs.vomag[i]
                worst_pyp_vm_bus = case4gs.int2ext(i)
        if abs(result4gs.voang[i]-case4gs.voang[i])> abs(pyp_va_dev):
                pyp_va_dev = result4gs.voang[i]-case4gs.voang[i]
                worst_pyp_va_bus = case4gs.int2ext(i)
7 7
78 print('\nPandaPower:')
79 print('Worst VM devition case4gs:', panda_vm_dev, 'pu, at bus:',
    worst_panda_vm_bus)
80 print('Worst VA devition case4gs:', panda_va_dev, 'deg, at bus:',
    worst_panda_va_bus)
81 print('PyPower:')
82 print('Worst VM devition case4gs:', pyp_vm_dev, 'pu, at bus:',
    worst_pyp_vm_bus)
83 print('Worst VA devition case4gs:', pyp_va_dev, 'deg, at bus:',
```

```
    worst_pyp_va__bus)
84
85
86 #compare case14
87 panda_vm_dev = 0
88 panda_va_dev = 0
89 worst_panda_vm_bus = 0
90 worst_panda_va_bus = 0
91 pyp_vm_dev = 0
92 pyp_va_dev = 0
93 worst__pyp_vm_bus = 0
94 worst__pyp_va_bus = 0
95
96 for i in range(0, result14.vomag.size):
97 ext = case14.int2ext(i)
98 if abs(result14.get('bus', ext, 'vm')-net14.res_bus.vm_pu[i])>
    abs(panda_vm_dev):
                panda_vm_dev = result14.get('bus', ext, 'vm')-net14.res_bus.
    vm_pu[i]
                worst_panda_vm_bus =ext
            if abs(result14.get('bus', ext, 'va')-net14.res_bus.va_degree[i])
    > abs(panda_va_dev):
                    panda_va_dev = result14.get('bus', ext, 'va')-net14.res_bus.
    va_degree[i]
                worst_panda_va_bus = ext
            if abs(result14.get('bus',ext,'vm')-case14.get('bus', ext, 'vm'))
    > abs(pyp_vm_dev):
            pyp_vm_dev = result14.get('bus',ext,'vm')-case14.get('bus',
    ext, 'vm')
            worst_pyp_vm_bus = ext
            if abs(result14.get('bus',ext,'va')-case14.get('bus', ext, 'va'))
        > abs(pyp_va_dev):
            pyp_va_dev = result14.get('bus',ext,'va')-case14.get('bus',
        ext, 'va')
            worst_pyp_va_bus = ext
110
111 print('\nPandaPower:')
112 print('Worst VM devition case14:', panda_vm_dev, 'pu, at bus:',
    worst_panda_vm_bus)
113 print('Worst VA devition case14:', panda_va_dev, 'deg, at bus:',
    worst_panda_va_bus)
114 print('PyPower:')
115 print('Worst VM devition case14:', pyp_vm_dev, 'pu, at bus:',
    worst_pyp_vm_bus)
116 print('Worst VA devition case14:', pyp_va_dev, 'deg, at bus:',
    worst_pyp_va__bus)
117
118
119 #compare case30
120 panda_vm_dev = 0
```

121 panda_va_dev = 0
122 worst_panda_vm_bus $=0$
123 worst_panda_va_bus $=0$
124 pyp_vm_dev $=0$
125 pyp_va_dev $=0$
126 worst_pyp_vm_bus $=0$
127 worst_pyp_va_bus $=0$
128
129 for i in range(0, result 30. vomag.size) :
130 ext $=$ case 30 .int2ext(i)
131 if abs(result30.get('bus', ext, 'vm')-net30.res_bus.vm_pu[i])>
abs (panda_vm_dev) :
132 panda_vm_dev = result30.get('bus', ext, 'vm')-net30.res_bus.
vm_pu[i]
worst_panda_vm_bus =ext
if abs(result30.get('bus', ext, 'va')-net30.res_bus.va_degree[i])
> abs(panda_va_dev):
panda_va_dev = result30.get('bus', ext, 'va')-net $30 . r e s \_b u s$.
va_degree[i]
worst_panda_va_bus = ext
if abs(result30.get('bus', ext, 'vm')-case30.get('bus', ext, 'vm')) > abs(pyp_vm_dev):

```
                pyp_vm_dev = result30.get('bus',ext,'vm')-case30.get('bus',
```

ext, 'vm')
worst_pyp_vm_bus = ext
if abs(result30.get('bus',ext,'va')-case30.get('bus', ext, 'va')) > abs(pyp_va_dev):
pyp_va_dev $=$ result 30 .get('bus', ext, 'va')-case30.get('bus', ext, 'va')
worst_pyp_va_bus = ext
143
144 print('\nPandaPower:')
145 print('Worst VM devition case30:', panda_vm_dev, 'pu, at bus:', worst_panda_vm_bus)
146 print('Worst VA devition case30:', panda_va_dev, 'deg, at bus:', worst_panda_va_bus)
147 print ('PyPower:')
148 print('Worst VM devition case30:', pyp_vm_dev, 'pu, at bus:', worst_pyp_vm_bus)
149 print('Worst VA devition case30:', pyp_va_dev, 'deg, at bus:', worst_pyp_va_bus)
150
151
152 \#compare case118
153 panda_vm_dev $=0$
154 panda_va_dev = 0
155 worst_panda_vm_bus = 0
156 worst_panda_va_bus $=0$
157 pyp_vm_dev $=0$
158 pyp_va_dev $=0$

```
worst_pyp_vm_bus = 0
```

worst_pyp_va_bus = 0
161
162 for i in range(0, result118.vomag.size):
163 ext $=$ case118.int2ext(i)
178 print('\nWorst VM devition case118:', panda_vm_dev, 'pu, at bus:',
worst_panda_vm_bus)
179 print('Worst VA devition case118:', panda_va_dev, 'deg, at bus:',
worst_panda_va_bus)
180 print ('PyPower:')
181 print('Worst VM devition case118:', pyp_vm_dev, 'pu, at bus:',
worst_pyp_vm_bus)
182 print('Worst VA devition case118:', pyp_va_dev, 'deg, at bus:',
worst_pyp_va_bus)
183
184
185 \#compare case300
186 panda_vm_dev $=0$
187 panda_va_dev = 0
188 worst_panda_vm_bus $=0$
189 worst_panda_va_bus $=0$
190 pyp_vm_dev $=0$
191 pyp_va_dev = 0
192 worst_pyp_vm_bus = 0
193 worst_pyp_va_bus = 0
194
195 for i in range(0, result 300. vomag.size):
196 ext $=$ case300.int2ext(i)
210 print ('\nPandaPower:')
211 print('Worst VM devition case300:', panda_vm_dev, 'pu, at bus:',
worst_panda_vm_bus)
212 print('Worst VA devition case300:', panda_va_dev, 'deg, at bus:',
worst_panda_va_bus)

213 print ('PyPower:')
214 print('Worst VM devition case300:', pyp_vm_dev, 'pu, at bus:', worst_pyp_vm_bus)
215 print('Worst VA devition case300:', pyp_va_dev, 'deg, at bus:', worst_pyp_va_bus)
216
217
218 \#compare case1354
219 panda_vm_dev = 0
220 panda_va_dev = 0
221 worst_panda_vm_bus $=0$
222 worst_panda_va_bus $=0$
223 pyp_vm_dev $=0$
224 pyp_va_dev $=0$
225 worst_pyp_vm_bus $=0$
226 worst_pyp_va_bus = 0
227
228 for i in range(0, result1354.vomag.size):
229 ext = case1354.int2ext(i)
230 if abs(result1354.get('bus', ext, 'vm')-net1354.res_bus.vm_pu[i]) > abs(panda_vm_dev):
panda_vm_dev $=$ result1354.get('bus', ext, 'vm')-net1354.
res_bus.vm_pu[i]
worst_panda_vm_bus =ext
243 print('\nWorst VM devition case1354 :', panda_vm_dev, 'pu, at bus:',
worst_panda_vm_bus)
244 print('Worst VA devition case1354 :', panda_va_dev, 'deg, at bus:',
worst_panda_va_bus)
245 print ('PyPower:')
246 print('Worst VM devition case1354:', pyp_vm_dev, 'pu, at bus:',
worst_pyp_vm_bus)
247 print('Worst VA devition case1354:', pyp_va_dev, 'deg, at bus:',
worst_pyp_va_bus)

248
249 \#compare case2869
250 panda_vm_dev = 0
251 panda_va_dev = 0
252 worst_panda_vm_bus = 0
253 worst_panda_va_bus = 0
254 pyp_vm_dev $=0$
255 pyp_va_dev = 0
256 worst_pyp_vm_bus = 0
257 worst_pyp_va_bus = 0
258
259
if abs(result1354.get('bus', ext, 'va')-net1354.res_bus.va_degree [i]) > abs(panda_va_dev):
panda_va_dev = result1354.get('bus', ext, 'va')-net1354.
res_bus.va_degree[i]
worst_panda_va_bus = ext
if abs(result1354.get('bus',ext,'vm')-case1354.get('bus', ext, '
vm')) > abs(pyp_vm_dev):
pyp_vm_dev = result1354.get('bus', ext,'vm')-case1354.get('bus
', ext, 'vm')
worst_pyp_vm_bus = ext
if abs(result1354.get('bus',ext,'va')-case1354.get('bus', ext, '
va')) > abs(pyp_va_dev):
', ext, 'va')
worst_pyp_va_bus = ext
\#compare case2869
panda_vm_dev $=0$
panda_va_dev $=0$
worst_panda_vm_bus $=0$
worst_panda_va_bus $=0$
pyp_vm_dev = 0
worst_pyp_vm_bus $=0$
worst_pyp_va_bus = 0
for i in range(0, result2869.vomag.size):
ext $=$ case2869.int2ext(i)
if abs(result2869.get('bus', ext, 'vm')-net2869.res_bus.vm_pu[i])
> panda_vm_dev:
panda_vm_dev = result2869.get('bus', ext, 'vm')-net2869.
res_bus.vm_pu[i]
worst_panda_vm_bus =ext
if abs(result2869.get('bus', ext, 'va')-net2869.res_bus.va_degree
[i]) > panda_va_dev:
panda_va_dev = result2869.get('bus', ext, 'va')-net2869.
res_bus.va_degree[i]
worst_panda_va_bus = ext
if abs(result2869.get('bus',ext,'vm')-case2869.get('bus', ext, '
vm')) > abs(pyp_vm_dev):
pyp_vm_dev $=$ result2869.get('bus',ext,'vm')-case2869.get('bus ', ext, 'vm')
worst_pyp_vm_bus = ext

271

272

273
274
275
worst_panda_vm_bus)
276 print('Worst VA devition case2869 :', panda_va_dev, 'deg, at bus:', worst_panda_va_bus)
277 print('PyPower:')
278 print('Worst VM devition case2869:', pyp_vm_dev, 'pu, at bus:', worst_pyp_vm_bus)
279 print('Worst VA devition case2869:', pyp_va_dev, 'deg, at bus:', worst_pyp_va_bus)
280
281 \#compare case9241
282 panda_vm_dev $=0$
283 panda_va_dev $=0$
284 worst_panda_vm_bus $=0$
285 worst_panda_va_bus $=0$
286 pyp_vm_dev $=0$
287 pyp_va_dev $=0$
288 worst_pyp_vm_bus $=0$
289 worst_pyp_va_bus $=0$
290
291 for i in range(0, result9241.vomag.size):
292
293 if abs(result9241.get('bus', ext, 'vm')-net9241.res_bus.vm_pu[i])
if abs (result2869.get('bus',ext,'va')-case2869.get('bus', ext, '
va')) > abs(pyp_va_dev):
', ext, 'va')
worst_pyp_va_bus = ext
print('\nWorst VM devition case2869 :', panda_vm_dev, 'pu, at bus:',
worst panda vm bus)
orst_PM-va_bus
for i in range(0, result9241.vomag.size):
ext $=$ case9241.int2ext(i)
> panda_vm_dev:
panda_vm_dev $=$ result 9241. get('bus', ext, 'vm')-net9241.
res_bus.vm_pu[i]
worst_panda_vm_bus =ext
if abs(result9241.get('bus', ext, 'va')-net9241.res_bus.va_degree
[i]) > panda_va_dev:
panda_va_dev = result9241.get('bus', ext, 'va')-net9241.
res_bus.va_degree [i]
worst_panda_va_bus $=$ ext
if abs(result9241.get('bus', ext, 'vm')-case9241.get('bus', ext, '
vm')) > abs(pyp_vm_dev):
pyp_vm_dev = result9241.get('bus', ext,'vm')-case9241.get('bus
', ext, 'vm')
worst_pyp_vm_bus = ext
if abs(result9241.get('bus',ext,'va')-case9241.get('bus', ext, '
va')) > abs(pyp_va_dev):
pyp_va_dev = result9241.get('bus',ext,'va')-case9241.get('bus
306 print ('\nWorst VM devition case9241 :', panda_vm_dev, 'pu, at bus:',
worst_panda_vm_bus)
307 print('Worst VA devition case9241:', panda_va_dev, 'deg, at bus:', worst_panda_va_bus)
308 print ('PyPower:')
309 print('Worst VM devition case9241:', pyp_vm_dev, 'pu, at bus:', worst_pyp_vm_bus)
310 print('Worst VA devition case9241:', pyp_va_dev, 'deg, at bus:', worst_pyp_va_bus)

## C. 3 reliability_decsolve.py

```
1 import sys
2 sys.path.append("..")
import topflow as tf
4
# #this test compares decsolve with acsolve. acsolve has been tested
    against other programs and should give reliable results
# Initialize topflow cases
7
8
9 for i in range(0,7):
    if(i == 0):
        print('\n##### case4gs #####')
        #initialize the case
        case = tf.example_case('case4gs')
    elif(i == 1):
        print('\n##### case14 #####')
        #initialize the case
        case = tf.example_case('case14')
    elif(i == 2):
        print('\n##### case30 #####')
        #initialize the case
        case = tf.example_case('case30')
    elif(i == 3):
        print('\n##### case118 #####')
        #initialize the case
        case = tf.example_case('case118')
    elif(i == 4):
        print('\n##### case300 #####')
        #initialize the case
        case = tf.example_case('case300')
    elif(i == 5):
        print('\n##### 1354pegase #####')
        #initialize the case
        case = tf.example_case('case1354pegase')
    elif(i == 6):
        print('\n##### 2869pegase #####')
        #initialize the case
        case = tf.example_case('case2869pegase')
    elif(i == 7):
        print('\n##### case9241pegase #####')
4 7
```

```
        #initialize the case
        case = tf.example_case('case9241pegase')
    #initialize the worst deviations
    worst_vm_fdxx = 0
    worst_va_fdxx = 0
    worst_vm_fdbx = 0
    worst_va_fdbx = 0
    worst_vm_fdxb = 0
    worst_va_fdxb = 0
    #run the load-flows
    resnr = tf.loadflow(case, version = 'NR', print_verbose = 0)
    resxx = tf.loadflow(case, version = 'FDXX', print_verbose = 0)
    resbx = tf.loadflow(case, version = 'FDBX', print_verbose = 0)
    resxb = tf.loadflow(case, version = 'FDXB', print_verbose = 0)
    #Compare the results, if the loadflow()-function returns 0 it
means the test diverged
    for i in range(0, resnr.vomag.size):
    #FDXX:
    if(resxx == 0):
        break
    if abs(resnr.vomag[i]-resxx.vomag[i])> abs(worst_vm_fdxx):
        worst_vm_fdxx = abs(resnr.vomag[i]-resxx.vomag[i])
    if abs(resnr.voang[i]-resxx.voang[i])> abs(worst_va_fdxx):
        worst_va_fdxx = abs(resnr.voang[i]-resxx.voang[i])
    #FDBX:
    if(resbx == 0):
        break
    if abs(resnr.vomag[i]-resbx.vomag[i])> abs(worst_vm_fdbx):
        worst_vm_fdbx = abs(resnr.vomag[i]-resbx.vomag[i])
    if abs(resnr.voang[i]-resbx.voang[i])> abs(worst_va_fdbx):
        worst_va_fdbx = abs(resnr.voang[i]-resbx.voang[i])
    #FDXB
    if(resxb == 0):
        break
    if abs(resnr.vomag[i]-resxb.vomag[i])> abs(worst_vm_fdxb):
        worst_vm_fdxb = abs(resnr.vomag[i]-resxb.vomag[i])
    if abs(resnr.voang[i]-resxb.voang[i])> abs(worst_va_fdxb):
        worst_va_fdxb = abs(resnr.voang[i]-resxb.voang[i])
    #Print the result:
    #FDXX
if(resxx == 0):
    print('\nFDXX: DIVERGENCE')
```

```
    else:
        print('\nFDXX: CONVERGENCE: number of iterations:', resxx.
        icount)
            print('Worst VM deviation:', worst_vm_fdxx)
            print('Worst VA deviation:', worst_va_fdxx)
        #FDBX
        if(resbx == 0):
        print('\nFDBX: DIVERGENCE')
        else:
        print('\nFDBX: CONVERGENCE: number of iterations:', resbx.
        icount)
        print('Worst VM deviation:', worst_vm_fdbx)
        print('Worst VA deviation:', worst_va_fdbx)
    #FDXB
    if(resxb == 0):
        print('\nFDXB: DIVERGENCE')
    else:
        print('\nFDXB: CONVERGENCE: number of iterations:', resxb.
    icount)
        print('Worst VM deviation:', worst_vm_fdxb)
        print('Worst VA deviation:', worst_va_fdxb)
```


## C. 4 speed_acsolve.py

1
2 import timeit
3 from pypower.api import case14 as pypcase14
4 from pypower.api import case30 as pypcase30
5 from pypower.api import runpf, ppoption
6
7 \#This tests compares the speed of running a Newton-Rapshon loadflow with Topflow
8 \#With the speed of pandapower and pypower
9
10 \#Measure the execution time of topflow
11 mysetup_tf = (
12
13 import sys
14 sys.path.append("..")
15 import topflow as tf
16
17 ' '
18 )
19
20 tf_runtime14 = timeit.timeit(setup = mysetup_tf +
21 ' '
22 case = tf.example_case("case14")
23 ' ' ',
24 stmt $=$ 'tf.loadflow(case, conv_tol $=0.00000001$, print_verbose $=0$ )',
25 number $=100) / 100$
26
27 tf_runtime30 = timeit.timeit(setup = mysetup_tf +
28 ''
29 case = tf.example_case("case30")
30 ' ' ',
31 stmt = 'tf.loadflow(case, conv_tol = 0.00000001, print_verbose = 0)',
32 number $=100) / 100$
33
34 tf_runtime1354 = timeit.timeit(setup = mysetup_tf +
35
36

38 stmt $=$ 'tf.loadflow(case, conv_tol $=0.00000001$, print_verbose $=0$ )',
39 number $=100) / 100$
40
41 tf_runtime2869 = timeit.timeit(setup = mysetup_tf +
42
43
44
45
46 number $=100) / 100$
47

```
48 tf_runtime9241 = timeit.timeit(setup = mysetup_tf +
'''
case = tf.example_case("case9241pegase")
''',
stmt = 'tf.loadflow(case, conv_tol = 0.00000001, print_verbose = 0)',
number = 100)/100
#Measure the exectution time of pypower
56 mysetup_pyp14 = (
57 '''
5 8 \text { from pypower.api import runpf, ppoption}
59 from pypower.api import case14
60 ppopt = ppoption(VERBOSE = 0, OUT_ALL = 0)
61 case14 = case14()
62 ''')
6 3
6 4
65 pyp_runtime14 = timeit.timeit(setup = mysetup_pyp14, stmt = 'ppr =
    runpf(case14, ppopt)[0]', number = 100)/100
6 6
6 7
68 mysetup_pyp30 = (
69 '''
70 from pypower.api import runpf, ppoption
71 from pypower.api import case30
72 ppopt = ppoption(VERBOSE = 0, OUT_ALL = 0)
73 case30 = case30()
74 '''
75 )
7 6
77 pyp_runtime30 = timeit.timeit(setup = mysetup_pyp30, stmt = 'ppr =
        runpf(case30, ppopt)[0]', number = 100)/100
78
7 9
80 #The cases which are not available in pypower are converted
81 #to pypower-cases from pandapower by using the pandapower.converter
82 mysetup_pyp1354 = (
83 '''
84 from pypower.api import runpf, ppoption
85 import pandapower.converter as pc
86 import pandapower.networks as pn
87 ppopt = ppoption(VERBOSE = 0, OUT_ALL = 0)
88 case1354 = pc.to_ppc(pn.case1354pegase())
89 '''
90 )
91
92 pyp_runtime1354 = timeit.timeit(setup = mysetup_pyp1354, stmt = 'ppr
        = runpf(case1354, ppopt)[0]', number = 100)/100
93
94 mysetup_pyp2869 = (
```

```
95
96 from pypower.api import runpf, ppoption
97 import pandapower.converter as pc
98 import pandapower.networks as pn
99 ppopt = ppoption(VERBOSE = 0, OUT_ALL = 0)
100 case2869 = pc.to_ppc(pn.case2869pegase())
101 '''
102 )
103
104 pyp_runtime2869 = timeit.timeit(setup = mysetup_pyp2869, stmt ='ppr =
                        runpf(case2869, ppopt)[0]', number = 100)/100
105
106 mysetup_pyp9241 = (
107 '''
108 from pypower.api import runpf, ppoption
109 import pandapower.converter as pc
110 import pandapower.networks as pn
111 ppopt = ppoption(VERBOSE = 0, OUT_ALL = 0)
112 case9241 = pc.to_ppc(pn.case9241pegase())
113 '''
114 )
115
116 pyp_runtime9241 = timeit.timeit(setup = mysetup_pyp9241, stmt ='ppr =
        runpf(case9241, ppopt)[0]', number = 100)/100
117
118
119 #Measure the execution time of pandapower
120 mysetup_pp = (
1 2 1
122 import pandapower as pp
123 import pandapower.networks as pn
124 net = pn.case14()
125 pp.runpp (net)
126 '''
127 )
128
129 pp_runtime14 = timeit.timeit(setup = mysetup_pp +
130 '''
131 net = pn.case14()
132 ''',
133 stmt = 'pp.runpp(net)',
134 number = 100)/100
135
136 pp_runtime30 = timeit.timeit(setup = mysetup_pp +
137
138 net = pn.case30()
139 ''',
140 stmt = 'pp.runpp(net)',
141 number = 100)/100
142
```

```
143 pp_runtime1354 = timeit.timeit(setup = mysetup_pp +
144 '''
145 net = pn.case1354pegase()
146 ''',
147 stmt = 'pp.runpp(net)',
148 number = 100)/100
149
150 pp_runtime2869 = timeit.timeit(setup = mysetup_pp +
151
152 net = pn.case2869pegase()
153 ''',
154 stmt = 'pp.runpp(net)',
155 number = 100)/100
156
157 pp_runtime9241 = timeit.timeit(setup = mysetup_pp +
158 '''
159 net = pn.case9241pegase()
160 ''',
161 stmt = 'pp.runpp(net)',
162 number = 100)/100
163
164
165 #Print the result to screen
166 print('\nRuntimes on case14:')
167 print('Topflow:', tf_runtime14)
168 print('pandapower:', pp_runtime14)
169 print('pypower:', pyp_runtime14)
170
171 print('\nRuntimes on case30:')
172 print('Topflow:', tf_runtime30)
173 print('pandapower:', pp_runtime30)
174 print('pypower:', pyp_runtime30)
175
176 print('\nRuntime on case1354pegase:')
177 print('Topflow:', tf_runtime1354)
178 print('pandapower:', pp_runtime1354)
179 print('pypower:', pyp_runtime1354)
180
181 print('\nRuntime on case2869pegase:')
182 print('Topflow:', tf_runtime2869)
183 print('pandapower:', pp_runtime2869)
184 print('pypower:', pyp_runtime2869)
185
186 print('\nRuntime on case9241pegase:')
187 print('Topflow:', tf_runtime9241)
188 print('pandapower:', pp_runtime9241)
189 print('pypower:', pyp_runtime9241)
```


## C. 5 test_acsolve_integration.py

```
import sys
sys.path.append("..")
3 import topflow as tf
4 import topflow.acsolve_wrapper as ac
import numpy as np
6 \text { from numpy.testing import assert_array_equal}
7 from numpy.testing import assert_array_almost_equal
8 import pytest
9 import math
#This is a integration-test for acsolve
13 @pytest.fixture
14 def case3():
15 case3 = tf.Case()
16 #Initialize the Case-object with the data of casel4
17 case3.icount == 0
18 case3.convergence == False
19 case3.sbase == 100
import sys
    # Bus data
    case3.nbuses = 3 # number of buses
    case3.slackbusnr= 2 #the internal number of the slacbus
    case3.numbus = np.array([0,1,2], dtype=np.int32)
    case3.buscod = np.array([1, 1, 3], dtype=np.int32) # bus type.
    1=PQ, 2=PV, 3=Slack, 4= disconnected
    case3.pload = np.array([1.0, 0.5, 0.0], dtype=np.double) # pload
    [nbuses]: active load vector (fixed)
    case3.qload = np.array([0.5, 0.5, 0.0], dtype=np.double) # qload
    [nbuses]: reactive load vector (fixed)
    case3.pinj = np.zeros(case3.nbuses, dtype=np.double) #
    calculated active power injection for all buses
    case3.qinj = np.zeros( case3.nbuses, dtype=np.double) #
    calculated reactive power injection for all buses
    case3.vomag = np.ones(case3.nbuses, dtype=np.double) # voltage
    magnetudes at all buses
        case3.voang = np.zeros(case3.nbuses, dtype=np.double) # voltage
        angles at all buses
        case3.gs = np.ones(case3.nbuses)
        case3.bs = np.ones(case3.nbuses)
        # Generator data
        case3.ngens = 1
        case3.genbus = np.array([-1, -1, 0], dtype=np.int32) # genbus[
        nbuses] if a generator: genbus[k]= generator nr. If not: genbus[k
        ] =-1
        case3.vgbus = np.ones(case3.ngens, dtype=np.int32)
        case3.pgen = np.array([0.0], dtype=np.double) # pgen[ngen]:
        active power generated at generator nr
```

10
12
20
21
22
23

```
    case3.qgen = np.array([0.0], dtype=np.double) # qgen[ngen]:
    reactive power generated at generator nr
    case3.qmax = np.array([9999])
    case3.qmin = np.array([-999])
    # Line data
    case3.nlines = 3 # number of lines
    case3.gii = np.array([5.17647059, 3.17647059, 6.0], dtype=np.
double) # real part of the diagonal element in the ybus
    case3.bii = np.array([-12.70588235, -10.70588235, -14.0], dtype=
np.double) # imag part of the diagonal element in the ybus
    case3.gij = np.array([1.17647059, 4.0, 2.0], dtype=np.double) #
    real part of the off-diagonal element in the ybus
    case3.bij = np.array([-4.70588235, -8.0, -6.0], dtype=np.double)
    # imag part of the off-diagonal element in the ybus
    case3.gi = np.zeros(case3.nlines, dtype = np.double)
    case3.bi = np.zeros(case3.nlines, dtype = np.double)
    case3.gj = np.zeros(case3.nlines, dtype = np.double)
    case3.bj = np.zeros(case3.nlines, dtype = np.double)
    case3.xinv = np.array([1 / 0.2, 1 / 0.1, 1 / 0.15]) # inverse of
    the line reactances
    case3.ratio = np.array([1.0, 1.0, 1.0], dtype=np.double) #
    transformer ratio for all the lines
    case3.shift_angle = np.zeros(3) #shift angles of the transformers
    case3.ifrom = np.array([0, 0, 1], dtype=np.int32) # from-bus for
    all the lines
    case3.ito = np.array([1, 2, 2], dtype=np.int32) # to-bus for all
    the lines
    case3.ibstat = np.array([True, True, True], dtype=np.bool) #
    status for all the lines. True= in service
    return case3
@pytest.fixture
def case14():
    #Create a empty Case-object
    case14 = tf.Case()
    #Initialize the Case-object with the data of case14
    case14.icount = 0
    case14.convergence = False
    case14.sbase = 100
    #Bus data
    case14.nbuses = 14
    case14.slackbusnr = 0
    case14.numbus = np.array([1, 2,3,4,5,6,7,8,9,10,11,12,13,14])
    case14.buscod = np.array([3, 2, 2,1,1,2,1,2,1,1,1,1,1,1])
    case14.basekv = np.zeros(14, dtype = np.double)
    case14.gs = np.zeros(14, dtype = np.double)
```

    case14.bs \(=\) np.array \(([0,0,0,0,0,0,0,0,0.19,0,0,0,0,0])\)
    case14.area \(=\) np.ones (14, dtype \(=\) np.int32)
    case14.zone \(=\) np.ones(14, dtype \(=n p\).int 32 )
    case14.pload \(=\) np.array
    \(([0,0.217,0.942,0.478,0.076,0.112,0,0,0.295,0.09,0.035,0.061,0.135,0.149])\)
    case14.qload \(=\) np.array
    \(([0,0.127,0.19,-0.039,0.016,0.075,0,0,0.166,0.058,0.018,0.016,0.058,0.05])\)
    case14.vomag = np.array
    \(([1.06,1.045,1.01,1.019,1.02,1.07,1.062,1.09,1.056,1.051,1.057,1.055,1.05,1.0\)
    case14.voang = np.array
    ( \([0,-4.98,-12.72,-10.33,-8.78,-14.22,-13.37,-13.36,-14.94,-15.1,-14.79,-15.07\)
    /180*math.pi
    case14.pinj \(=\) np.zeros(case14.nbuses, dtype \(=n p . d o u b l e)\)
    case14.qinj \(=\) np.zeros(case14.nbuses, dtype \(=n p . d o u b l e)\)
    casel4.mismloc \(=\) np.zeros(2, dtype \(=\) np.int 32 )
    \#Generator data
    case14.ngens \(=5\)
    case14.genbus \(=\) np.array \(([0,1,2,-1,-1,3,-1,4,-1,-1,-1,-1,-1,-1])\)
    case14.genstat = np.array ([True, True, True, True, True])
    case14.genbase \(=\) np.array \(([100,100,100,100,100])\)
    case14.vgbus \(=\) np.array ([1.06,1.045,1.01,1.07,1.09])
    case14.pgen \(=\) np.array \(([2.324,0.40,0,0,0])\)
    case14. qgen \(=n p . \operatorname{array}([-0.169,0.424,0.234,0.122,0.174])\)
    case14.qmax \(=\) np.array \(([0.10,0.50,0.40,0.24,0.24])\)
    case14.qmin \(=\) np.array \(([0,-0.40,0,-0.06,-0.06])\)
    case14.pmax \(=\) np.array \(([3.324,1.40,1.00,1.00,1.00])\)
    case14.pmin \(=n p . z e r o s(5\), dtype \(=n p . d o u b l e)\)
    \#Branch data
    case14.nlines = 20
    casel4.ifrom \(=\) np.array
    \(([0,0,1,1,1,2,3,3,3,4,5,5,5,6,6,8,8,9,11,12]\), dtype \(=\) np.int 32\()\)
    case14.ito \(=\) np.array
    \(([1,4,2,3,4,3,4,6,8,5,10,11,12,7,8,9,13,10,12,13]\), dtype \(=\mathrm{np}\).
    int32)
    case14.gij = np.array([4.9991316,1.02589745,1.13501919,
    \(1.68603315,1.70113967,1.98597571,6.84098066,0.0,0.0,0.0\),
    1.95502856,
        \(1.52596744,3.0989274,0.0,0.0\),
        3.90204955, 1.42400549, 1.88088475, 2.48902459, 1.13699416])
        case14.bij = np.array([-15.26308652, -4.23498368, -4.78186315,
        -5.11583833, -5.1939274,-5.06881698, -21.57855398, -4.78194338,
        -1.79797907,
        -3.96793905, -4.09407434, -3.17596397,
        \(-6.10275545,-5.67697985,-9.09008272,-10.36539413,-3.02905046\),
        -4.40294375,
    $$
-2.25197463, \quad-2.31496348])
$$

case14.xinv = 1/np.array
( $[0.05917,0.22304,0.19797,0.17632,0.17388,0.17103,0.04211,0.20912,0.55618,0.2$
case14.b = np.array
$([0.0528,0.0492,0.0438,0.034,0.0346,0.0128,0,0,0,0,0,0,0,0,0,0,0,0,0,0])$

```
    case14.gi = np.zeros(20, dtype = np.double)
    case14.bi = np.zeros(20, dtype = np.double)
    case14.gj = np.zeros(20, dtype = np.double)
    case14.bj = np.zeros(20, dtype = np.double)
    case14.ratio = np.array
    ([1,1,1,1,1,1,1,0.978,0.969,0.932,1,1,1,1,1,1,1,1,1,1])
    case14.shift_angle = np.zeros(20, dtype = np.double)
    case14.ibstat = np.array([True,True,True,True,True,True,True,True
    ,True, True, True, True, True, True, True, True, True, True, True, True])
    case14.gii = np.array([6.02502906, 9.52132361, 3.1209949,
    10.51298952, 9.56801778, 6.57992341, 0.0, 0.0, 5.32605504,
    5.78293431, 3.83591332, 4.01499203, 6.72494615, 2.56099964])
    case14.bii = np.array([-19.44707021, -30.2721154,-9.82238013,
    -38.65417121, -35.53363946, -17.34073281,
    -19.54900595,-5.67697985, -24.09250638, -14.76833788, -8.49701809,
    -5.42793859, -10.66969355, -5.34401393])
```

    return case14
    def test_acsolve(case3, case14):
\#Test case3
res3 $=$ tf.loadflow (case 3 )
assert_array_almost_equal(res3.vomag, np.array ([0.88496373,
0.88444577, 1.0]), 6)
assert_array_almost_equal(res3.voang, np.array ([-0.0786143,
-0.0654881, 0.0]), 6)
\#Test case14
for flat in [False, True]:
\#Start with flat start or not; the result should be the same
res14 = tf.loadflow(case14, flat_start = flat)
assert_array_almost_equal(res14.vomag, np.array([1.06, 1.045,
1.01, 1.01767086, 1.01951386, 1.07, 1.06151954,
$1.09,1.05593173$,
$1.05098463,1.05690652,1.05518856$,
1.03552995]), 6)
assert_array_almost_equal(res14.voang, np.array([0.0,
$-0.08696258,-0.22209489,-0.17999408,-0.15313263,-0.24820233$,
-0.23316948,
$-0.23316948,-0.26072638,-0.26349739,-0.25814505,-0.26311858$,
-0.26452692,

$$
-0.27983988]), 6)
$$

## C. 6 test_acsolve_unit.py

1
2 import sys
3 sys.path.append("..")
4 import topflow as tf
5 import topflow.acsolve_wrapper as ac
6 import numpy as np
7 from numpy.testing import assert_array_equal
8 from numpy.testing import assert_allclose
9 import pytest
10 import math
11
\#This file provides the unit-tests of the sub-functions of acsolve
13
14 @pytest.fixture
15 def case3():
16 case3 = tf.Case()
17 case3.sbase $=100$ \#MVa base
18
19
20
21
\# Bus data
case3.nbuses $=3$ \# number of buses
case3.slackbusnr= 2 \#the internal number of the slacbus
case3.numbus $=$ np.array $([0,1,2]$, dtype=np.int 32$)$
case3.buscod $=$ np.array ([1, 1, 3], dtype=np.int32) \# bus type.
$1=P Q, \quad 2=P V, 3=$ Slack, $4=$ disconnected
case3.pload = np.array([1.0, 0.5, 0.0], dtype=np.double) \# pload
[nbuses]: active load vector (fixed)
case3.qload $=$ np.array ([0.5, 0.5, 0.0], dtype=np.double) \# qload [nbuses]: reactive load vector (fixed)
case3.pinj = np.zeros(case3.nbuses, dtype=np.double) \#
calculated active power injection for all buses
case3.qinj = np.zeros ( case3.nbuses, dtype=np.double) \#
calculated reactive power injection for all buses
case3.vomag = np.ones(case3.nbuses, dtype=np.double) \# voltage magnetudes at all buses
case3.voang $=$ np.zeros(case3.nbuses, dtype=np.double) \# voltage angles at all buses
\# Generator data
case3.genbus = np.array $([-1,-1,0]$, dtype=np.int32) \# genbus[ nbuses] if a generator: genbus[k]= generator nr. If not: genbus[k ] $=-1$
case3.vgbus $=-n p$. ones (case3.nbuses, dtype=np.int 32 )
case3.pgen = np.array([0.0], dtype=np.double) \# pgen[ngen]:
active power generated at generator nr
case3.qgen $=$ np.array([0.0], dtype=np.double) \# qgen[ngen]: reactive power generated at generator nr
\# Line data
case3.nlines $=3$ \# number of lines
case3.gii = np.array ([5.17647059, 3.17647059, 6.0], dtype=np.
double) \# real part of the diagonal element in the ybus
case3.bii = np.array([-12.70588235, -10.70588235, -14.0], dtype=
np.double) \# imag part of the diagonal element in the ybus
case3.gij = np.array([1.17647059, 4.0, 2.0], dtype=np.double) \#
real part of the off-diagonal element in the ybus
case3.bij = np.array ([-4.70588235, -8.0, -6.0], dtype=np.double)
\# imag part of the off-diagonal element in the ybus
case3.xinv = np.array ([1 / 0.2, 1 / 0.1, 1 / 0.15]) \# inverse of
the line reactances
case3.ratio = np.array([1.0, 1.0, 1.0], dtype=np.double) \#
transformer ratio for all the lines
case3.shift_angle = np.zeros(3) \#shift angles of the transformers
case3.ifrom = np.array ([0, 0, 1], dtype=np.int32) \# from-bus for
all the lines
case3.ito = np.array([1, 2, 2], dtype=np.int32) \# to-bus for all
the lines
case3.ibstat = np.array([True, True, True], dtype=np.bool) \#
status for all the lines. True= in service
return case3
def test_flatstart():
nbuses $=4$
buscod $=$ np.array ([1, 1, 3, 2])
vomag $=$ np.array ([1.01, 1.02, 1.03, 1.04])
voang = np.array ([0.1, 0.2, 0.3, 0.4])
vomag_ref $=$ np.array $([1.01,1.02,1.03,1.04])$
voang_ref $=$ np.array $([0.1,0.2,0.3,0.4])$
for flat_start in [False, True]:
ac.flatstart(nbuses, buscod, vomag, voang, flat_start)
if(flat_start):
assert_allclose(vomag, np.array ([1.00, 1.00, 1.03, 1.04]))
assert_allclose(voang, np.array ([0.0, 0.0, 0.3, 0.0]))
else:
assert_allclose(vomag, vomag_ref)
assert_allclose(voang, voang_ref)
def test_netinj(case3):
\#The NR-solution of case3:
case3.vomag $=$ np.array ([0.88496373, 0.88444577, 1.0])
case3.voang $=$ np.array ([-0.0786143, -0.0654881, 0.0])
for pqv in $[1,2,3]:$
\#set initial values for pinj and qinj as zero to check the
operation of pqv:
case3.pinj = np.zeros(case3.nbuses)
case3.qinj = np.zeros(case3.nbuses)
pinj_init $=$ np.zeros(case3.nbuses)
qinj_init $=$ np.zeros(case3.nbuses)
ac.netinj(pqv, case3.nbuses, case3.nlines, case3.pinj, case3.
qinj, case3.vomag, case3.voang, case3.gii, case3.bii, case3.gij,
case3.bij, case3.ratio, case3.shift_angle, case3.ifrom, case3.ito,
case3.ibstat)
\#The missmatch between calculated injections and known load
should be smaller than $10^{\wedge}-8$
if(pqv==1):
assert_allclose(case3.pinj[:2], -1*case3.pload[:2])
assert_allclose(case3.qinj, qinj_init)
elif(pqv == 2):
assert_allclose(case3.pinj, pinj_init)
assert_allclose(case3.qinj[:2], -1*case3.qload[:2])
else:
assert_allclose (case3.pinj[:2], -1*case3.pload[:2])
assert_allclose(case3.qinj[:2], -1*case3.qload[:2])
def test_missmatch():
\#Set up the test case
nbuses $=4$
buscod $=$ np.array $([1,1,3,2])$
genbus $=$ np.array $([-1,-1,0,1])$
pgen $=$ np.array $([0,0.010])$
qgen $=$ np.array $([0,-0.010])$
pload $=$ np.array $([0.010,0.020,0.0,0.0])$
qload $=$ np.array $([0.010,0.020,0.0,0.0])$
alfa $=$ beta $=$ np.zeros(4)
qdelta $=$ pdelta $=0.0$
\#net injections:
pinj $=$ np.array $([-0.010,0.0,0.100,0.005])$
qinj $=$ np.array $([0.0,-0.020,0.100,-0.010])$
\#the mismatches overwrites the net injections in the lists pinj
and qinj.
\#The exceptions are PV-buses in qinj and the slack-bus in both
pinj and qinj. These posisitons keeps the net injections
\# The mismatches should be xgen - xload - xinj for PV-buses and -
xload - xinj for $P Q$ buses.
\#Therefore, expected mismatches are:
pmism $=$ np.array $([0.0,-0.020,0.100,0.005])$
qmism $=$ np.array $([-0.010,0.0,0.100,-0.010])$
\#Test all variants of pqv:
for pqv in $[1,2,3]$ :
\#reset xinj values, and make copies
pinj $=$ np.array $([-0.010,0.0,0.100,0.005])$
qinj $=$ np.array $([0.0,-0.020,0.100,-0.010])$
pinj_init $=$ np.array $([-0.010,0.0,0.100,0.005])$
qinj_init $=n p . \operatorname{array}([0.0,-0.020,0.100,-0.010])$
\#Run the mismat function
ac.mismat (pqv, nbuses, genbus, buscod, pgen, qgen, pload,
qload, pinj, qinj, alfa, beta, pdelta, qdelta)
\#Excpected results:
if (pqv == 1):
assert_allclose(pinj, pmism)
assert_allclose(qinj, qinj_init)
elif(pqv == 2):
assert_allclose(pinj, pinj_init)
assert_allclose(qinj, qmism)
else:
assert_allclose(pinj, pmism)
assert_allclose(qinj, qmism)
def test_maxmism():
nbuses $=4$
buscod $=$ np.array $([2,3,1,-2])$
mismloc $=$ np.array $([0,0])$
pinj $=$ np.array $([0.0,0.150,-0.100,0.010])$
qinj $=$ np.array $([0.150,0.0,0.010,-0.100])$
for pqv in $[1,2,3]:$
mismloc $=$ np.array $([0,0])$
ac.maxmism(pqv, nbuses, buscod, mismloc, pinj, qinj)
if(pqv == 1):
assert_array_equal(mismloc, np.array([2,0]))
elif(pqv == 2):
assert_array_equal(mismloc, np.array([0,3]))
else:
assert_allclose(mismloc, np.array([2,3]))
def test_enforce_qlim():
nbuses $=6$
qmism $=1$ \#worst missmatch at 1
buscod $=$ np.array $([3,1,-2,-2,2,2])$
genbus $=$ np.array $([0,-1,1,2,3,4])$
numbus $=$ np.array $([0,1,2,3,4,5])$
pinj $=$ np.array ([1.1,0.2,0.3,0.0,0,0.1])
qinj $=$ np.array $([0,0.3,0.1,0.2,0,0])$ \#missmatches
qgen $=$ np.array $([1.4,-0.2,0.5,0.9,-0.3])$
qmax $=$ np.array $([10,0.5,0.5,0.5,0.5])$
qmin $=$ np.array $([-0.2,-0.2,-0.2,-0.2,-0.2])$
vomag $=$ np.array ([1.02,1.0, 0.92, 1.01, 1.03, 1.04])
vgbus $=$ np.array $([1.02,0.98,0.95,1.03,1.04])$
\#Can see that:
\#bus 2 is var-limited on the lower q-limit
\#bus 3 is var-limited on the upper q-limit
\#vomag[2]<vgbus[genbus[2]] and vomag[3]>vgbus[genbus[3]]
\#therefor both 2 and 3 should be put back as PV-buses
\#bus 4 violates the upper limit by 0.4
\#bus 5 violates the lover limit by -0.1
\#Therefor bus 4 and 5 should be put inn as var-limited buses
\#the new worst mismatch should be 0.4 at bus 4 after the function
call
qmism = ac.enforce_qlim(False, nbuses,qmism,buscod,genbus, numbus,
pinj,qinj,qgen, qmin, qmax, vomag, vgbus)
assert qmism == 4
assert_allclose (buscod, np.array([3,1,2,2,-2,-2]))
assert qgen[genbus[4]] == qmax[genbus[4]]
assert qgen[genbus[5]] == qmin[genbus[5]]
assert qinj[2] $==0.0$
assert qinj[3] $=0.0$
def test_bujac (case3):
melem $=5$ * (case3.nbuses +2 * case3.nlines) \#allocate enough
space for the sparse matrix
jacbi = np.zeros((melem), dtype=np.double)
jcol = -np.ones((melem), dtype=np.int32)
ipv = -np.ones((melem), dtype=np.int32)
ac.bujac (
case3.nbuses,
case3.nlines,
jacbi,
case3.vomag,
case3.voang,
case3.gii,
case3.bii,
case3.gij,
case3.bij,
case3.ratio,
jcol,
ipv,
case3.ifrom,
case3.ito,
case3.ibstat,
)
assert_allclose(jacbi, np.array([12.70588235, 10.70588235, 14.0,
12.70588235, 10.70588235, 14.0,
5.17647059, 3.17647059, 6.0,
$-4.70588235,-4.70588235,-1.17647059$,
-1.17647059, 1.17647059,

```
-5.17647059, 1.17647059, -3.17647059, -4.70588235,
                -4.70588235,-8.0,-8.0, -4.0,
-4.0, 4.0,4.0, -6.0, -8.0, -8.0, -6.0, -6.0,
                            -2.0, -2.0, 2.0, 2.0, -6.0,
-6.0, 0.0, 0.0, 0.0, 0.0, 0.0, 0.0, 0.0, 0.0, 0.0]))
```

    assert_allclose (jcol, np.array ([0, 1, 2, 3, 4, 5, 3, 4, 5, 1, 0,
    \(4,3,1,0,0,1,4,3,2,0,5,3,2\),
            \(0,2,5,3,2,1,5,4,2,1,5\),
    \(4,-1,-1,-1,-1,-1,-1,-1,-1,-1]))\)
    assert_allclose(ipv, np.array([6, 7 , 8, 13 , 15, 24, 9, 10, 20,
    \(11,12,19,28,14,17,16,18,23,32,21,22,-1,29,26\),
                                    25, 27, -1, 33, 30, 31, \(-1,-1,34\),
    \(35,-1,-1,-1,-1,-1,-1,-1,-1,-1,-1,-1])\) )
    def test_jacsize():
jcol = np.array ([1, 2, 3, 4, 5, 6, 7, -1, -1,-1,-1])
jacsize $=$ ac.jacsize(jcol)
assert jacsize == 7
def test_coo_conv(case3):
melem $=5$ * (case3.nbuses +2 * case3.nlines)
buscod = np.array ([1, 1, 3])
jacbi $=$ np.array ([12.70588235, 10.70588235, 14.0, 12.70588235,
10.70588235, 14.0,
$-4.70588235,-4.70588235,-1.17647059$,
-1.17647059, 1.17647059,
$-5.17647059,1.17647059,-3.17647059,-4.70588235$,
$-4.70588235,-8.0,-8.0,-4.0$,
$-4.0,4.0,4.0,-6.0,-8.0,-8.0,-6.0,-6.0$,
$-2.0,-2.0,2.0,2.0,-6.0$,
$-6.0,0.0,0.0,0.0,0.0,0.0,0.0,0.0,0.0,0.0])$
ipv $=$ np.array $([6,7,8,13,15,24,9,10,20,11,12,19$,
28, 14, 17, 16, 18, 23, 32, 21, 22, -1, 29, 26,
25, 27, $-1,33,30,31,-1,-1,34$,
35, $-1,-1,-1,-1,-1,-1,-1,-1,-1,-1,-1])$
jcol $=$ np.array $([0,1,2,3,4,5,3,4,5,1,0,4,3,1,0,0$,
$1,4,3,2,0,5,3,2$,
$0,2,5,3,2,1,5,4,2,1,5$,
$4,-1,-1,-1,-1,-1,-1,-1,-1,-1])$
isa $=-n p$. ones(melem, dtype $=n p . i n t 32)$
row $=n p . z e r o s((m e l e m)$, dtype $=n p . i n t 32)$
col = np.zeros((melem), dtype = np.int32)
data $=n p . z e r o s((m e l e m), ~ d t y p e=n p . d o u b l e)$
ac.coo_conv(case3.nbuses, case3.ngens, case3.buscod, isa, ipv,
row, jcol, col, jacbi, data)
284 def test_update_voltages():
$0,0,0,0,0,0,0])$ ) $0,0,0,0,0,0])$ )
$0.0,0.0,0.0,0.0,0.0,0.0,0.0$,
$0.0,0.0,0.0,0.0,0.0,0.0,0.0])$ )
def test_set_rhs():
nbuses $=4$
buscod = np.array ([1, $2,3,-2])$
for pqv in $[1,2,3]:$
if(pqv == 1):
rhs = np.zeros(nbuses) 0.4/1.04]))
elif(pqv == 2):
rhs $=$ np.zeros (nbuses)
-0.4/1.04])) else:
-0.4]))
nbuses $=4$
assert_array_equal (row, np.array ([0, 0, 0, 0, 1, 1, 1, 1, 2, 2, $2,2,3,3,3,3,0,0,0,0,0,0,0,0,0,0,0,0$,
assert_array_equal (col, np.array ([0, 2, 1, 3, 1, 3, 0, 2, 2, 1, $0,3,3,0,1,2,0,0,0,0,0,0,0,0,0,0,0,0$, $0,0,0,0,0,0,0,0,0,0,0$,
assert_allclose(data, np.array([12.70588235, 5.17647059, -4.70588235, -1.17647059, 10.70588235, 3.17647059, -4.70588235, -1.17647059,
12.70588235, 1.17647059, -5.17647059, -4.70588235,
$-3.17647059,-4.70588235,0.0,0.0,0.0$,
$0.0,0.0,0.0,0.0,0.0,0.0$,
$0.0,0.0,0.0,0.0,0.0,0.0$,
ngens $=3$ \#number of generators (var-limited included)
nlimgens $=1$ \#number of limited generators
vomag = np.array ([1.01, 1.02, 1.03, 1.04])
pinj = np.array ([0.1, 0.2, 0.0, 0.4]) \#mismatches
qinj $=$ np.array ([-0.1, 0.0, 0.0, -0.4]) \#mismatches
ac.set_rhs(pqv, nbuses, buscod, vomag, pinj, qinj, rhs) assert_allclose(rhs, np.array([0.1/1.01, 0.2/1.02, 0.0,
ac.set_rhs(pqv, nbuses, buscod, vomag, pinj, qinj, rhs) assert_allclose(rhs, np.array([-0.1/1.01, 0.0, 0.0,
rhs $=$ np.zeros $(2 *$ nbuses - ngens + nlimgens -1$)$
ac.set_rhs(pqv, nbuses, buscod, vomag, pinj, qinj, rhs)
assert_allclose(rhs, np.array ([0.1, 0.2, 0.4, -0.1,

```
    ngens = 3 #number of generators (var-limited included)
    nlimgens = 1 #number of limited generators
    buscod = np.array([1,2,3,-2])
    vomag_init = np.array([1.01, 1.02, 1.03, 1.04])
    voang_init = np.array([0.1, 0.2, 0.3, 0.4])
    for pqv in [1,2,3]:
        vomag = np.array([1.01, 1.02, 1.03, 1.04])
        voang = np.array([0.1, 0.2, 0.3, 0.4])
        if(pqv == 1):
            #the correction vector for the FDLF should be of
dimension: nbuses
        #all buses are included (also slack and generators)
        correction = np.array([0.1,0.2,0.3, 0.4])
        ac.update_voltages(pqv, nbuses, buscod, vomag, voang,
correction)
            assert_allclose(vomag, vomag_init)
            assert_allclose(voang, np.array([0.2, 0.4, 0.3, 0.8]))
        elif(pqv == 2):
    #the correction vector for the FDLF should be of
dimension: nbuses
    #all buses are included (also slack and generators)
    correction = np.array([0.01,0.02,0.03, 0.04])
    ac.update_voltages(pqv, nbuses, buscod, vomag, voang,
correction)
            assert_allclose(vomag, np.array([1.02, 1.02, 1.03, 1.08])
)
    assert_allclose(voang, voang_init)
        else:
            #the correction vector for the NR should be of dimension
2*nbuses-ngens+nlimgens(var-limited generators)
    #slack bus is not part of correction-vector, neither is
generors for the magnitude-corrections)
    correction = np.array([0.1, 0.2, 0.4, 0.01, 0.04])
    ac.update_voltages(pqv, nbuses, buscod, vomag, voang,
correction)
    assert_allclose(vomag, np.array([1.02, 1.02, 1.03, 1.08])
)
    assert_allclose(voang, np.array([0.2, 0.4, 0.3, 0.8]))
```


## C. 7 test_case.py

```
1
import sys
3
4 sys.path.append("..")
5
6 \text { import math}
7 \text { import topflow as tf}
import numpy as np
9 from numpy.testing import assert_array_equal
10 from numpy.testing import assert_allclose
11
#This file contains the unit-tests of the Case-calss
13
14 def test_loadxl():
15 case14 = tf.example_case('case14')
16 assert case14.icount == 0
17 assert case14.convergence == False
18 assert case14.sbase == 100
1 9
20
21
22
23
24 assert_array_equal(case14.buscod , np.array
    ([3,2,2,1,1,2,1,2,1,1,1,1,1,1]))
    assert_array_equal(case14.basekv , np.zeros(14))
    assert_array_equal(case14.gs , np.zeros(14))
    assert_array_equal(case14.bs , np.array
    ([0,0,0,0,0,0,0,0,0.19,0,0,0,0,0]))
    assert_array_equal(case14.area , np.ones(14, dtype = np.int32))
    assert_allclose(case14.zone , np.ones(14, dtype = np.int32))
    assert_allclose(case14.pload , np.array
    ([0,0.217,0.942,0.478,0.076,0.112,0,0,0.295,0.09,0.035,0.061,0.135,0.149])
    )
    assert_allclose(case14.qload , np.array
    ([0,0.127,0.19,-0.039,0.016,0.075,0,0,0.166,0.058,0.018,0.016,0.058,0.05])
    )
    assert_allclose(case14.vomag , np.array([1.06 , 1.045 ,
    1.01 , 1.01767085, 1.01951386,
    1.09 , 1.05593172, 1.05098463,
    1.05038171, 1.03552995]))
    assert_allclose(case14.voang , np.array([ 0.
    -0.08696259, -0.22209489, -0.17999408, -0.15313264,
        -0.24820234, -0.23316948,
```

```
        -0.23316948, -0.26072638, -0.26349739,
        -0.25814505, -0.26311859,
        -0.26452692, -0.27983989]))
        #Generator data
        assert case14.ngens == 5
        assert_array_equal(case14.genbus , np.array
        ([0,1,2,-1,-1,3,-1,4,-1,-1,-1,-1,-1,-1]))
        assert_array_equal(case14.genstat , np.array([True,True,True,True
        ,True]))
        assert_array_equal(case14.genbase , np.array
        ([100,100,100,100,100]))
        assert_allclose(case14.vgbus , np.array
        ([1.06,1.045,1.01,1.07,1.09]))
        assert_allclose(case14.pgen , np.array([2.324,0.40,0,0,0]))
        assert_allclose(case14.qgen , np.array
        ([-0.169,0.424,0.234,0.122,0.174]))
        assert_array_equal(case14.qmax , np.array
        ([0.10,0.50,0.40,0.24,0.24]))
        assert_array_equal(case14.qmin , np.array
        ([0,-0.40,0,-0.06,-0.06, ]))
    assert_allclose(case14.pmax , np.array
    ([3.324,1.40,1.00,1.00,1.00]))
    assert_array_equal(case14.pmin , np.zeros(5))
        #Branch data
        case14.nlines = 20
        assert_array_equal(case14.ifrom , np.array
        ([0,0,1,1,1,2,3,3,3,4,5,5,5,6,6,8,8,9,11,12]))
        assert_array_equal(case14.ito , np.array
        ([1,4,2,3,4,3,4,6,8,5,10,11,12,7,8,9,13,10,12,13]) )
        assert_allclose(case14.xinv , 1/np.array
        ([0.05917,0.22304,0.19797,0.17632,0.17388,0.17103,0.04211,0.20912,0.55618,0.2
        )
        assert_allclose(case14.b , np.array
        ([0.0528,0.0492,0.0438,0.034,0.0346,0.0128,0,0,0,0,0,0,0,0,0,0,0,0,0,0])
        )
        assert_array_equal(case14.gi , np.zeros(20))
        assert_array_equal(case14.bi , np.zeros(20))
        assert_array_equal(case14.gj , np.zeros(20))
        assert_array_equal(case14.bj , np.zeros(20))
        assert_allclose(case14.ratio , np.array
        ([1,1,1,1,1,1,1,0.978,0.969,0.932,1,1,1,1,1,1,1,1,1,1]))
        assert_array_equal(case14.shift_angle , np.zeros(20))
        assert_array_equal(case14.ibstat , np.array([True,True,True,True,
        True, True, True, True, True, True, True, True, True, True, True, True, True,
        True,True,True]))
```


## D setup.py

```
1 import os
2 from pathlib import Path
3 from setuptools import setup, find_packages
4 #This is the setup-file of Topflow
5 #It specifies dependencies,
6
7 setup(
8 name="topflow",
    version="0.0.1",
    packages=find_packages('topflow'),
    # Project uses the following python packages, so ensure that
    these get
    # installed or upgraded on the target machine:
    install_requires=["numpy", "scipy", "scons", "openpyxl"],
    # metadata to display on PyPI
    author="aasmunsa",
    author_email="aa.selen@hotmail.com",
    description="Load flow",
21
22 )
23
24 #dir_path = str(Path(
```

$\qquad$

``` file
``` \(\qquad\)
``` ). parent.absolute())
25 #os.chdir(dir_path + "/topflow")
26 #os.system('scons')
```


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